

On the Weak Line Cover Numbers

Oksana Firman

Alexander Wolff

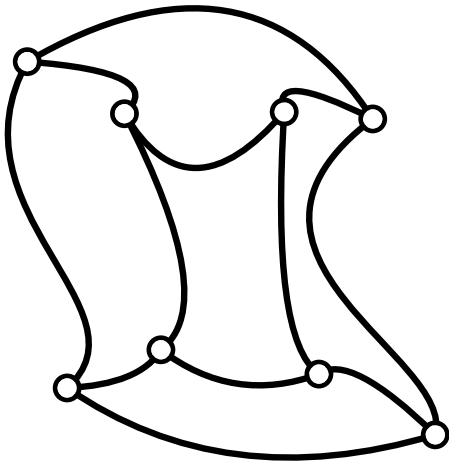
Julius-Maximilians-Universität Würzburg, Germany

Alexander Ravsky

Pidstryhach Institute for Applied Problems of Mechanics and Mathematics,
National Academy of Science of Ukraine, Lviv, Ukraine

Visual complexity of a drawing of a graph

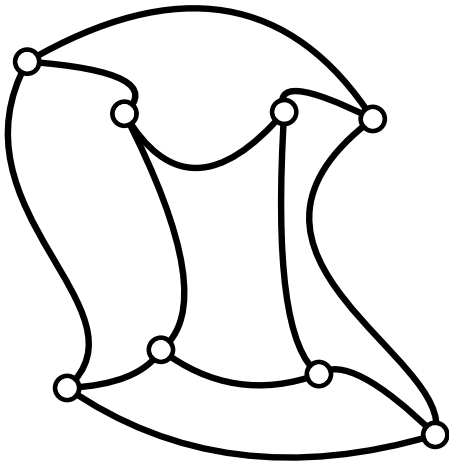
Given: graph G



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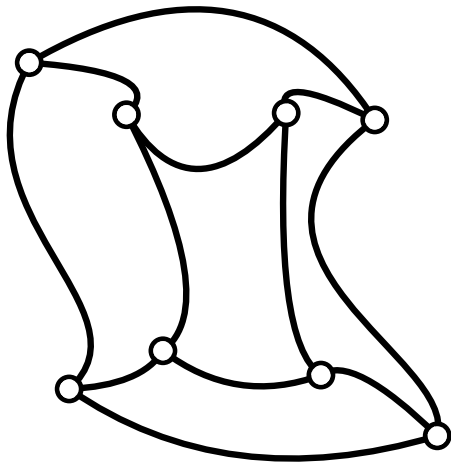
Use as few **objects** as possible to draw G



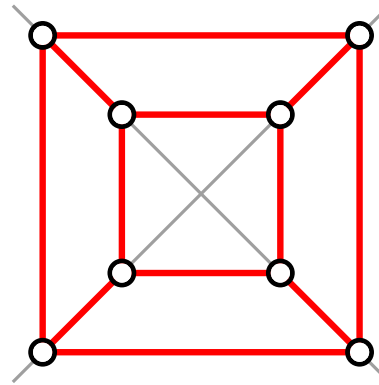
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12 segments

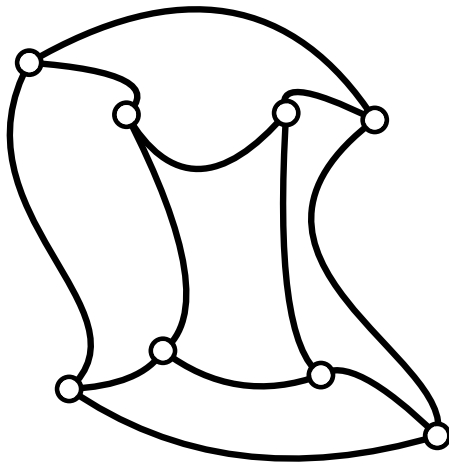


[Dujmović et al., CGTA'07]

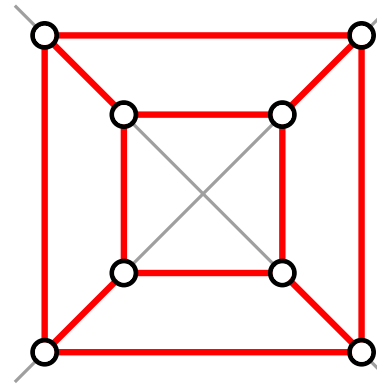
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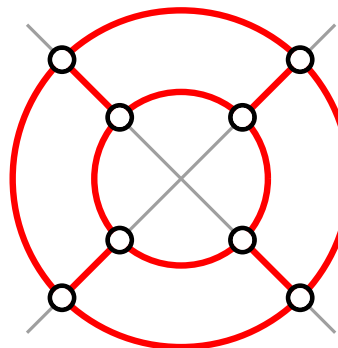


12 segments



[Dujmović et al., CGTA'07]

6 arcs

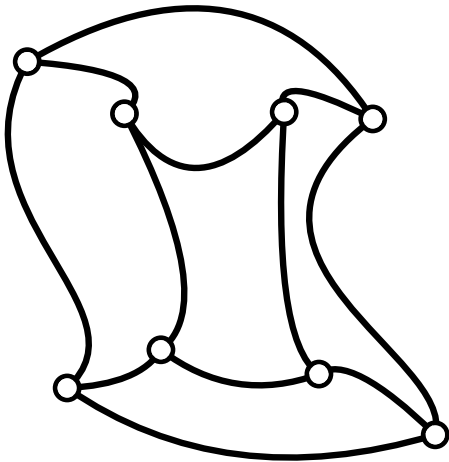


[Schulz, JGAA'15]

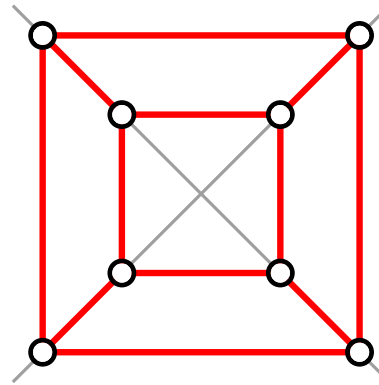
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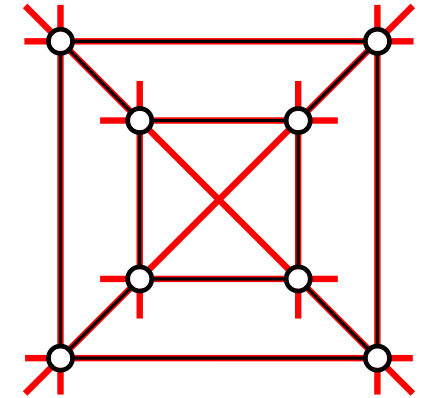


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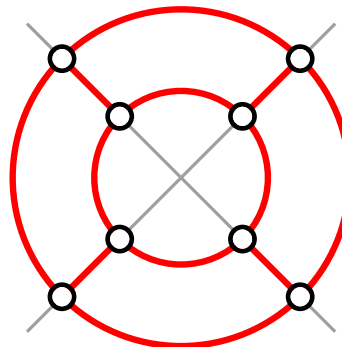
[Dujmović et al., CGTA'07]

10 straight lines



[Chaplick et al., GD'16]

6 arcs

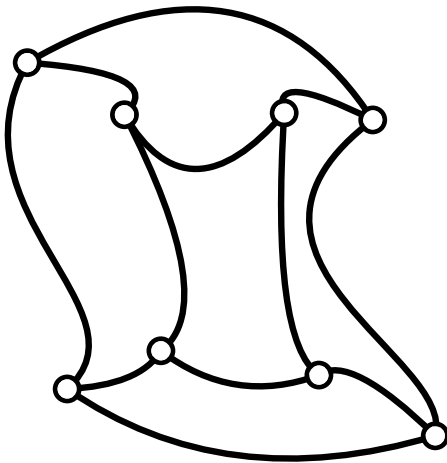


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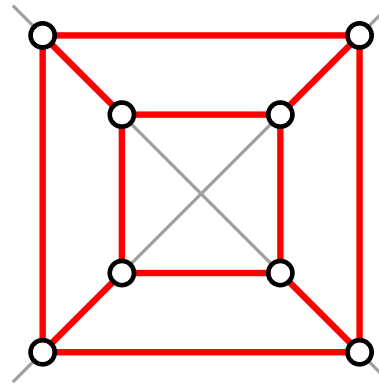
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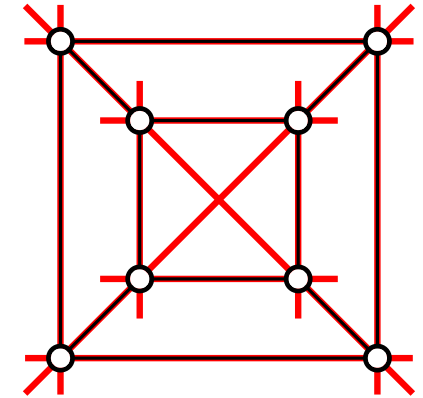


12 segments



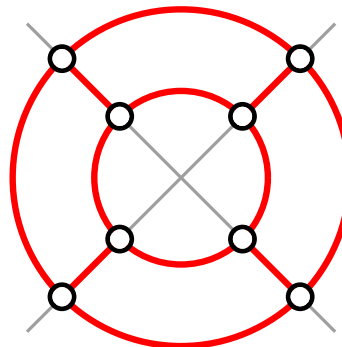
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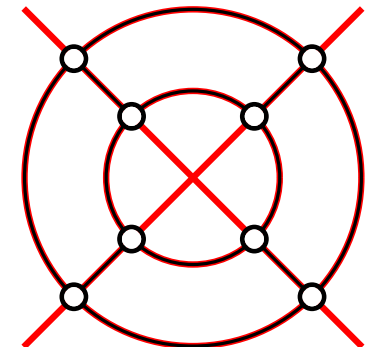
[Chaplick et al., GD'16]

6 arcs



[Schulz, JGAA'15]

4 circles

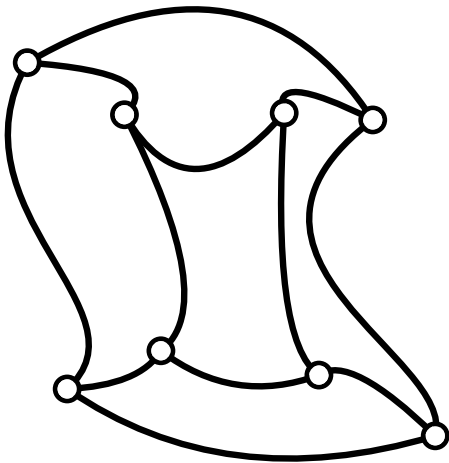


[Kryven et al., CALDAM'18]

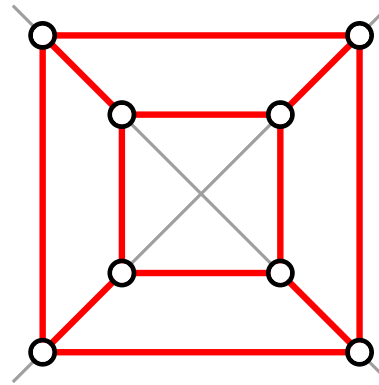
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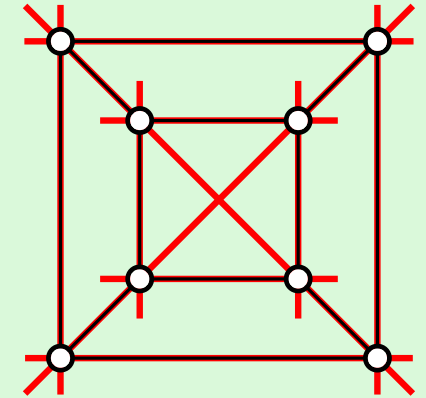


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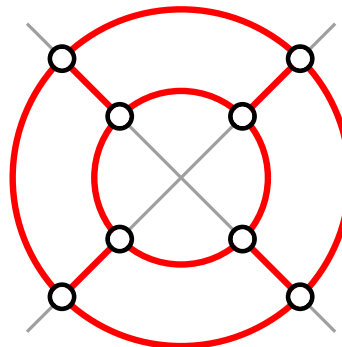
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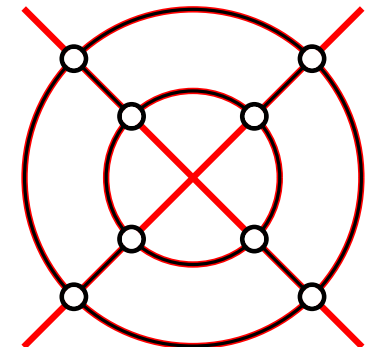
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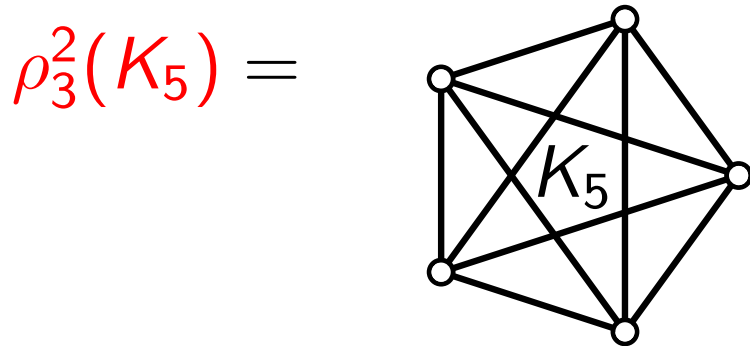
Affine cover numbers

[Chaplick et al., GD'16]

Let G be a graph and let $1 \leq m < d$.

All drawings are straight-line and crossing-free.

Def. The m -dimensional affine cover number $\rho_d^m(G)$ is the minimum number of m -dimensional planes in \mathbb{R}^d such that the vertices and the edges of a drawing of G are contained in the union of these planes.



Affine cover numbers

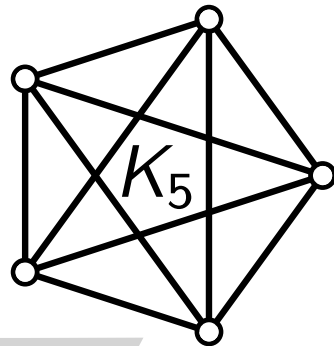
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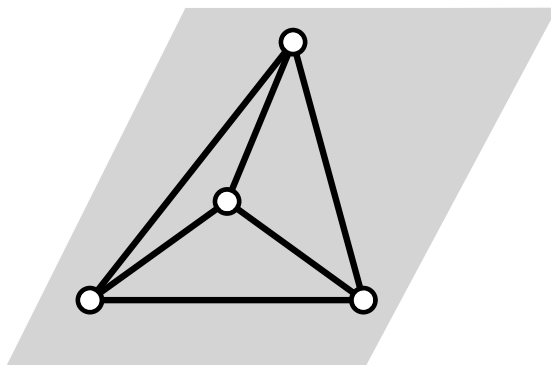
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$$\rho_3^2(K_5) =$$



○



Affine cover numbers

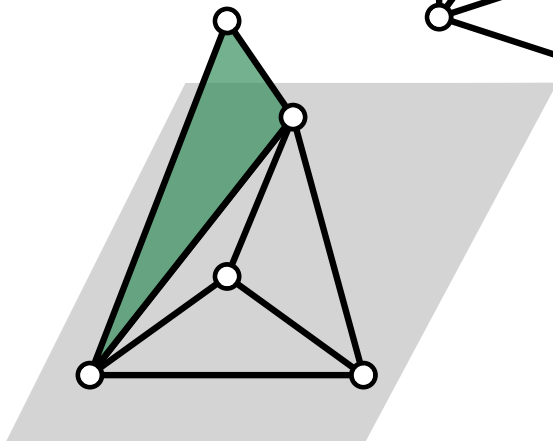
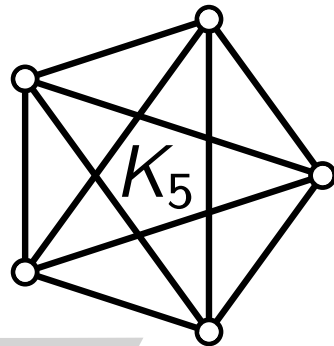
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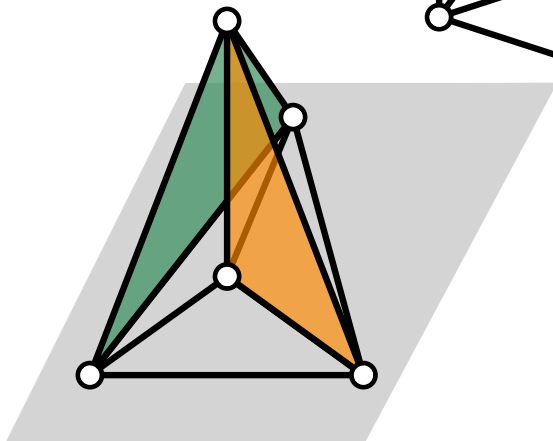
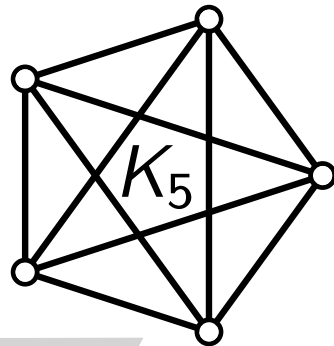
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$$\rho_3^2(K_5) = 3$$



Affine cover numbers

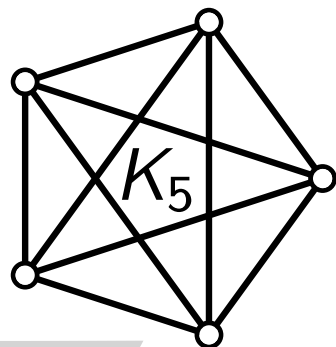
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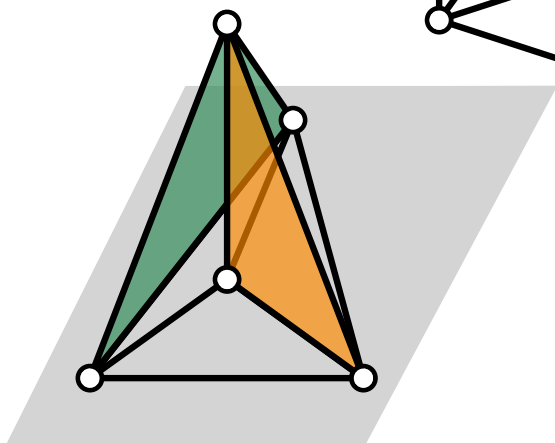
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$$\rho_3^1(K_5) =$$



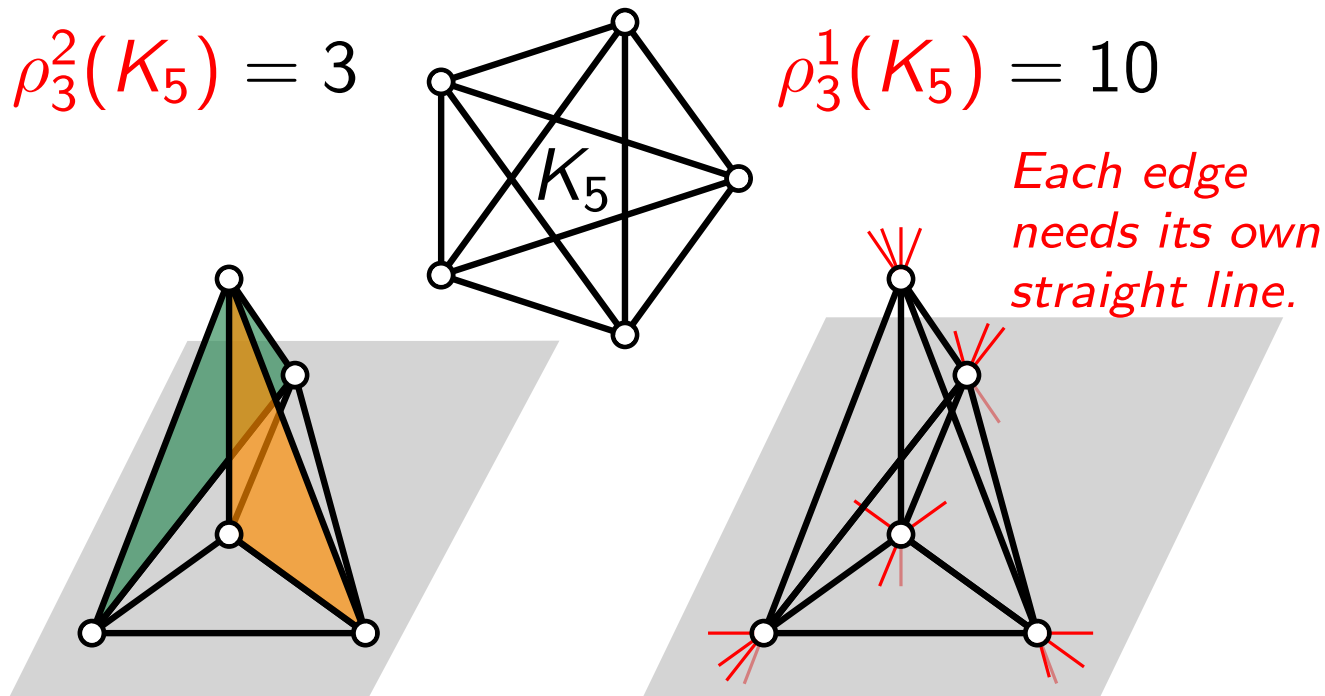
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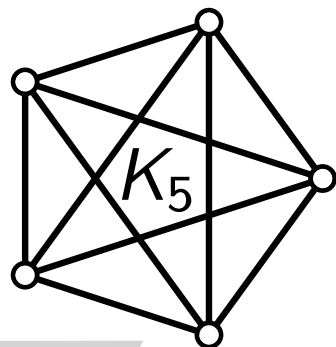
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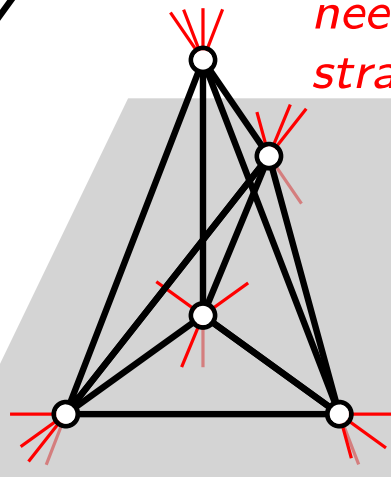
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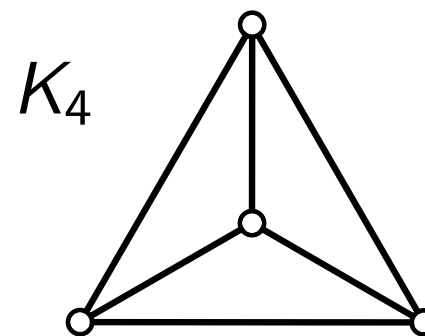


$$\rho_3^1(K_5) = 10$$

Each edge
needs its own
straight line.



$$\rho_2^1(K_4) =$$



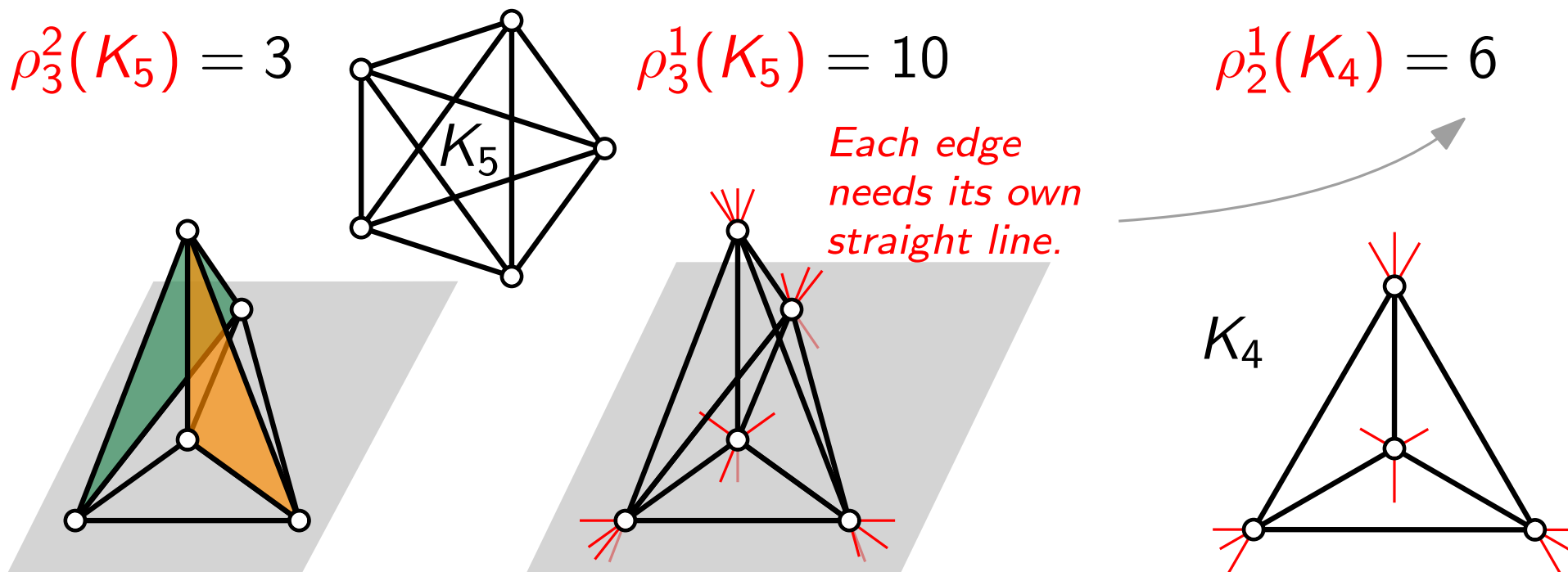
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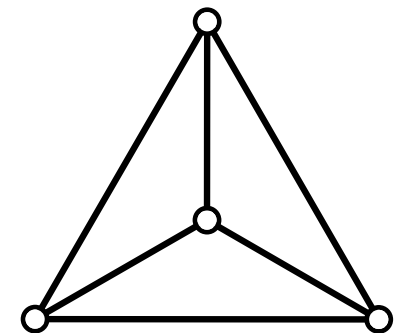
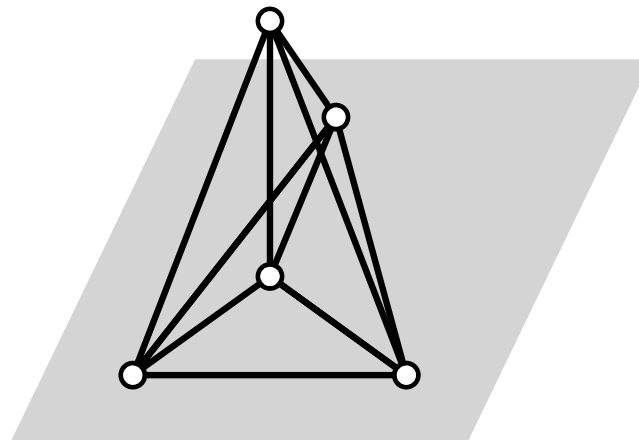
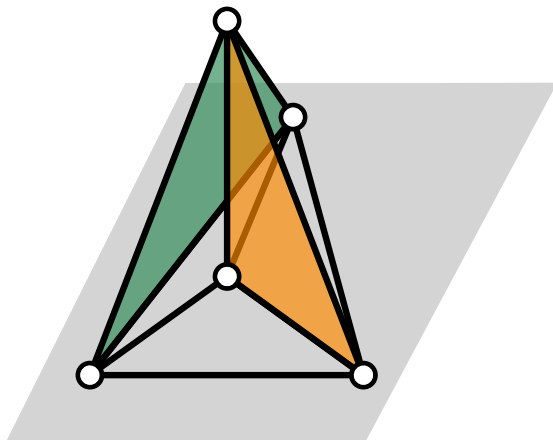
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$$\rho_3^2(K_5) = 3$$

$$\rho_3^1(K_5) = 10$$

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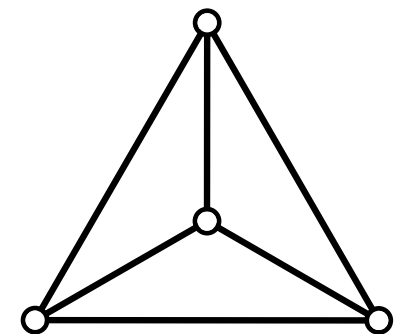
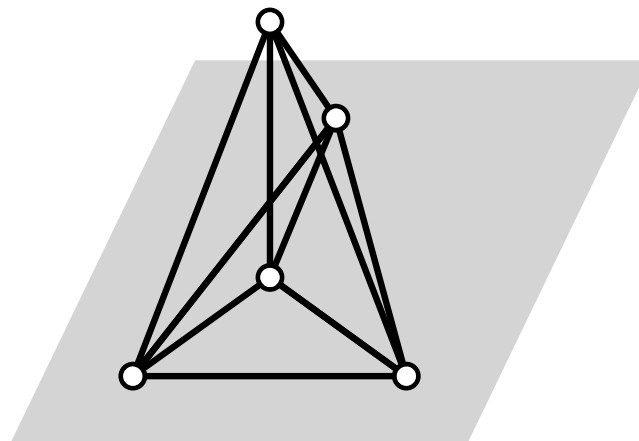
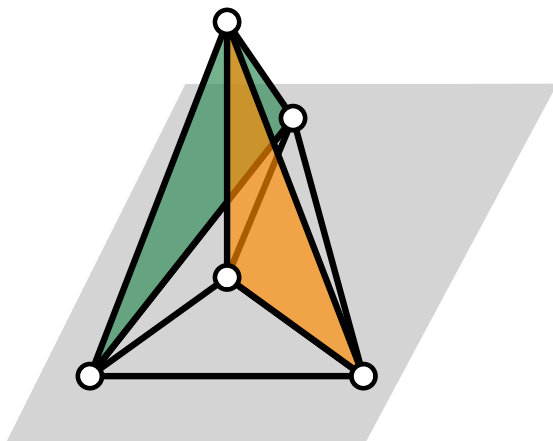
All drawings are straight-line and crossing-free.

Def. The ^{weak} m -dimensional affine cover number $\pi_d^m(G)$ ~~$\rho_d^m(G)$~~ is the minimum number of m -dimensional planes in \mathbb{R}^d such that the vertices ~~and the edges~~ of a drawing of G are contained in the union of these planes.

$$\rho_3^2(K_5) = 3$$

$$\rho_3^1(K_5) = 10$$

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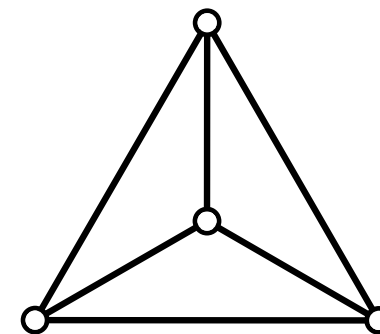
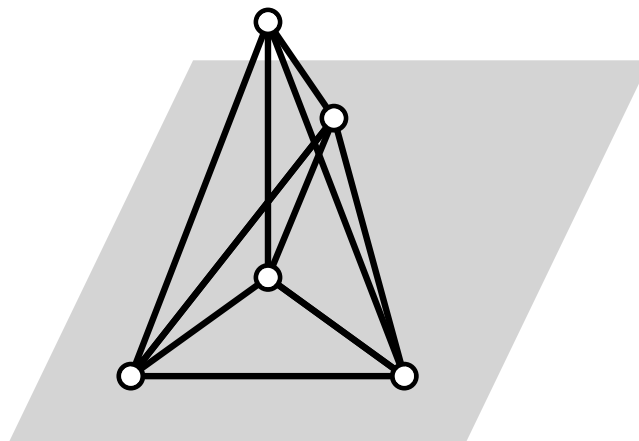
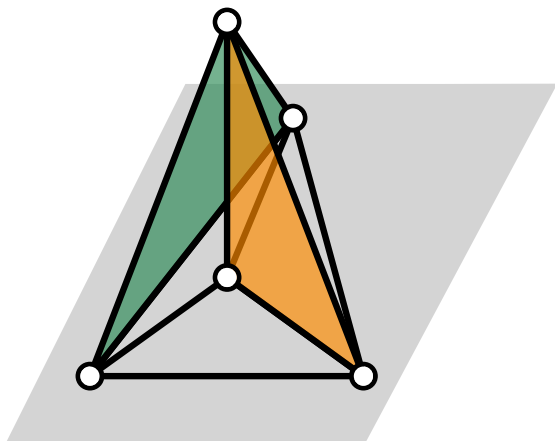
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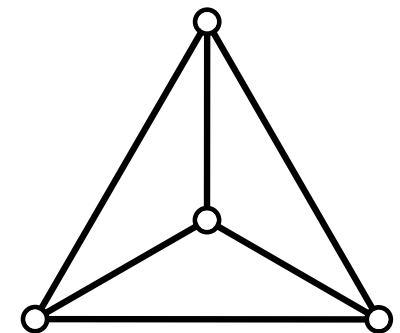
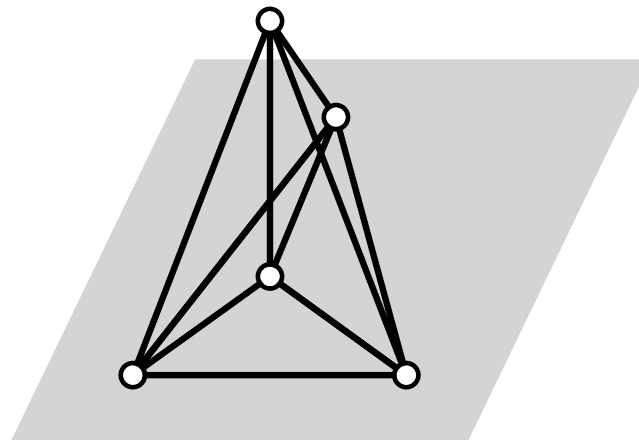
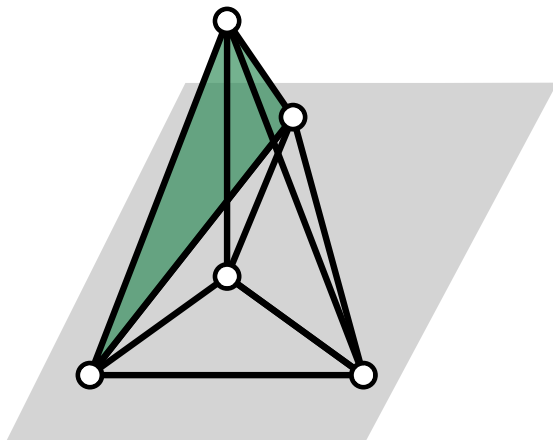
$$\pi_3^2(K_5) = 2$$

$$\rho_3^1(K_5) = 10$$

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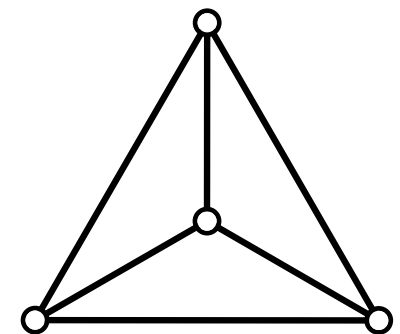
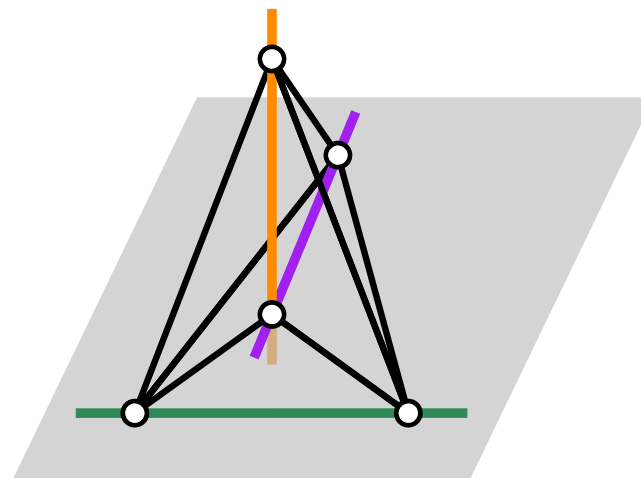
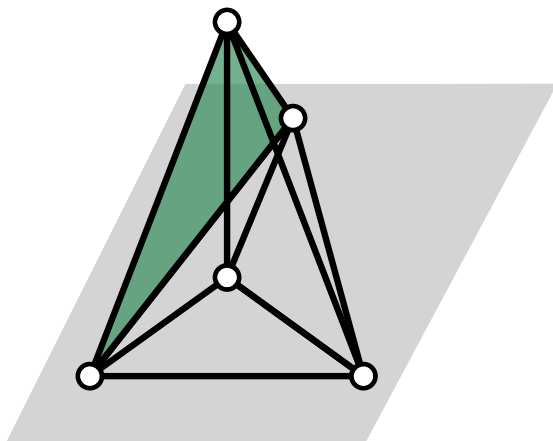
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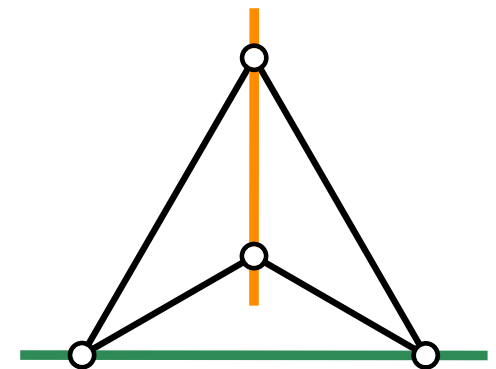
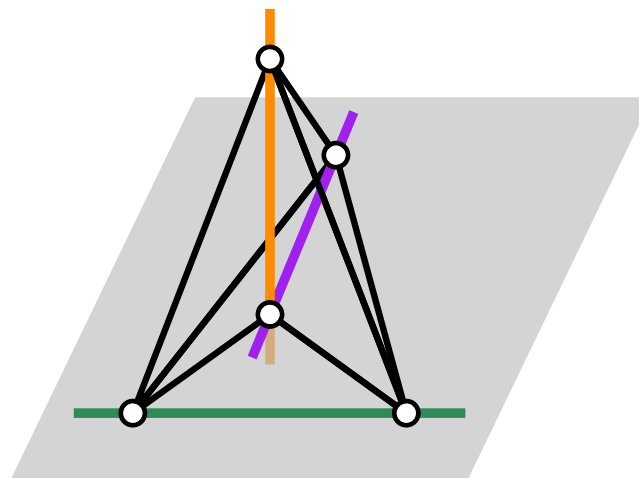
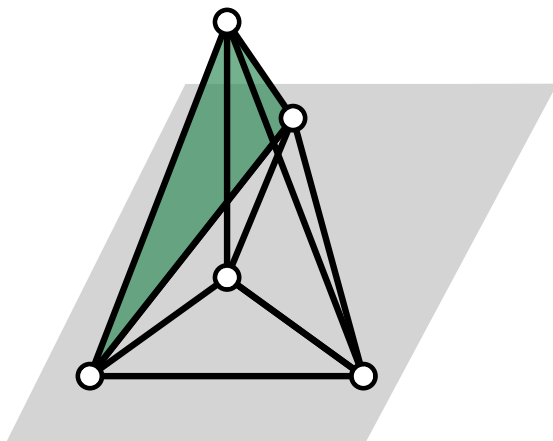
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$$\pi_2^1(K_4) = 2$$



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$$\rho_3^2$$

$$\pi_3^2$$

$$\rho_3^1$$

$$\pi_3^1$$

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Thm. Collapse of the Multidimensional Affine Hierarchy

For any integers $1 \leq l < 3 \leq d$ and for any graph G , it holds that $\pi_d^l(G) = \pi_3^l(G)$ and $\rho_d^l(G) = \rho_3^l(G)$.

Complexity of affine cover numbers

[Chaplick et al., WADS'17]

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ρ_3^2 NP-hard

ρ_3^1 NP-hard

ρ_2^1 NP-hard

π_3^2 NP-hard

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π_2^1

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Let G be a graph and let $1 \leq m < d$.

All drawings are straight-line and crossing-free.

Def. The ^{weak} m -dimensional affine cover number $\pi_d^m(G)$ is the minimum number of m -dimensional planes in \mathbb{R}^d such that the vertices ~~and the edges~~ of a drawing of G are contained in the union of these planes.

ρ_3^2 NP-hard

ρ_3^1 NP-hard

ρ_2^1 NP-hard

π_3^2 NP-hard

π_3^1 NP-hard

π_2^1 ? OPEN

Thm. Collapse of the Multidimensional Affine Hierarchy

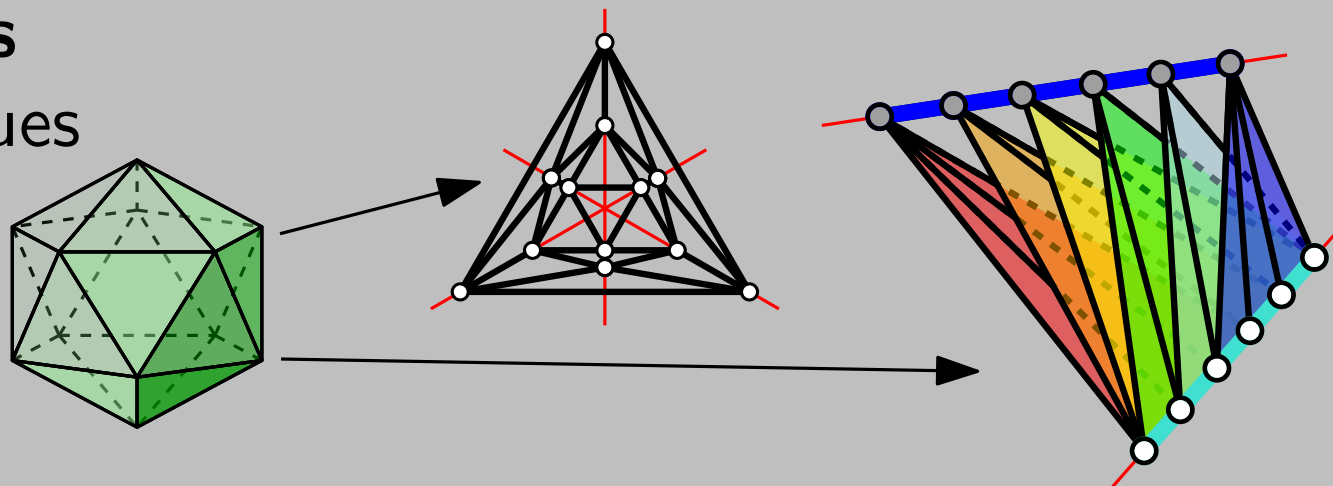
For any integers $1 \leq l < 3 \leq d$ and for any graph G , it holds that $\pi_d^l(G) = \pi_3^l(G)$ and $\rho_d^l(G) = \rho_3^l(G)$.

Overview

- Notation

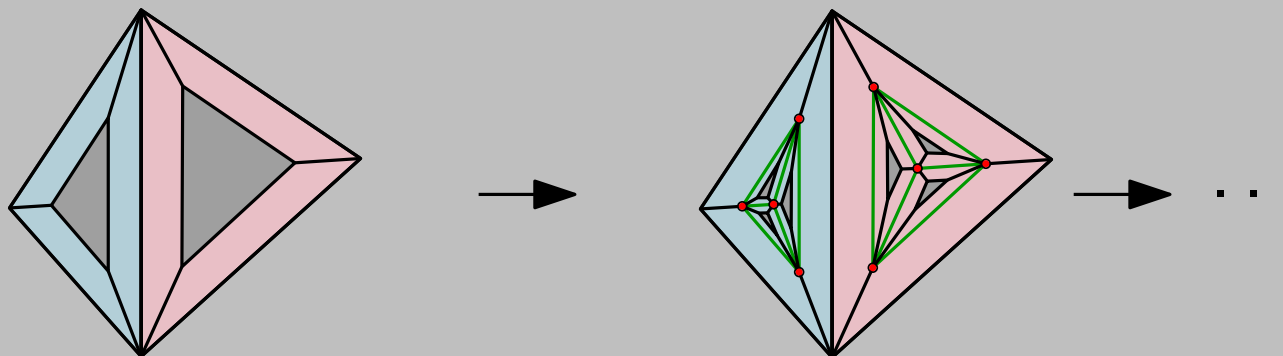
- **Platonic solids**

π_2^1 - and π_3^1 -values



- **Main contribution**

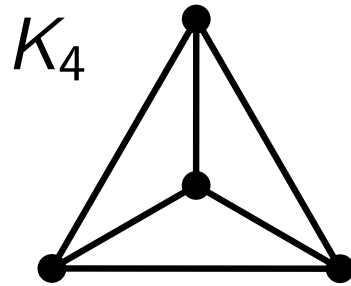
Infinite family of planar graphs with unbounded π_2^1 -value



Notation

Linear vertex arboricity $\text{lva}(G)$ of a graph G :
smallest size r of a partition of $V(G) = V_1 \cup \dots \cup V_r$
such that every V_i induces a linear forest.

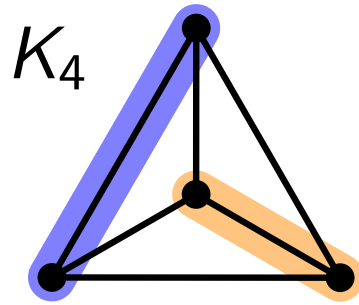
$$\text{lva}(K_4) =$$



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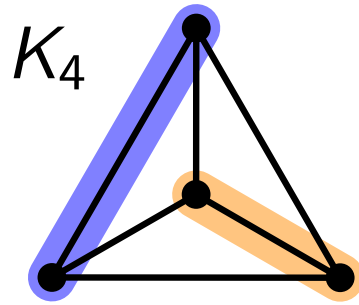
$$\text{lva}(K_4) = 2$$



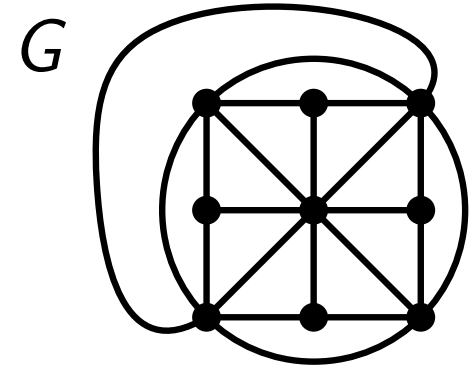
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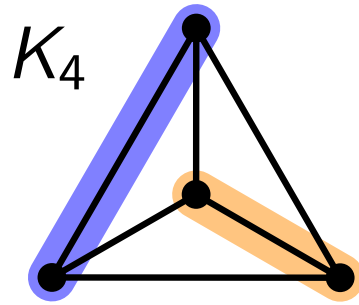
$$\text{lva}(G) =$$



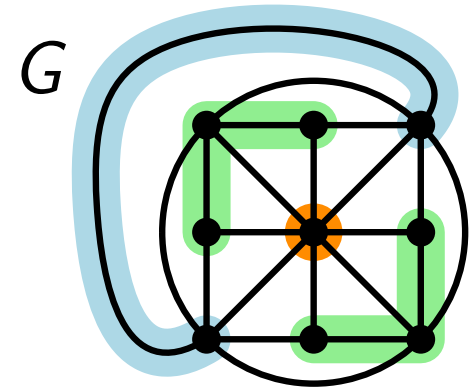
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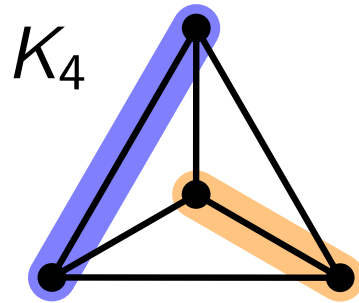
$$\text{lva}(G) = 3$$



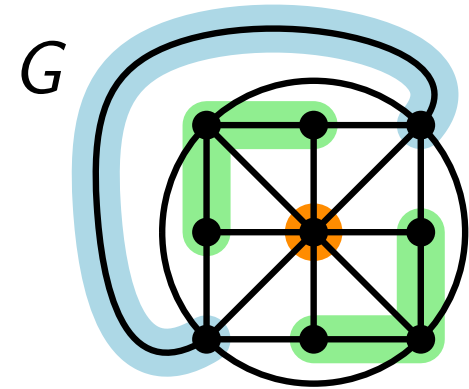
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$$\text{lva}(K_4) = 2$$



$$\text{lva}(G) = 3$$



Treewidth $\text{tw}(G)$ of a graph G :

upper bound

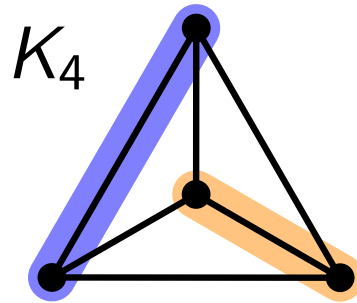
$$\text{tw}(G) \leq k$$

if G is a subgraph of a k -tree.

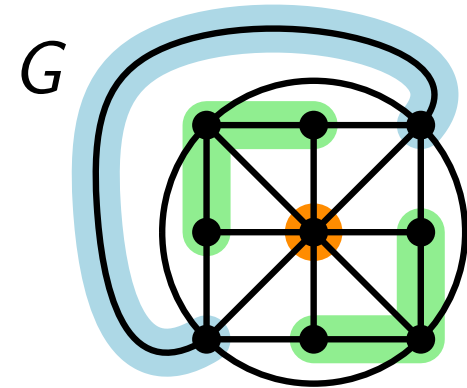
Notation

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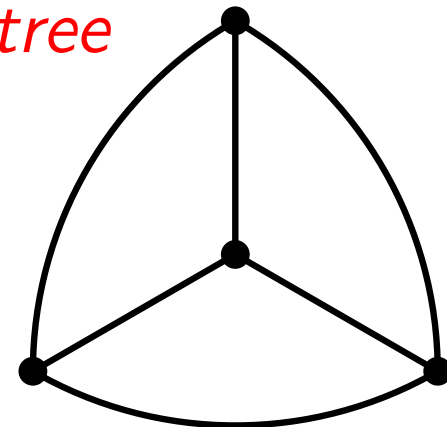
Treewidth $\text{tw}(G)$ of a graph G :

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3-tree

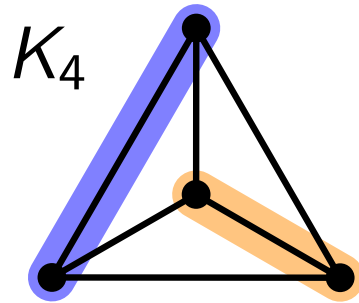


start
with
 $K_4 \dots$

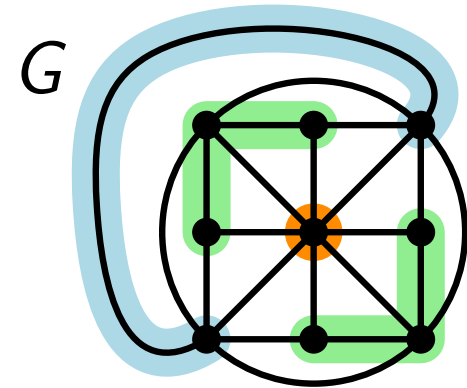
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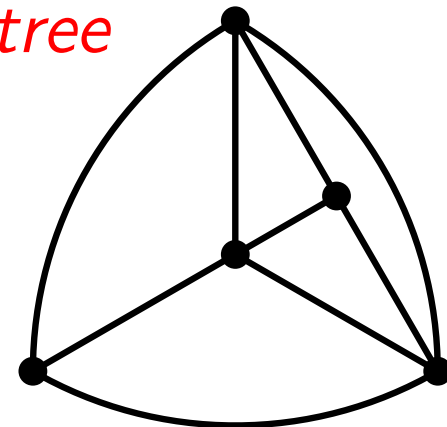
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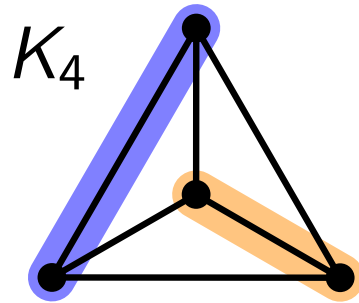


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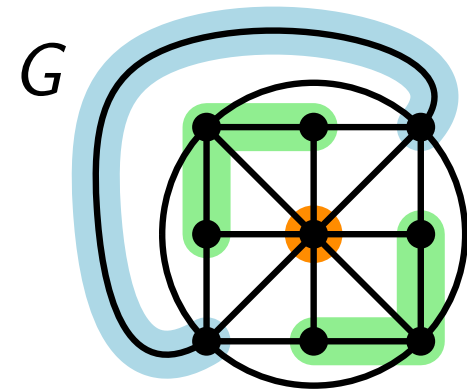
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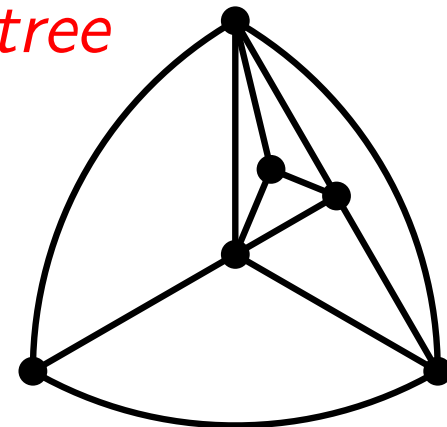
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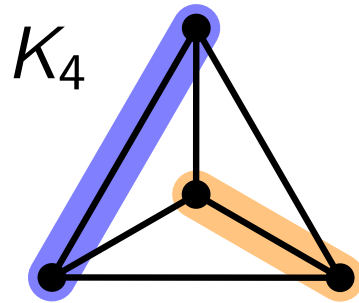


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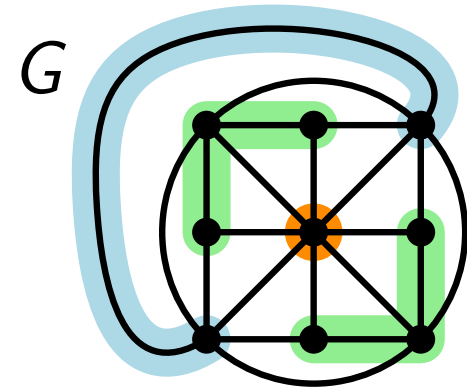
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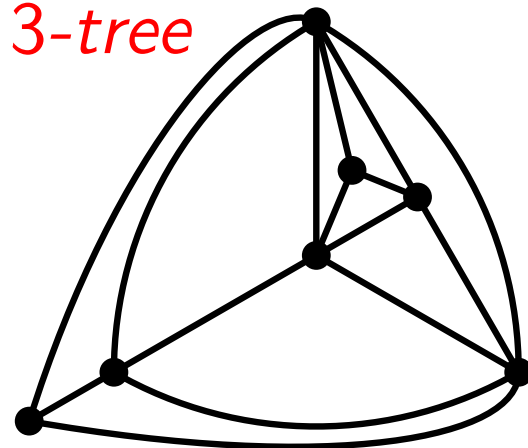
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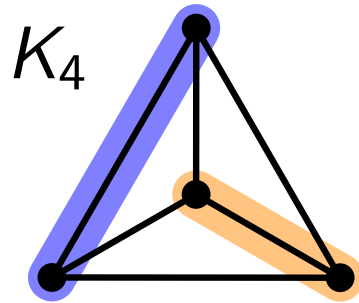


start
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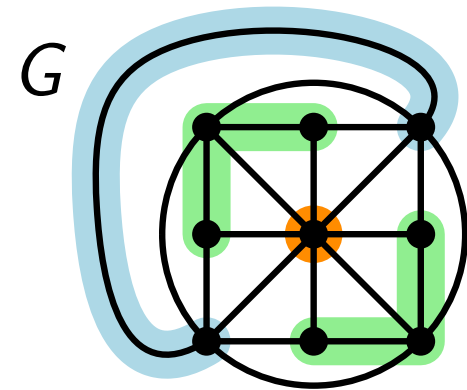
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Linear vertex arboricity $\text{lva}(G)$ of a graph G :
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$$\text{lva}(K_4) = 2$$



$$\text{lva}(G) = 3$$



Treewidth $\text{tw}(G)$ of a graph G :

upper bound

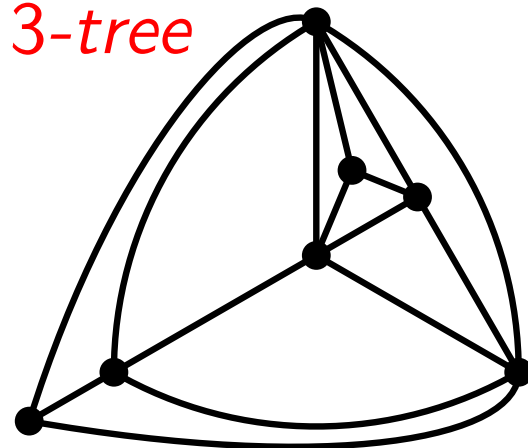
$$\text{tw}(G) \leq k$$

if G is a subgraph of a k -tree.

lower bound

$$\text{tw}(G) \geq \text{mindeg}(G).$$

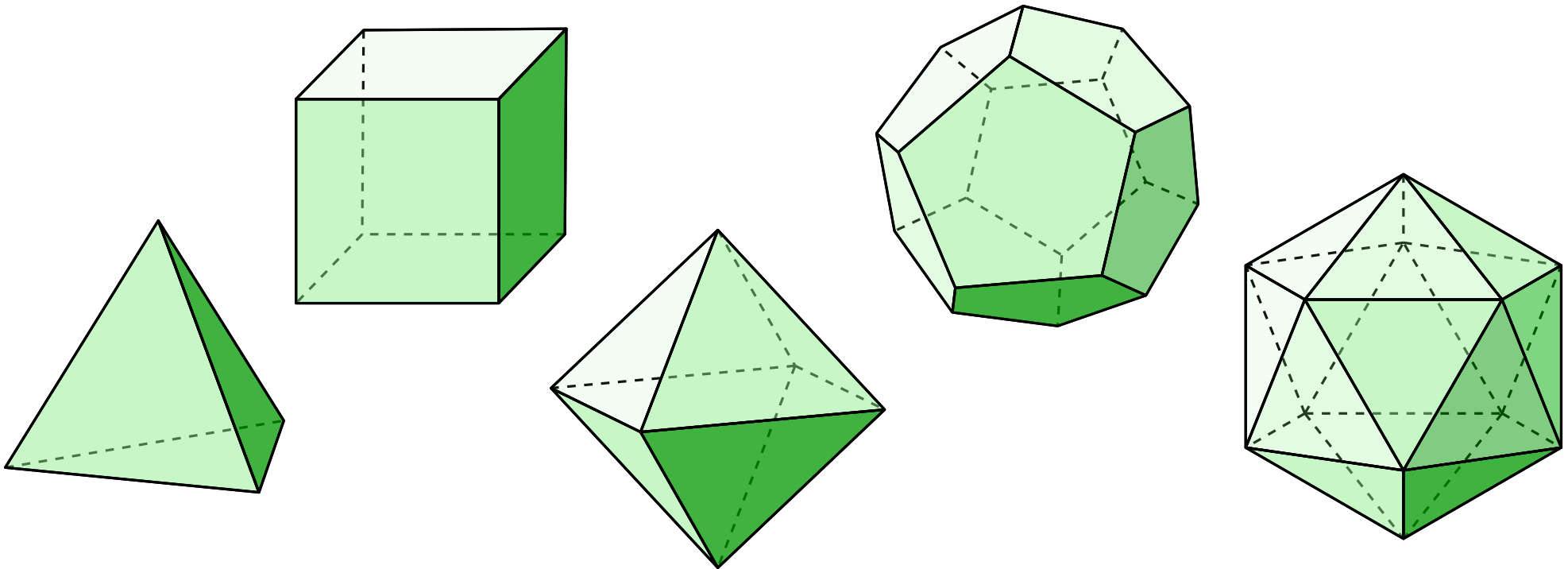
3-tree



start
with
 $K_4 \dots$

Platonic solids

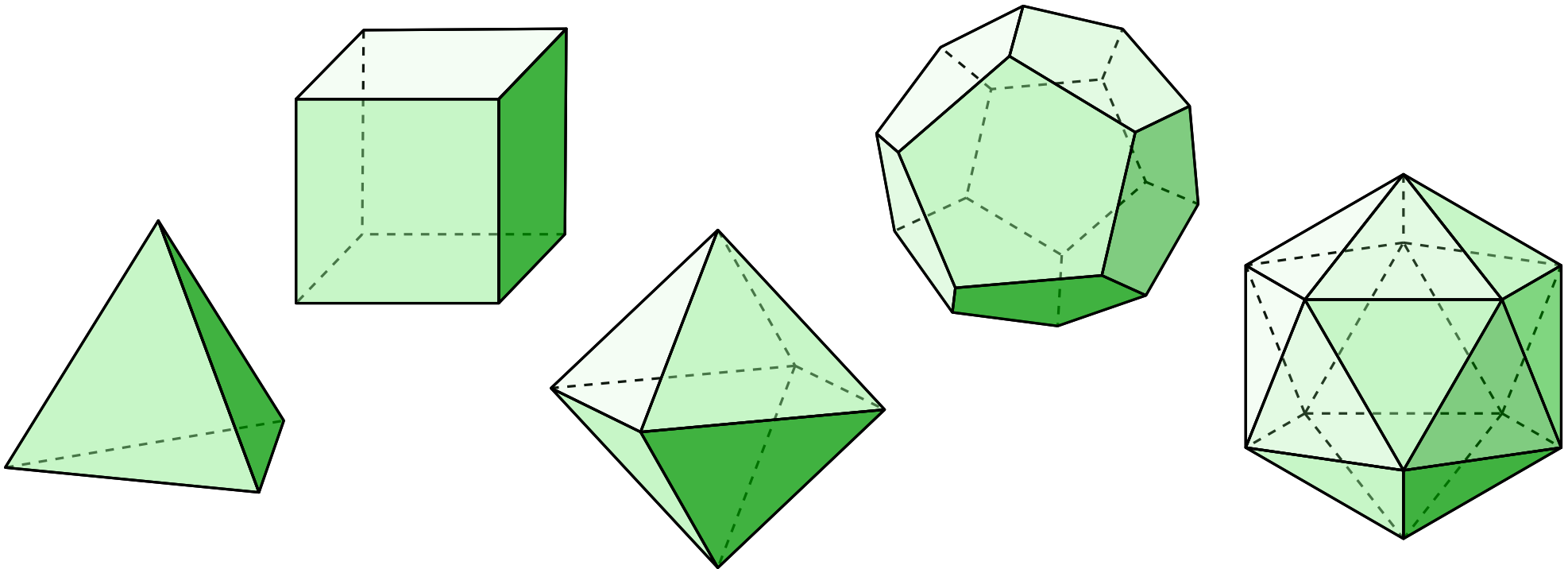
$G = (V, E)$	$ V $	$ E $	$ F $	$\rho_2^1(G)$	$\rho_3^1(G)$	$\pi_2^1(G)$	$\pi_3^1(G)$
tetrahedron	4	6	4				
cube	8	12	6				
octahedron	6	12	8				
dodecahedron	20	30	12				
icosahedron	12	30	20				



Platonic solids

[Kryven et al., CALDAM'18]

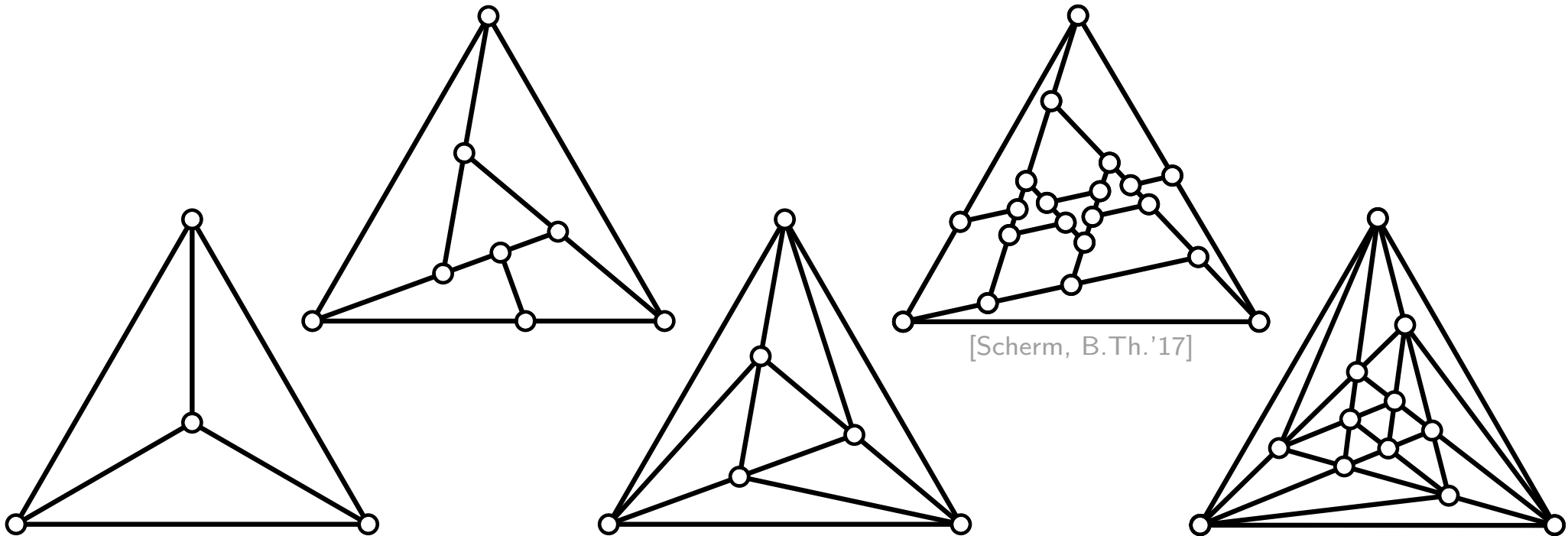
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Platonic solids

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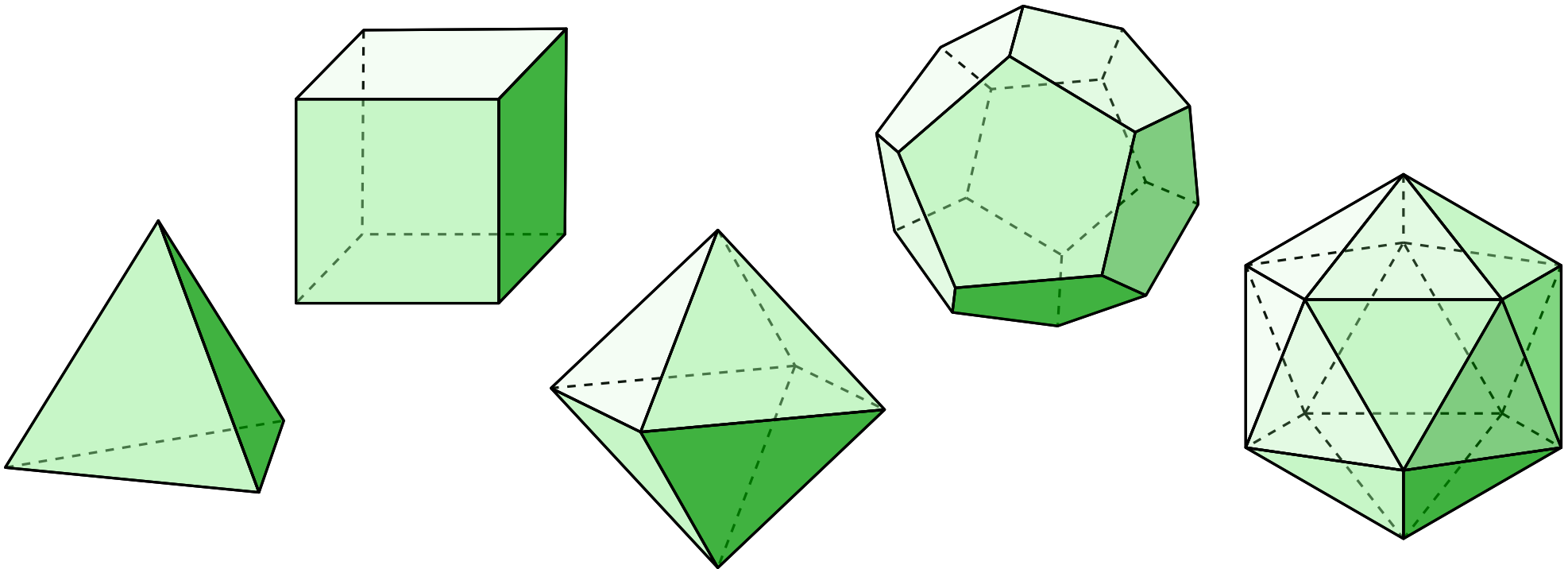
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cube	8	12	6	7	7		
octahedron	6	12	8	9	9		
dodecahedron	20	30	12	9...10	9...10		
icosahedron	12	30	20	13...15	13...15		



Platonic solids

[Kryven et al., CALDAM'18]

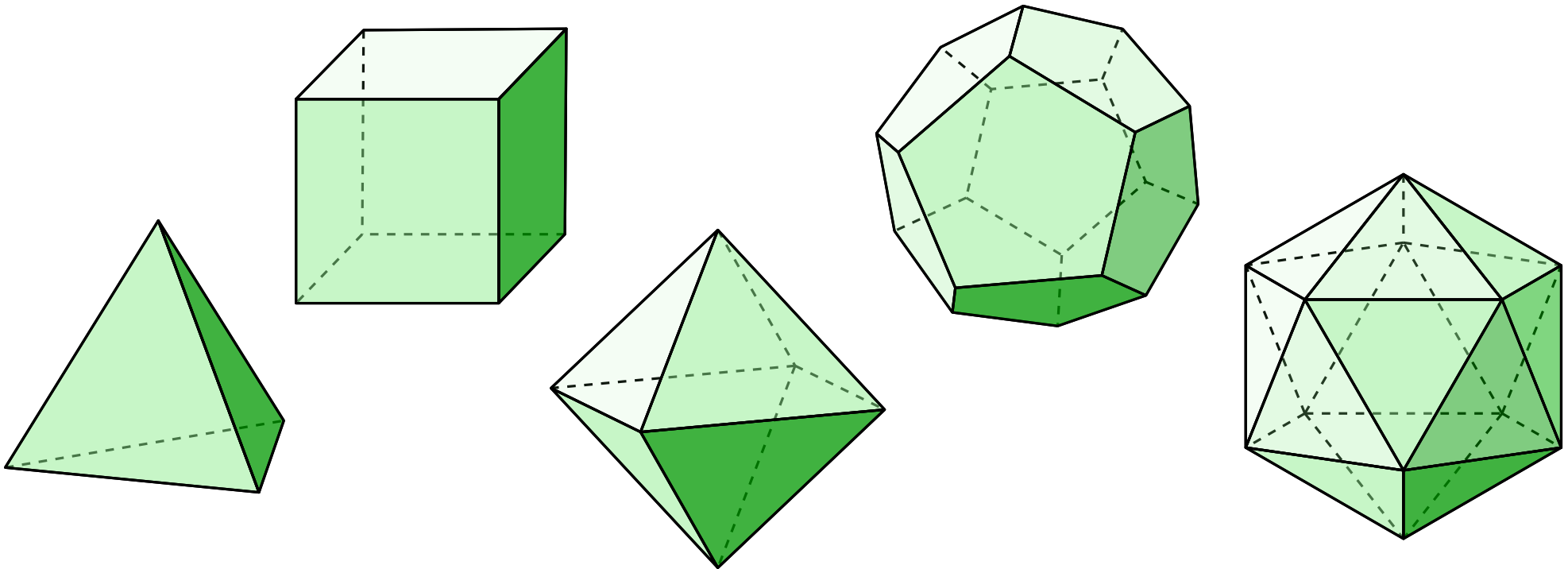
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Platonic solids

[Kryven et al., CALDAM'18]

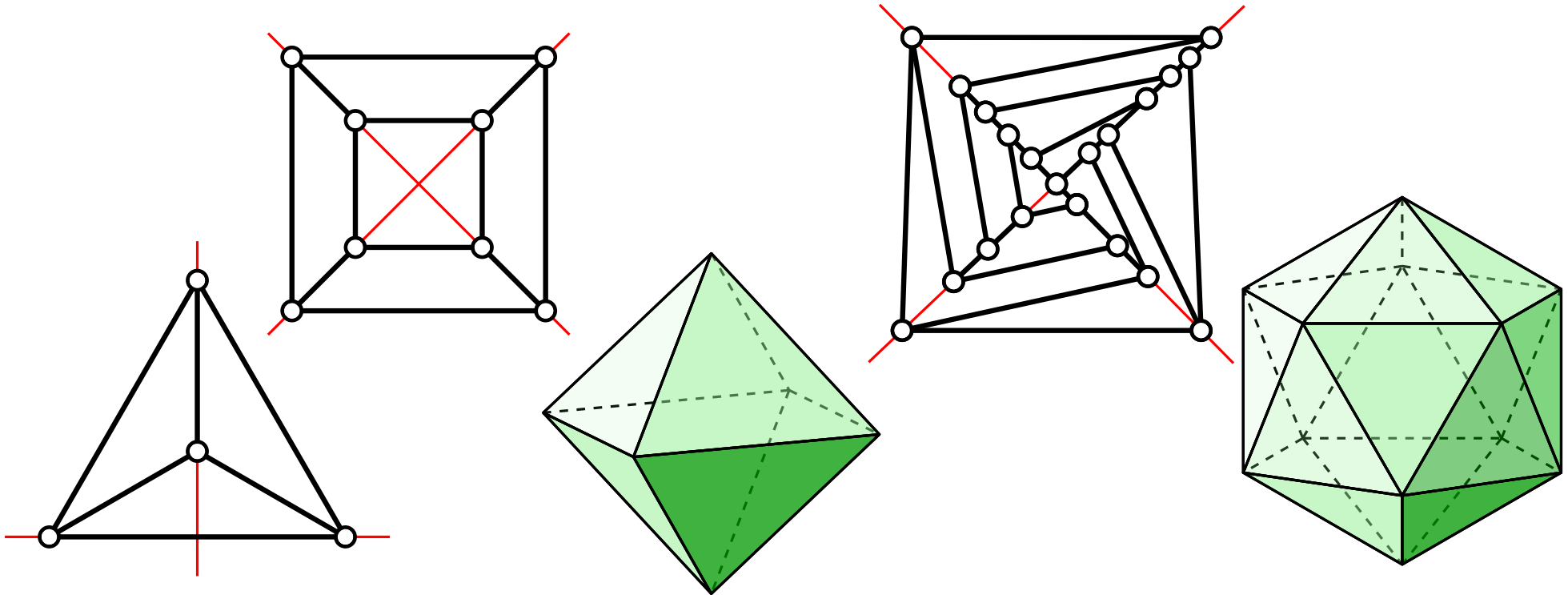
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Platonic solids

[Kryven et al., CALDAM'18]

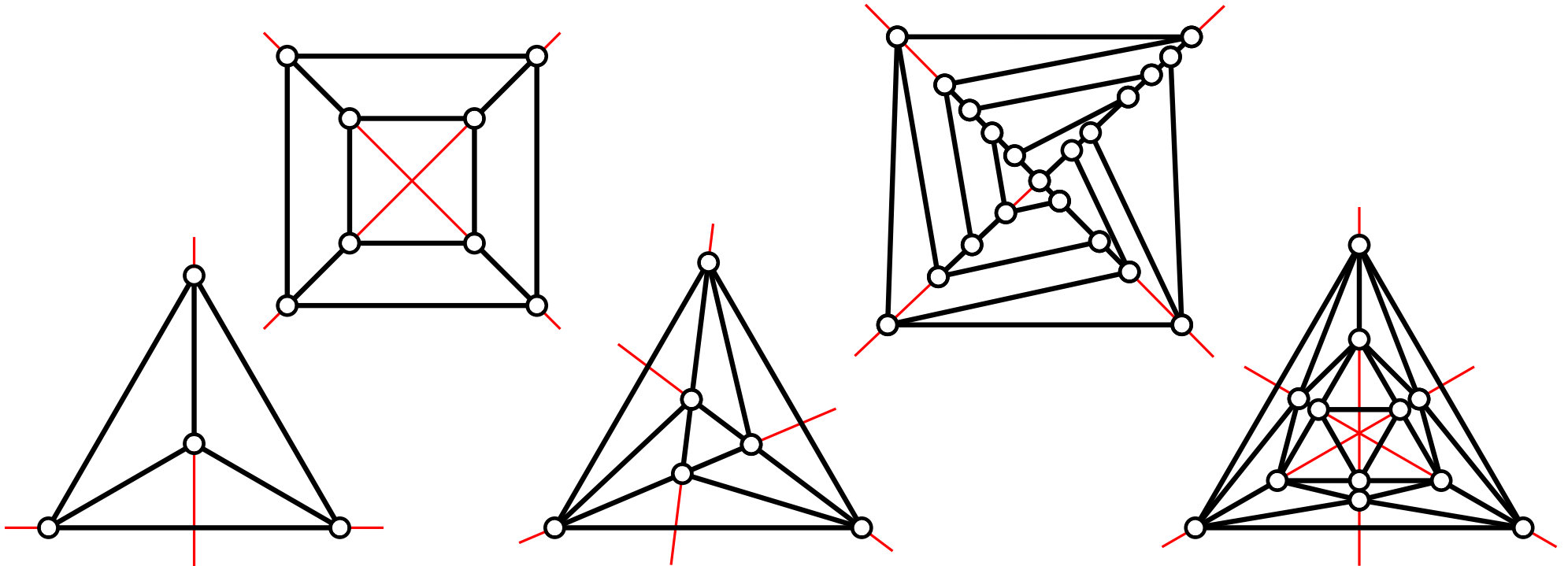
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Platonic solids

[Kryven et al., CALDAM'18]

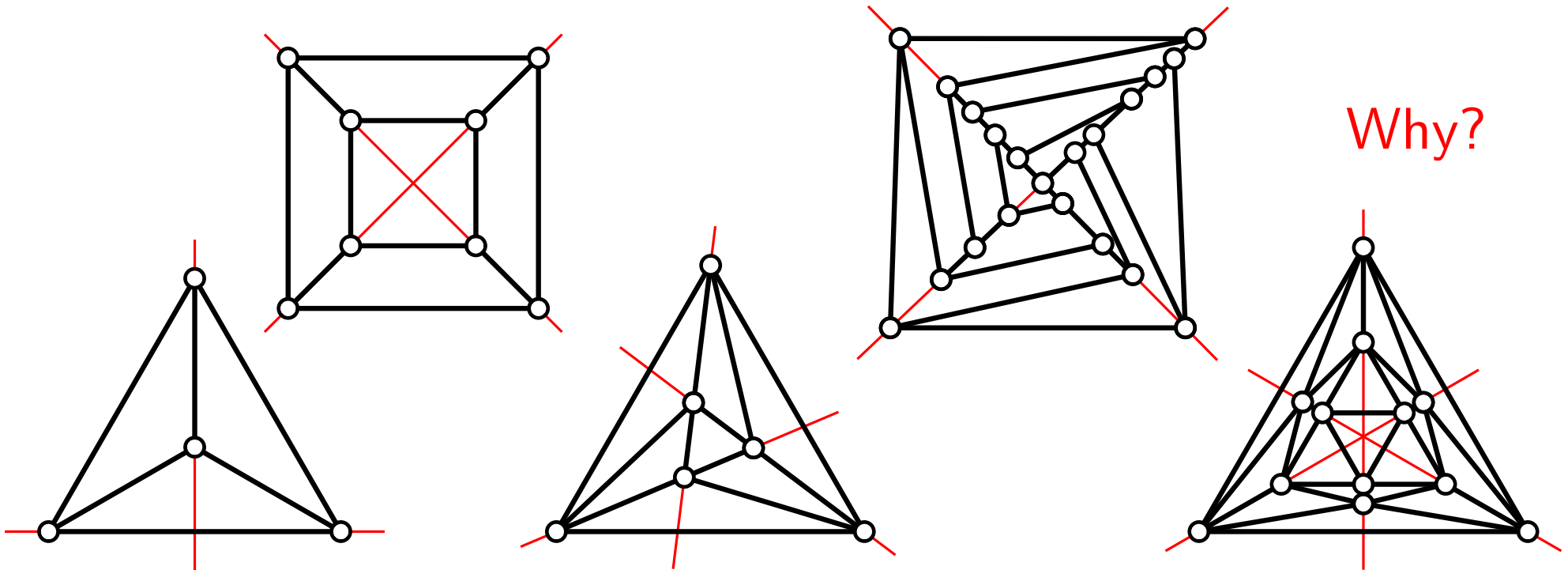
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Platonic solids

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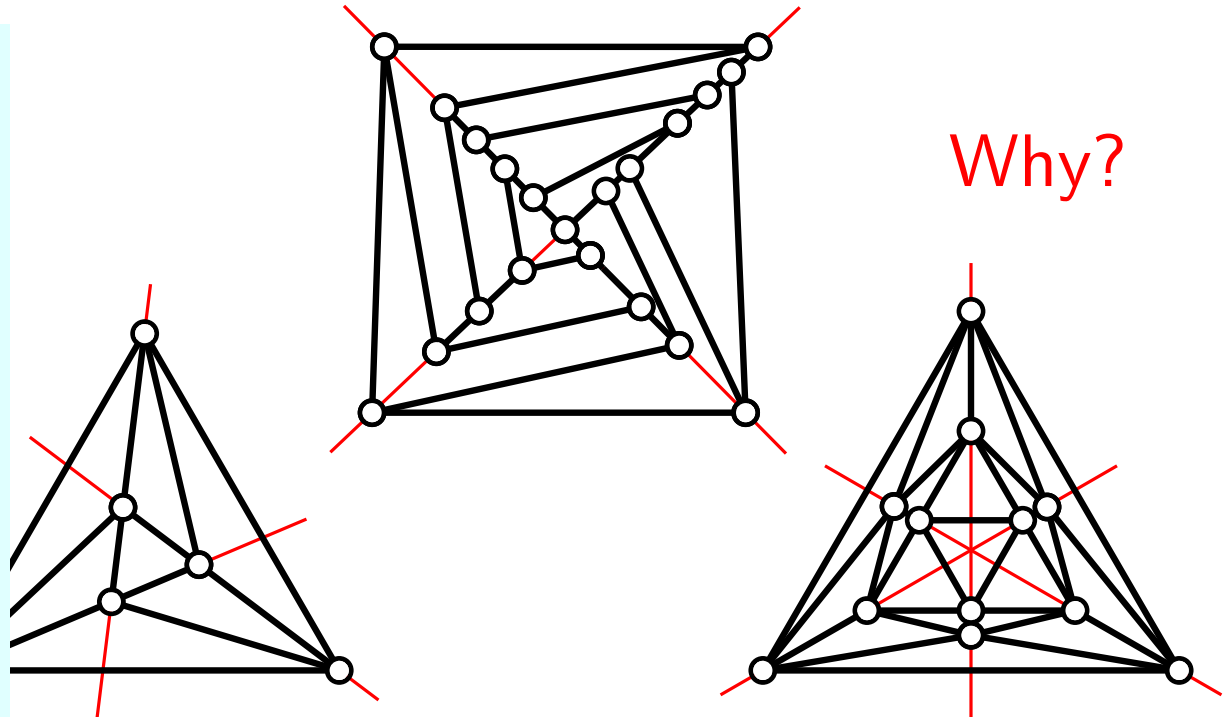


Platonic solids

[Kryven et al., CALDAM'18]

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Proof

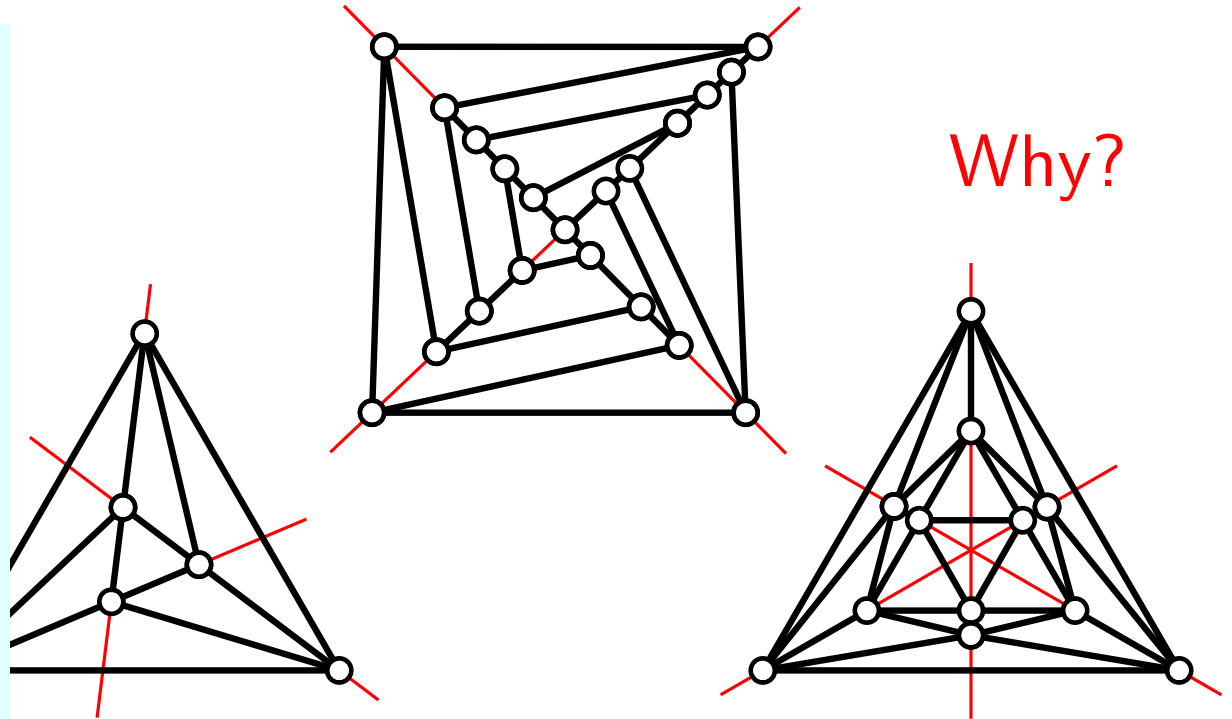
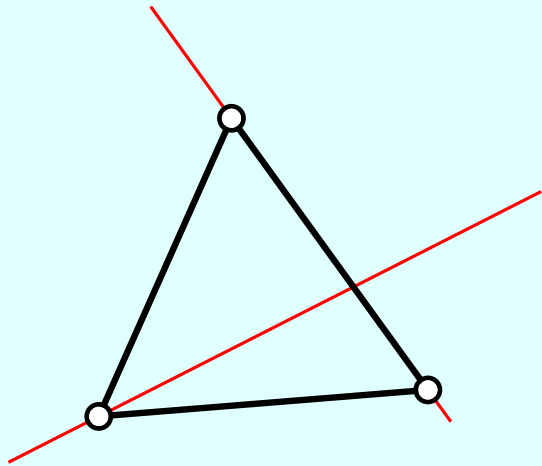


Platonic solids

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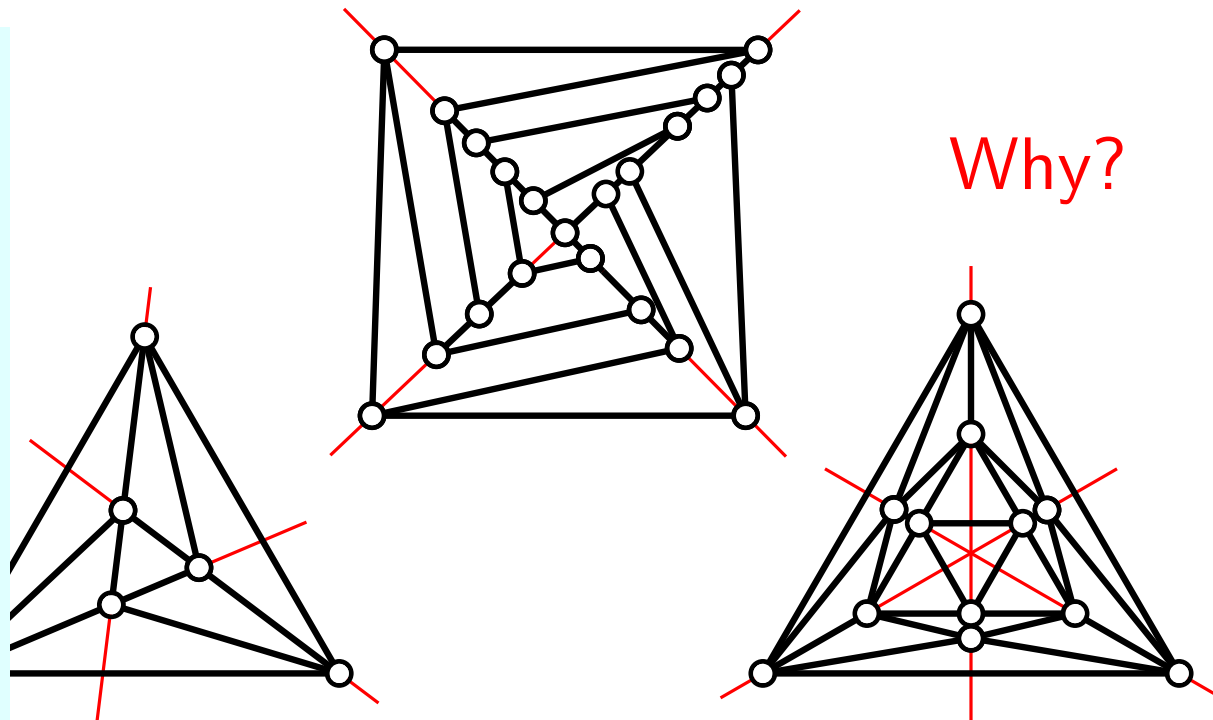
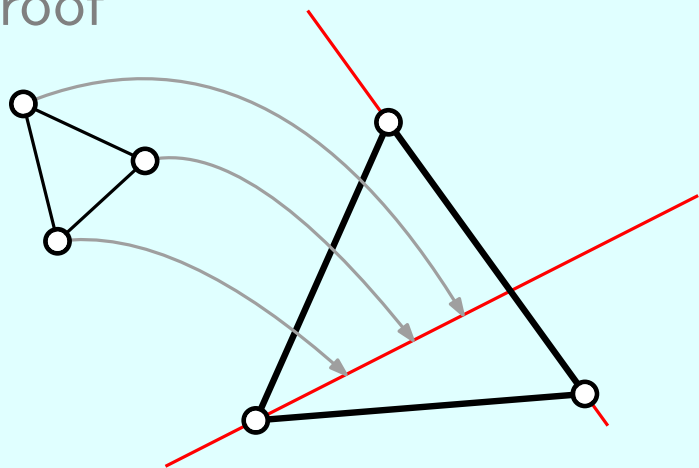


Platonic solids

[Kryven et al., CALDAM'18]

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Proof

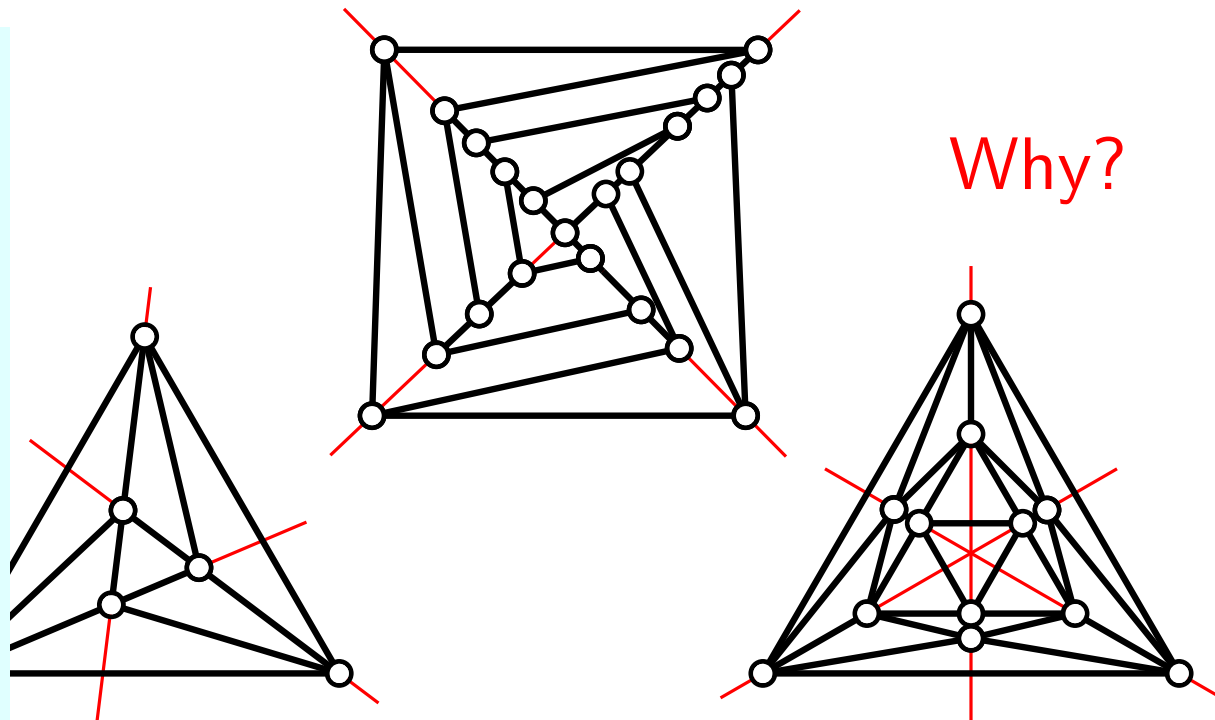
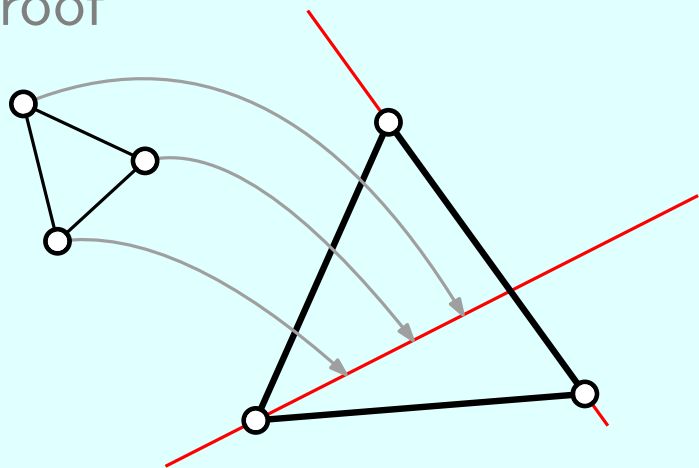


Platonic solids

[Kryven et al., CALDAM'18]

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Proof



Platonic solids

[Kryven et al., CALDAM'18]

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Platonic solids

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Platonic solids

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$$\pi_3^1(G) = \text{lva}(G)$$

[Chaplick et al., GD'16]

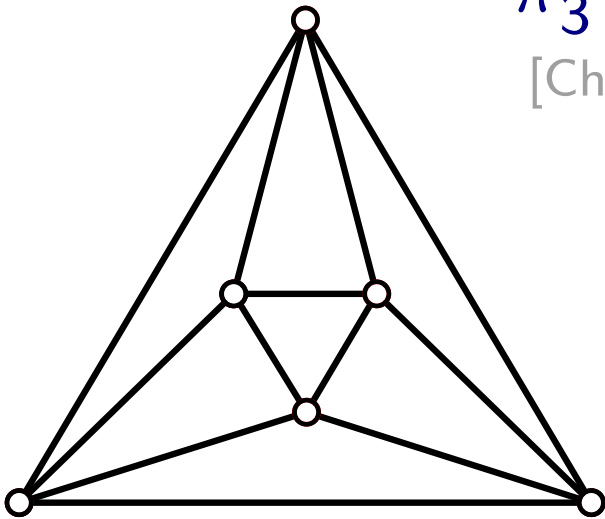
Platonic solids

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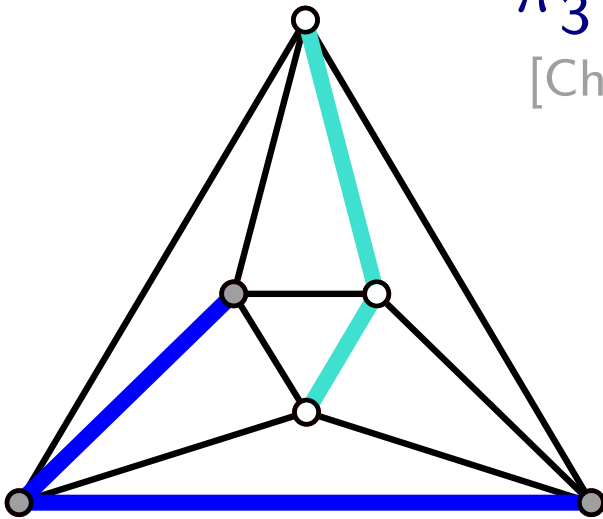
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$$\pi_3^1(G) = \text{lva}(G)$$

[Chaplick et al., GD'16]



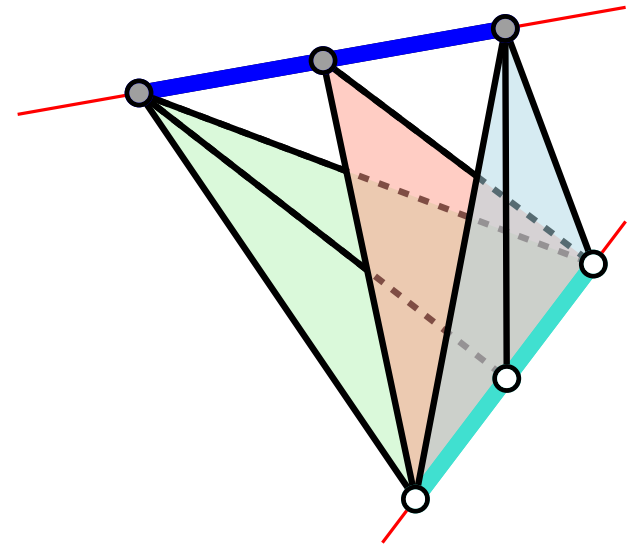
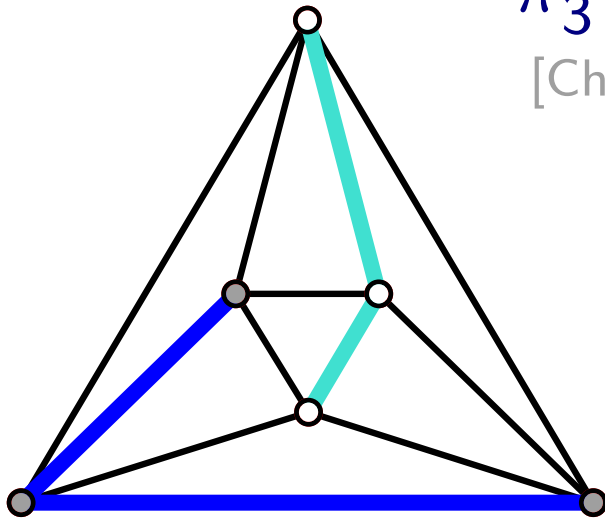
Platonic solids

[Kryven et al., CALDAM'18]

$G = (V, E)$	$ V $	$ E $	$ F $	$\rho_2^1(G)$	$\rho_3^1(G)$	$\pi_2^1(G)$	$\pi_3^1(G)$
tetrahedron	4	6	4	6	6	2	2
cube	8	12	6	7	7	2	2
octahedron	6	12	8	9	9	3	2
dodecahedron	20	30	12	9...10	9...10	2	2
icosahedron	12	30	20	13...15	13...15	3	

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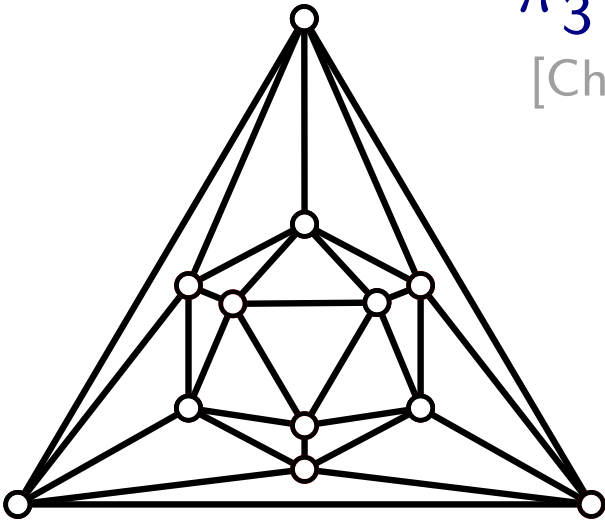
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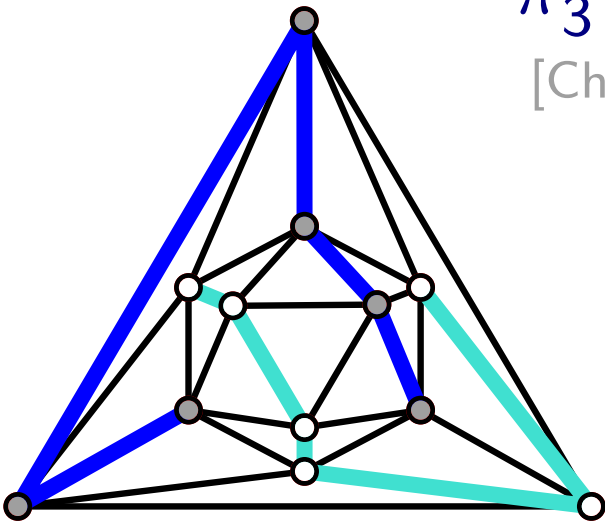
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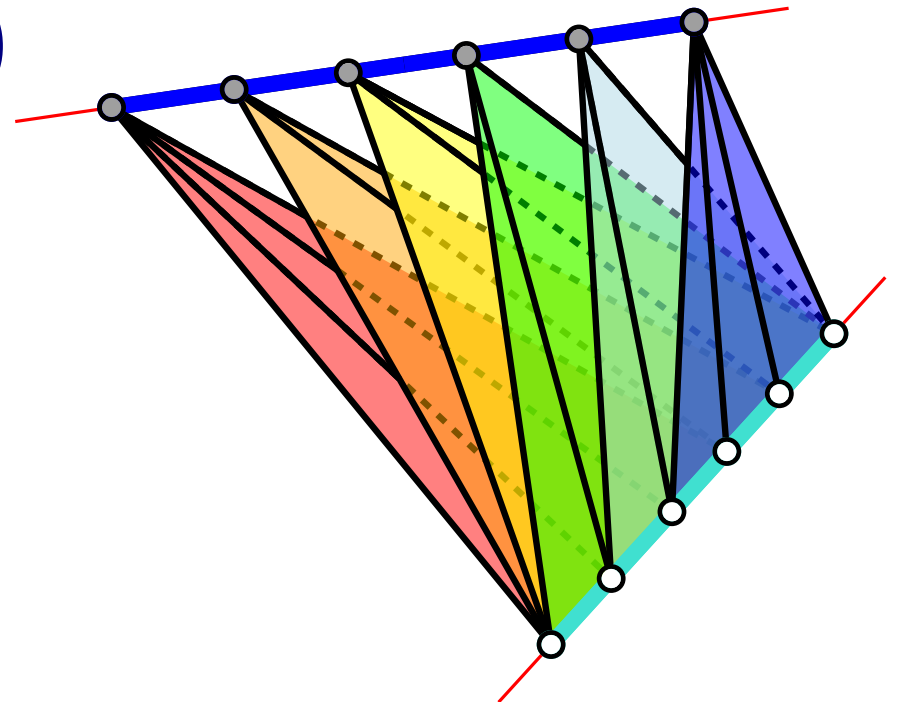
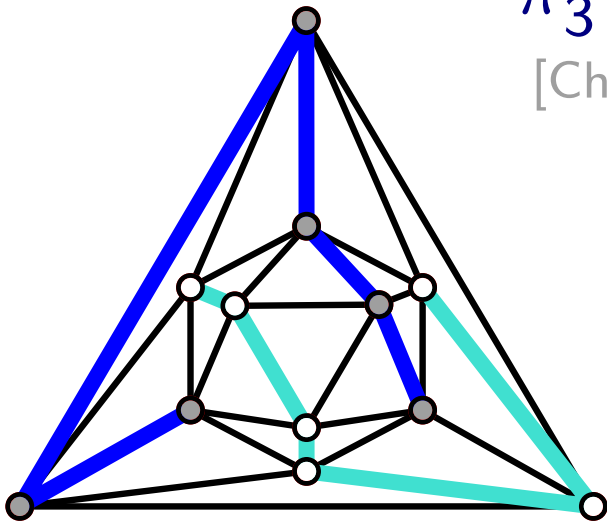
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Motivation

$$\pi \frac{1}{2}$$

How large can it be?

Motivation

π_2^1

How large can it be?

Q: Is the π_2^1 -value unbounded for some graph families?

Motivation

π_2^1

How large can it be?

Q: Is the π_2^1 -value unbounded for some graph families?

Yes!

Motivation

π_2^1

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[Ravsky and Verbitsky, WG'11]
[Da Lozzo et al., GD'16]

Yes!

- $\text{tw}(G_i) = 5$

Motivation

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- $\text{maxdeg}(G_i) \leq 12$
 $\pi_2^1(G_i) \geq n^{0.01}$

[Chaplick et al., GD'16]

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π_2^1

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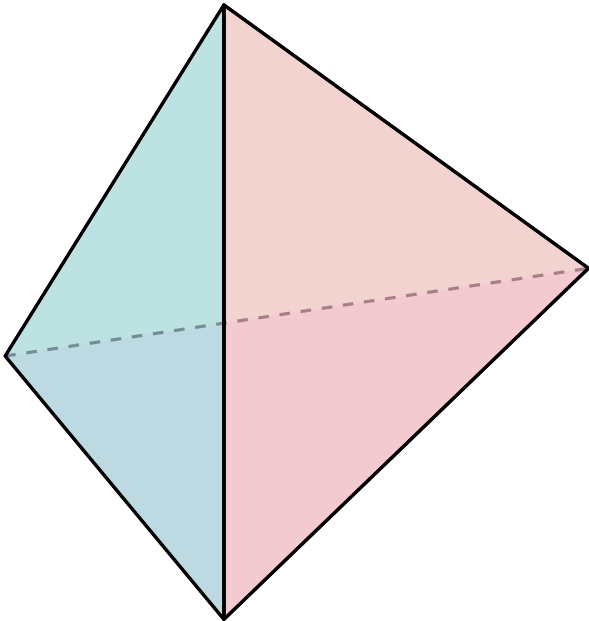
- $\text{tw}(G_i) = 3$
 $\text{maxdeg}(G_i) = 6$
 $\pi_2^1(G_i) \in \Omega(\log n_i)$

New!

Yes!

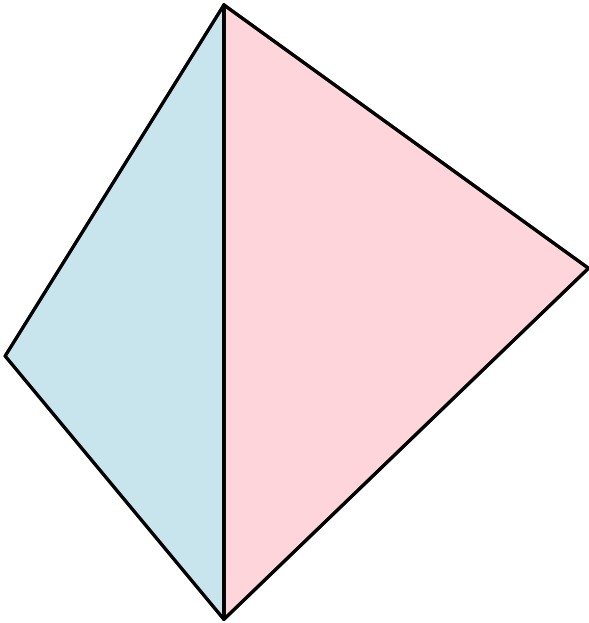
Main contribution

We construct an infinite family of graphs



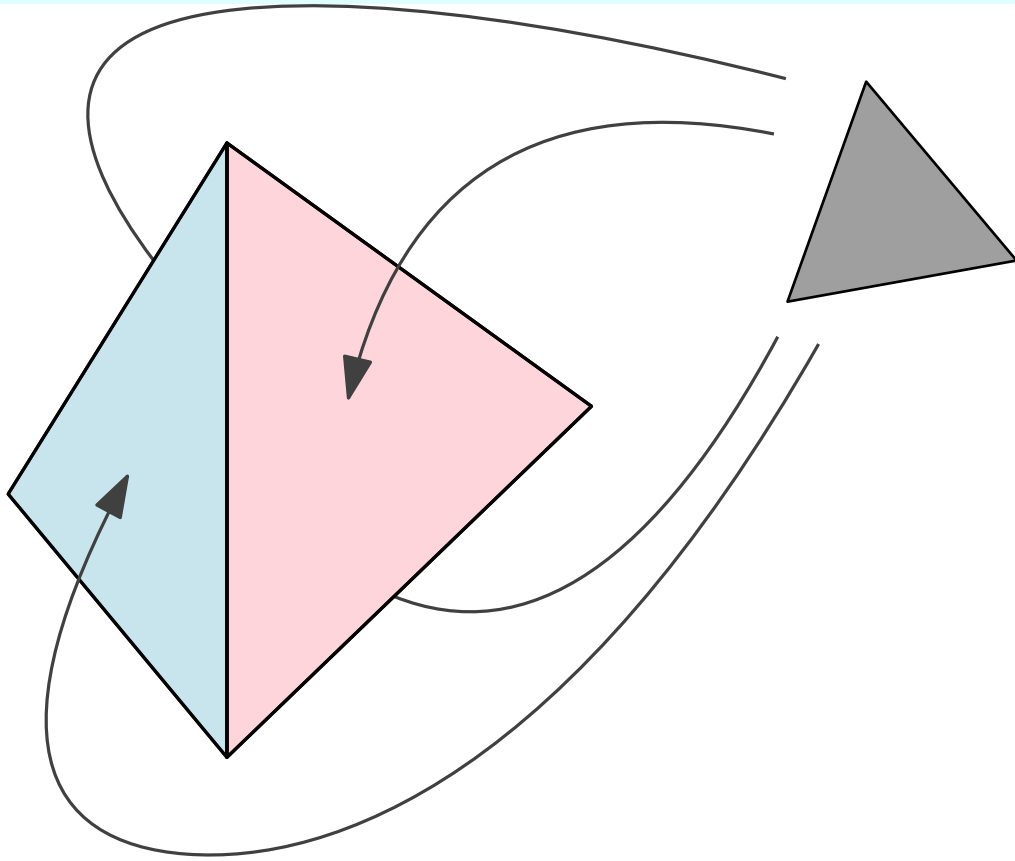
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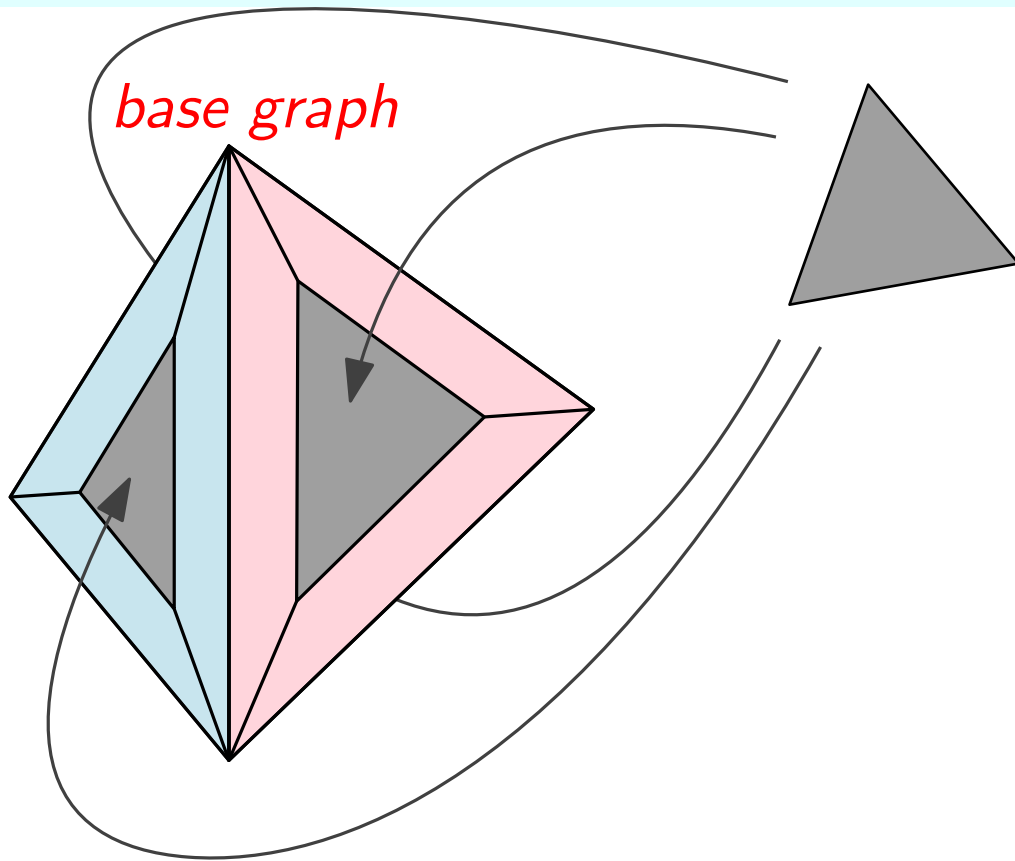
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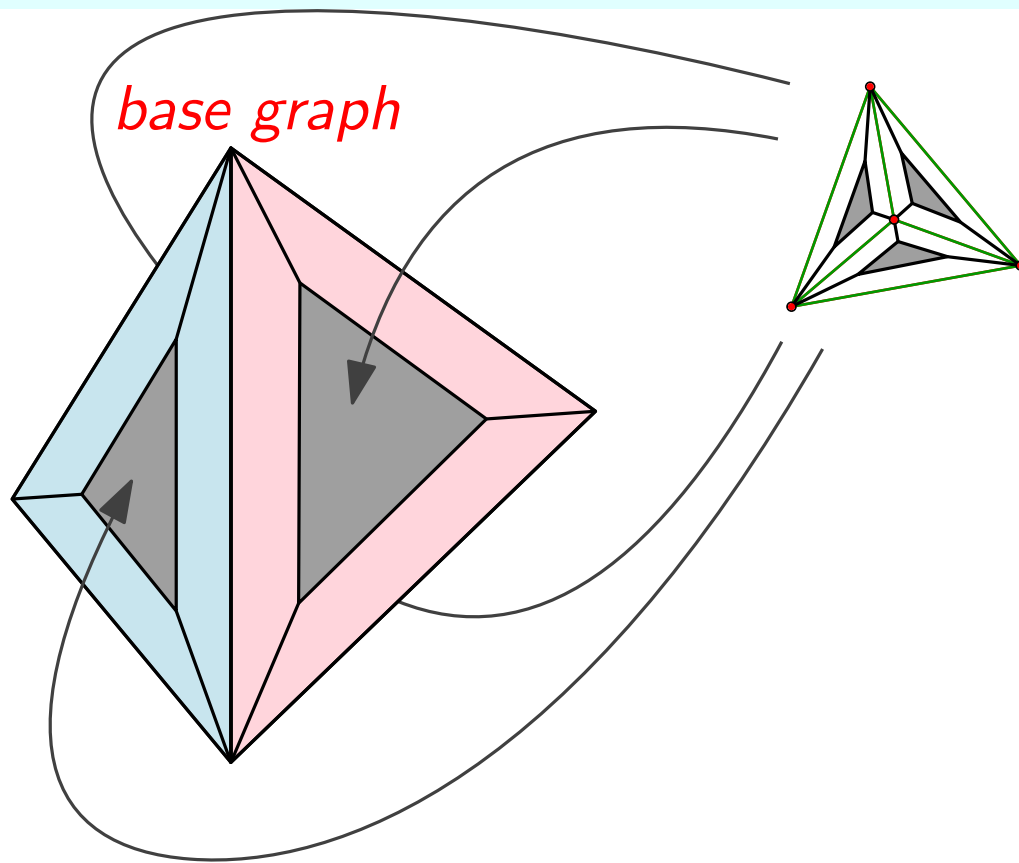
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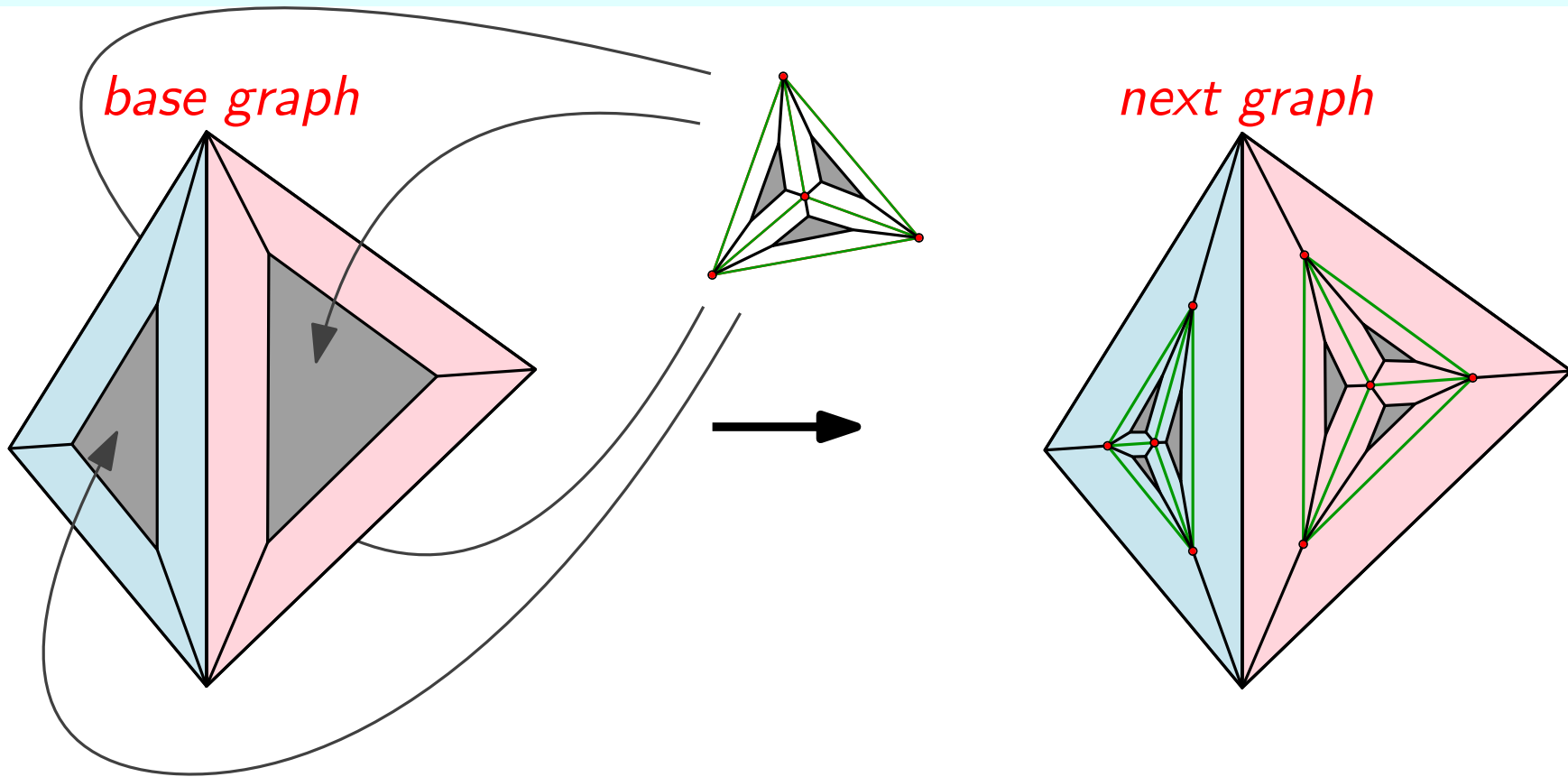
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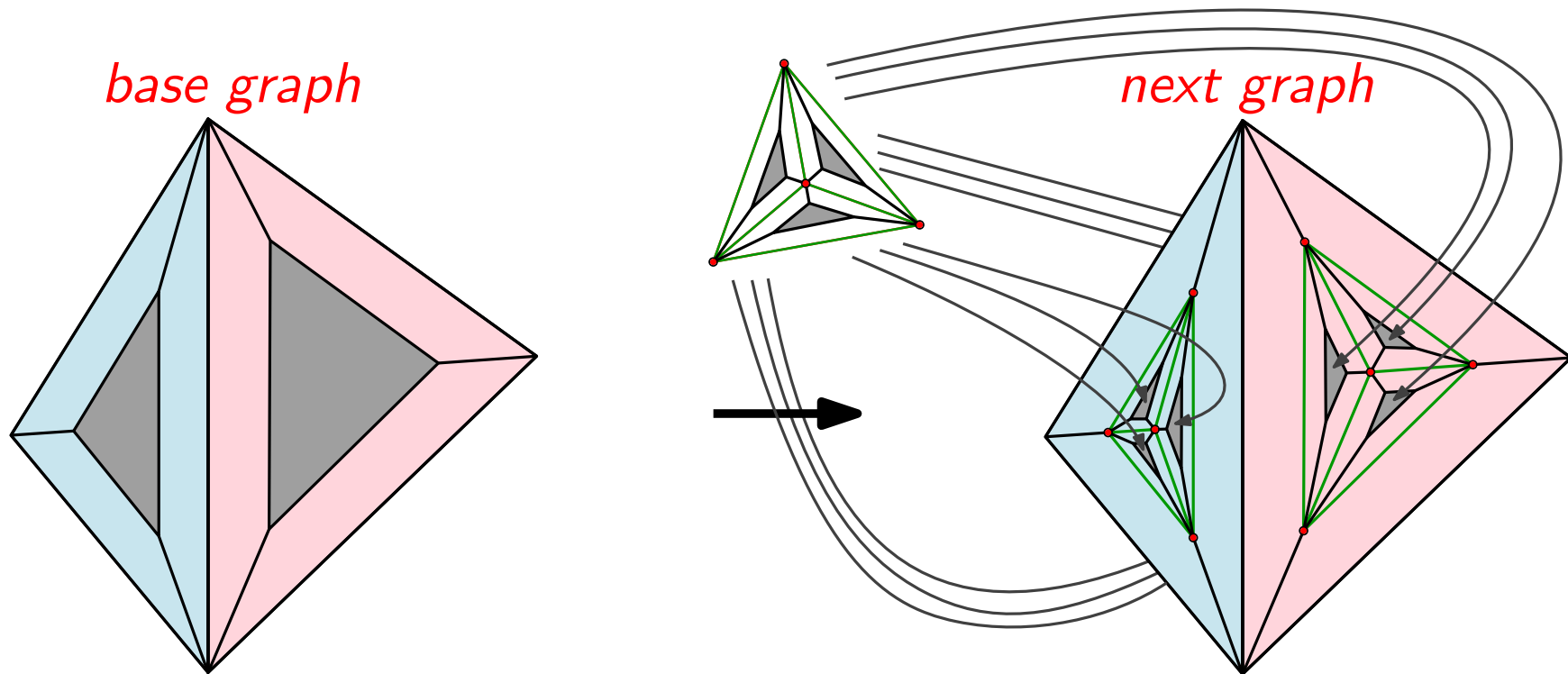
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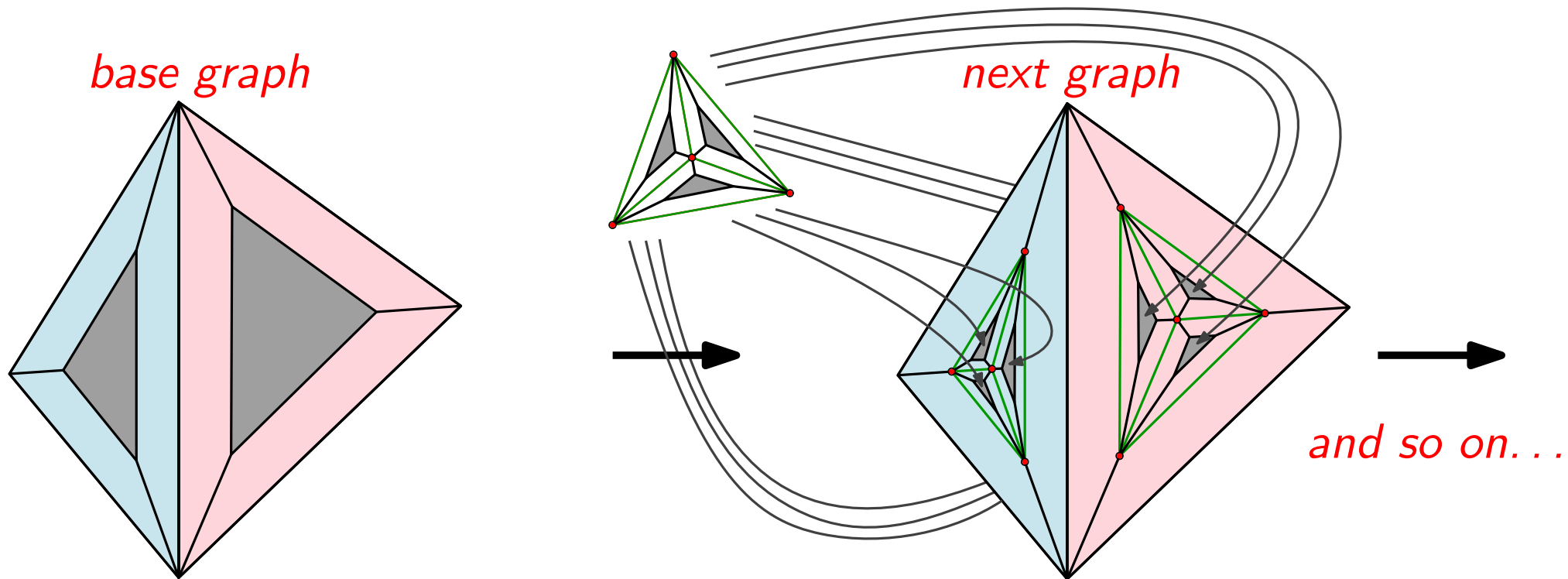
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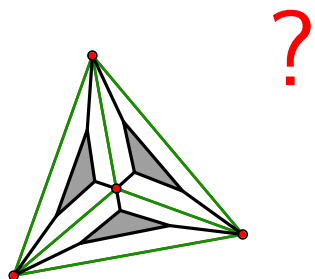


Main contribution

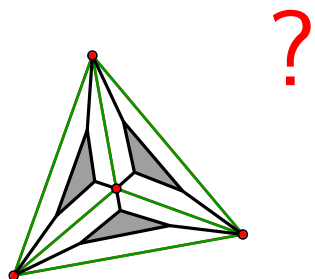
We construct an infinite family of graphs



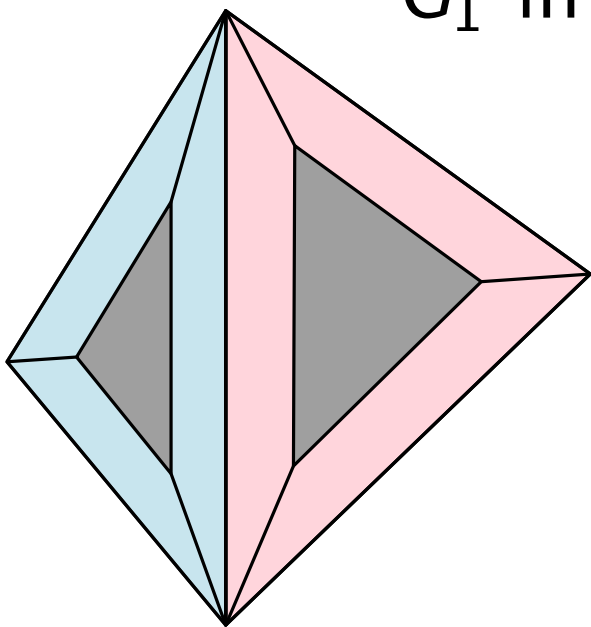
Why do we use this graph?



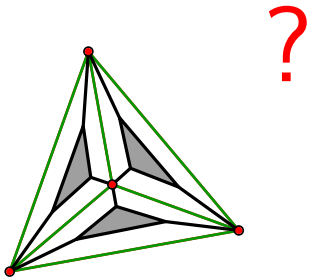
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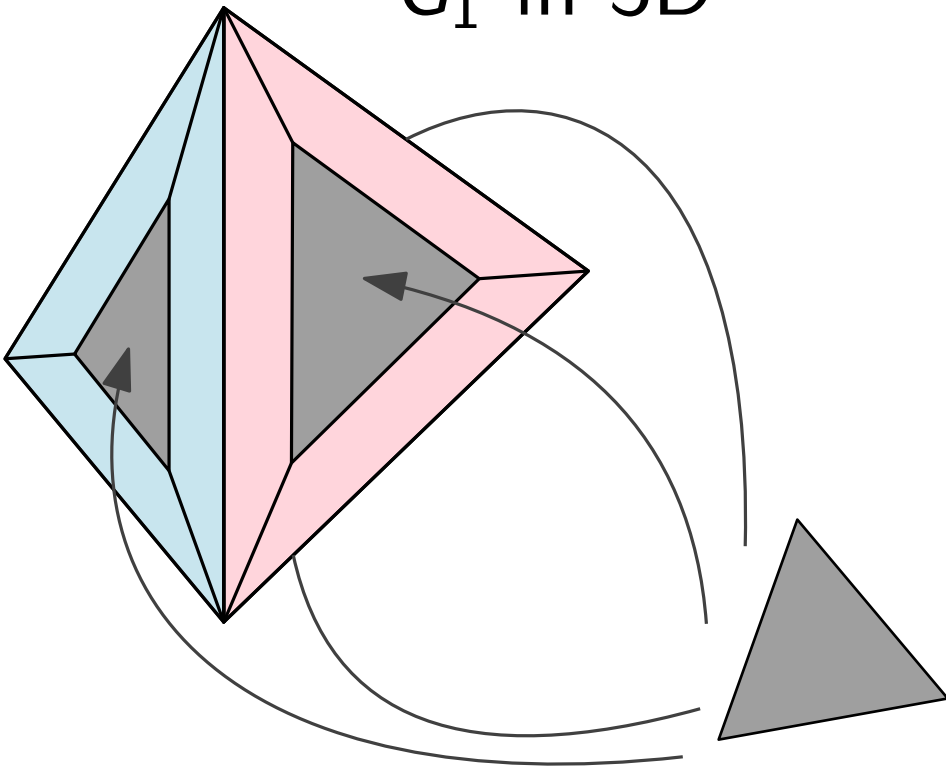
G_1 in 3D



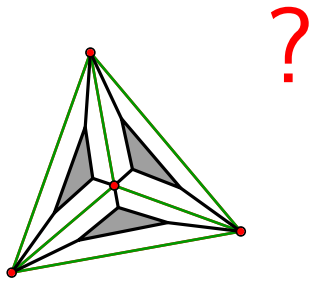
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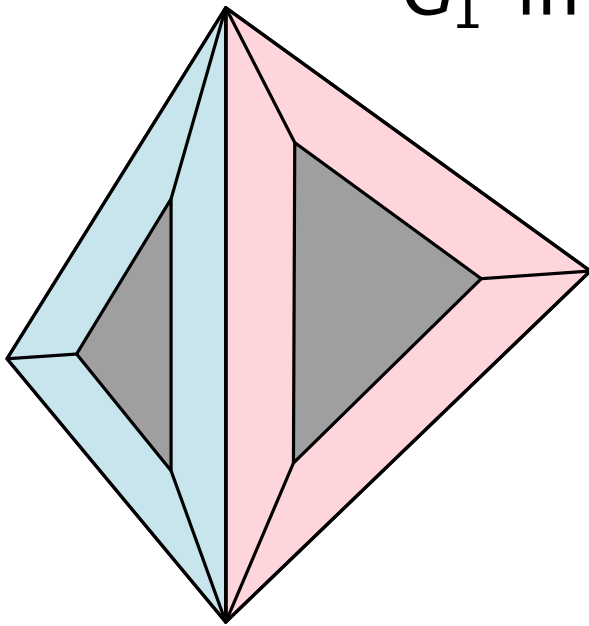
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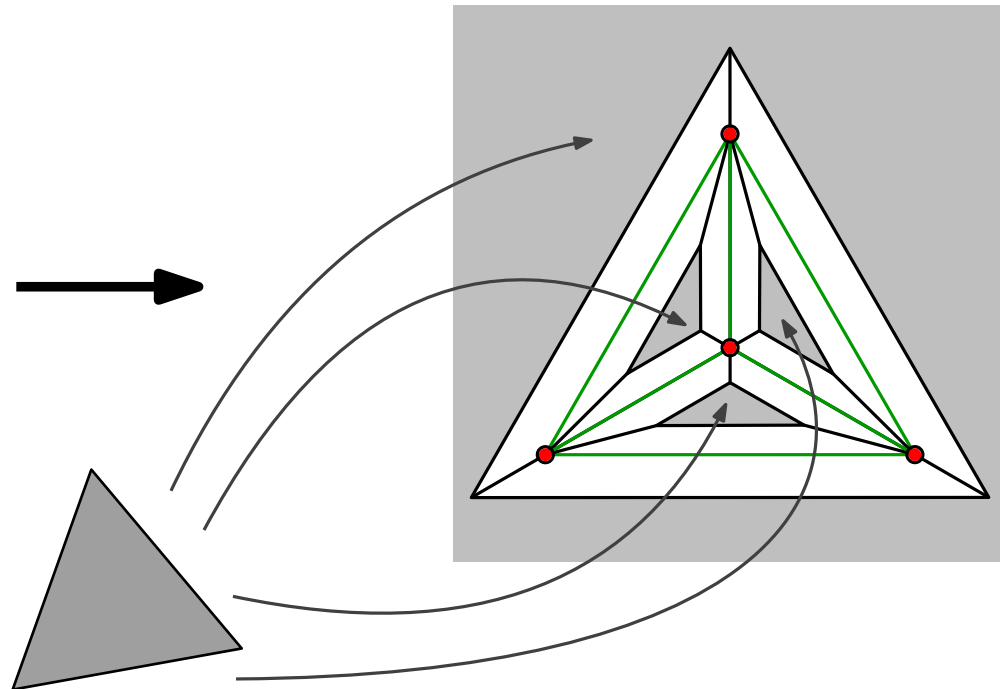
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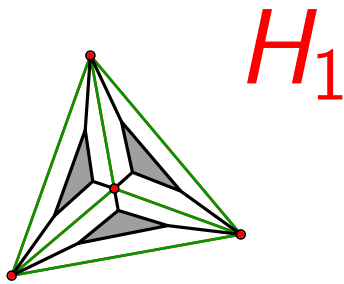
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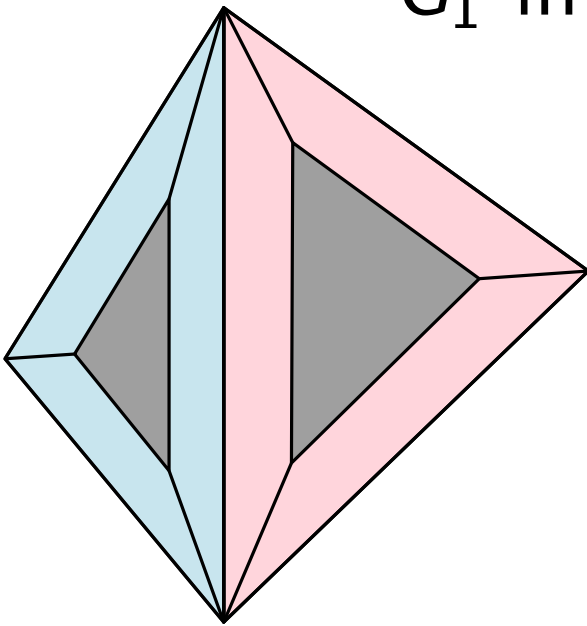
G_1 in 2D



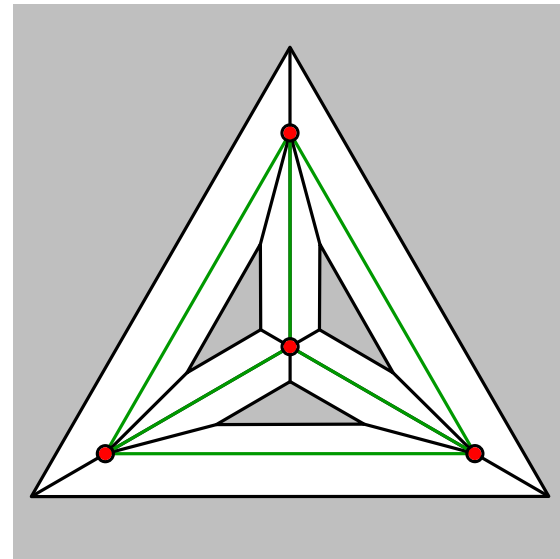
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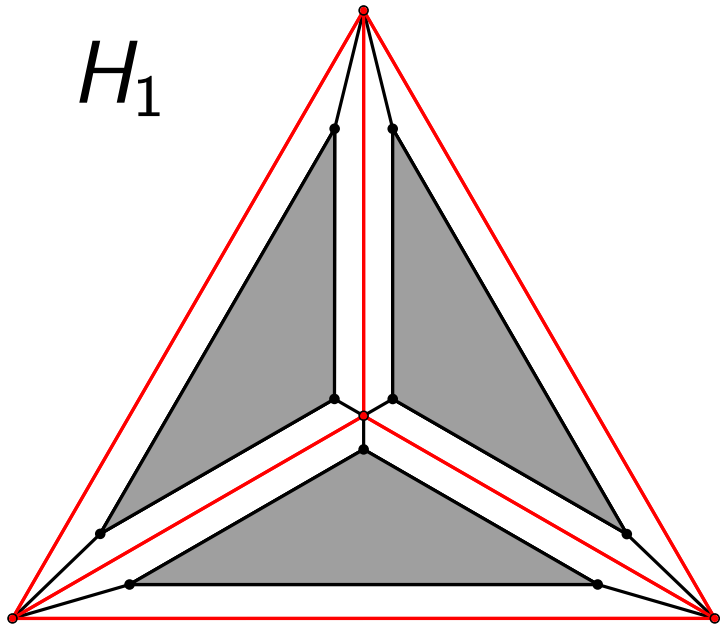
G_1 in 3D



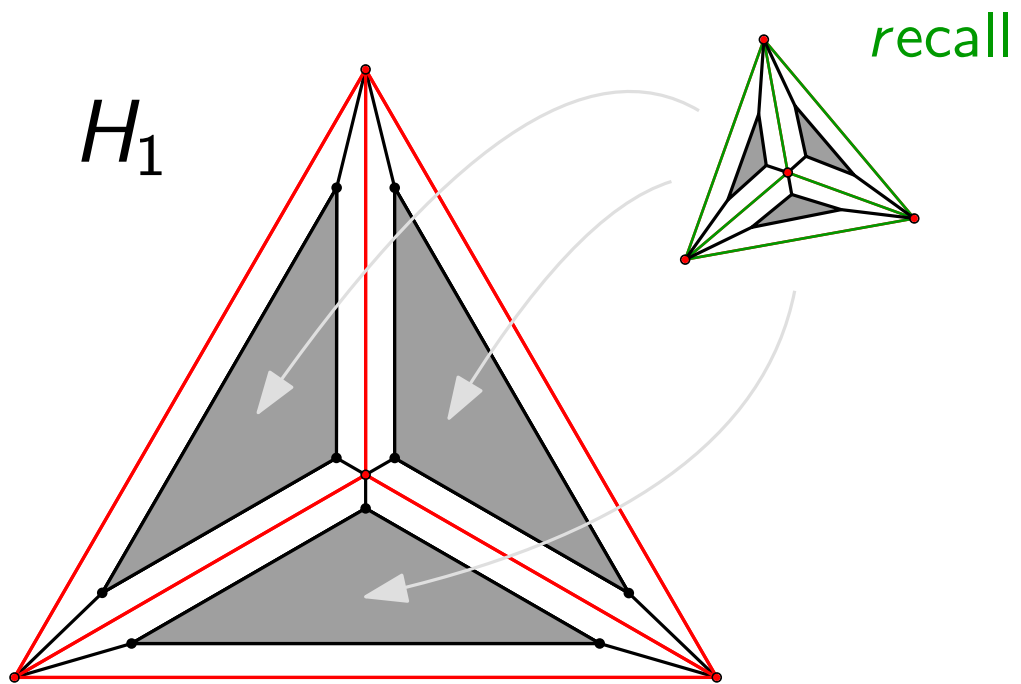
G_1 in 2D



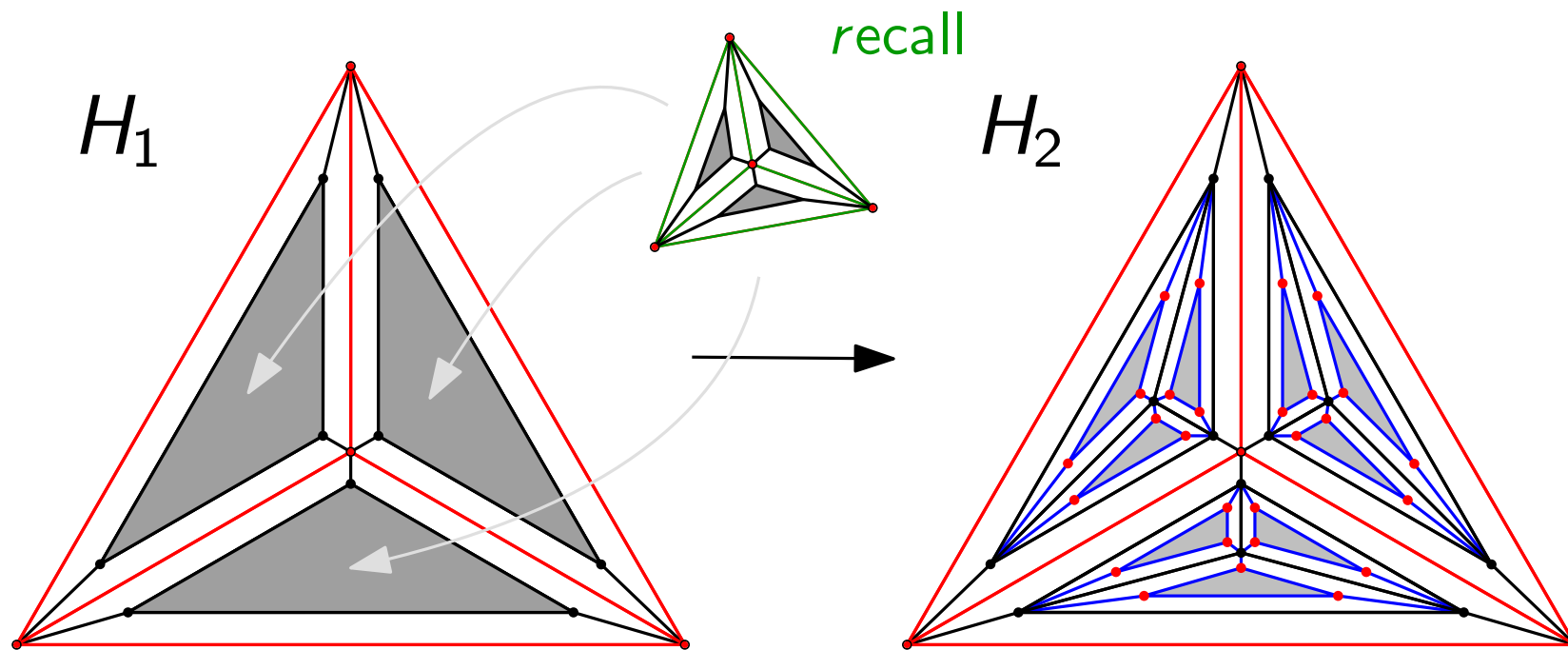
Properties of the family of graphs



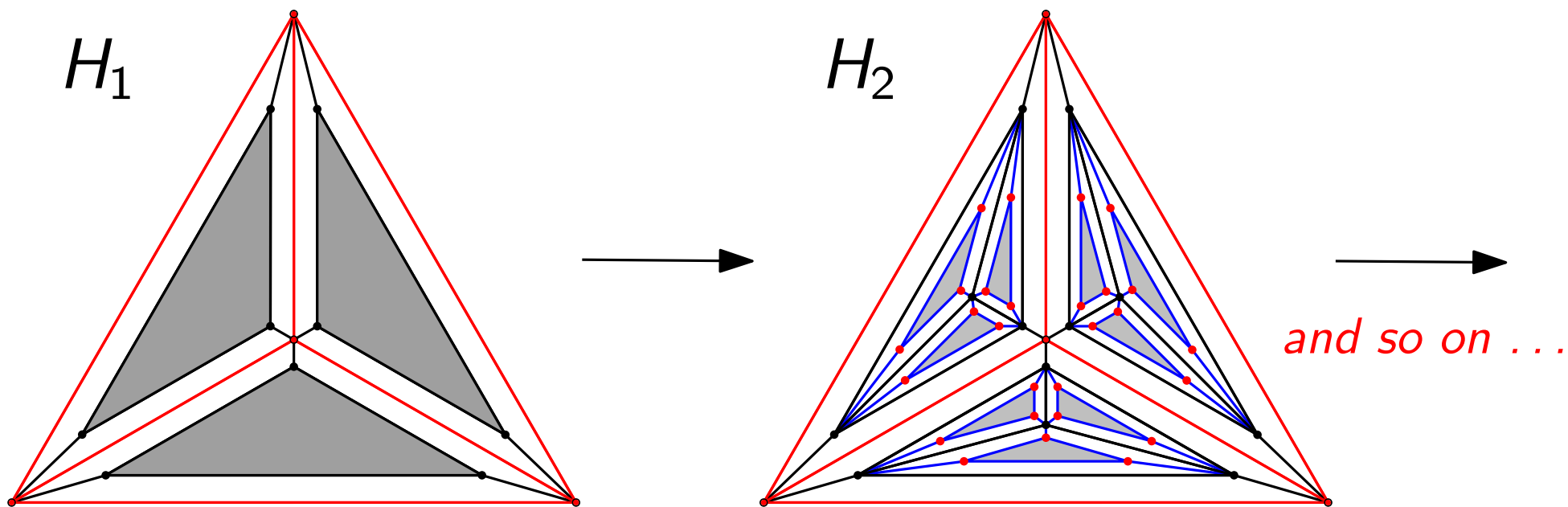
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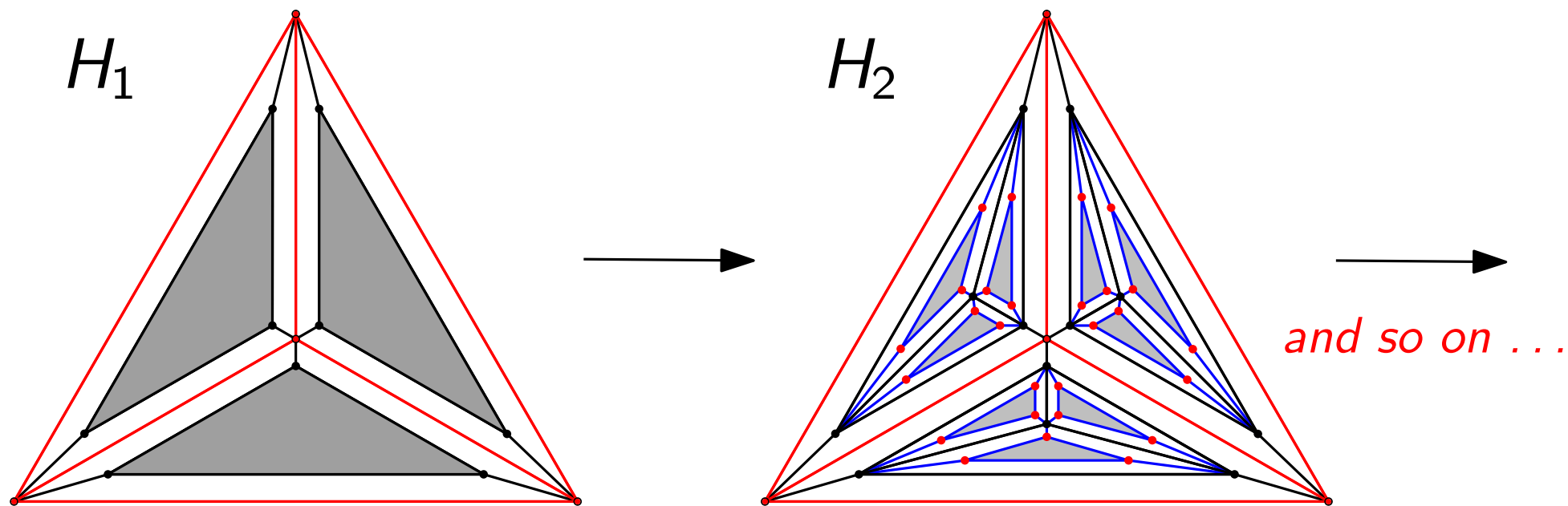
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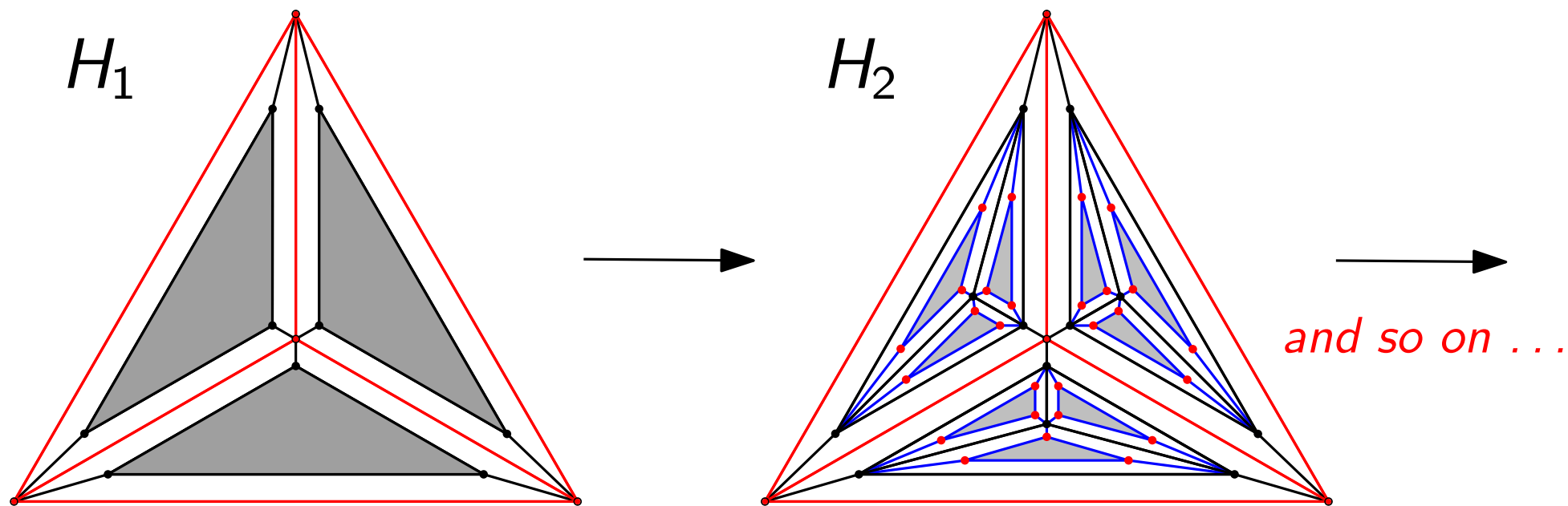
Properties of the family of graphs



maximum degree

6

Properties of the family of graphs



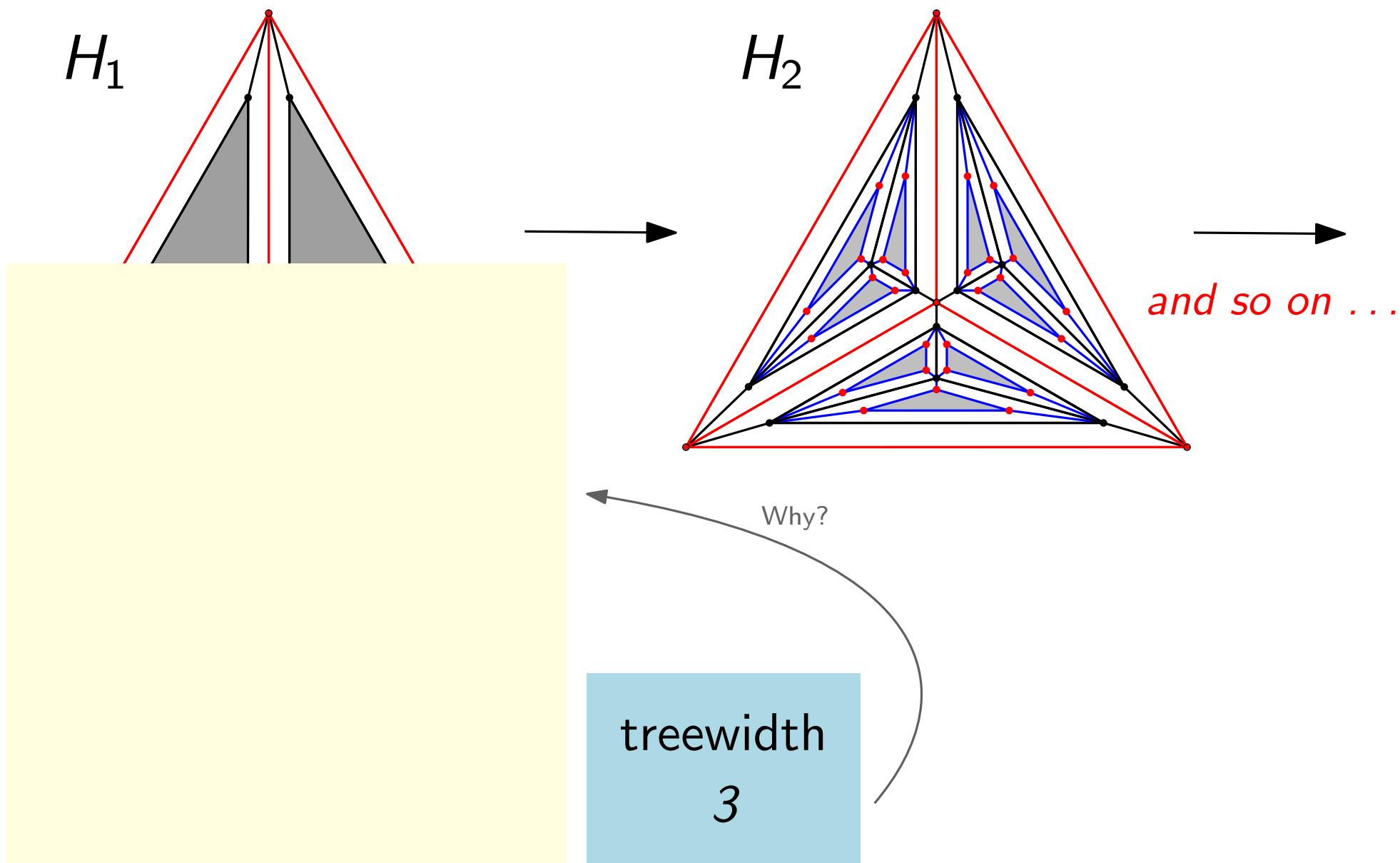
maximum degree

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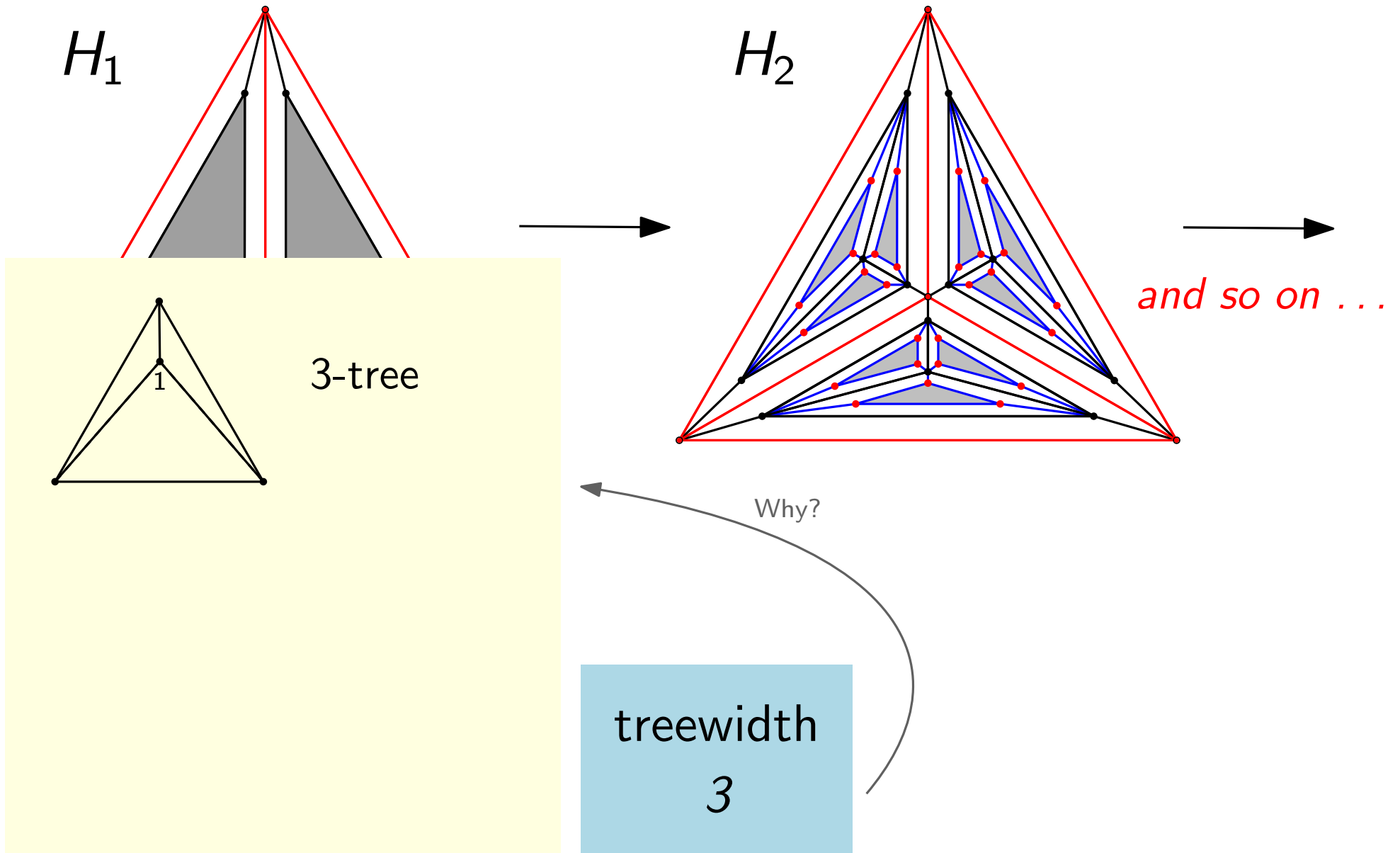
treewidth

3

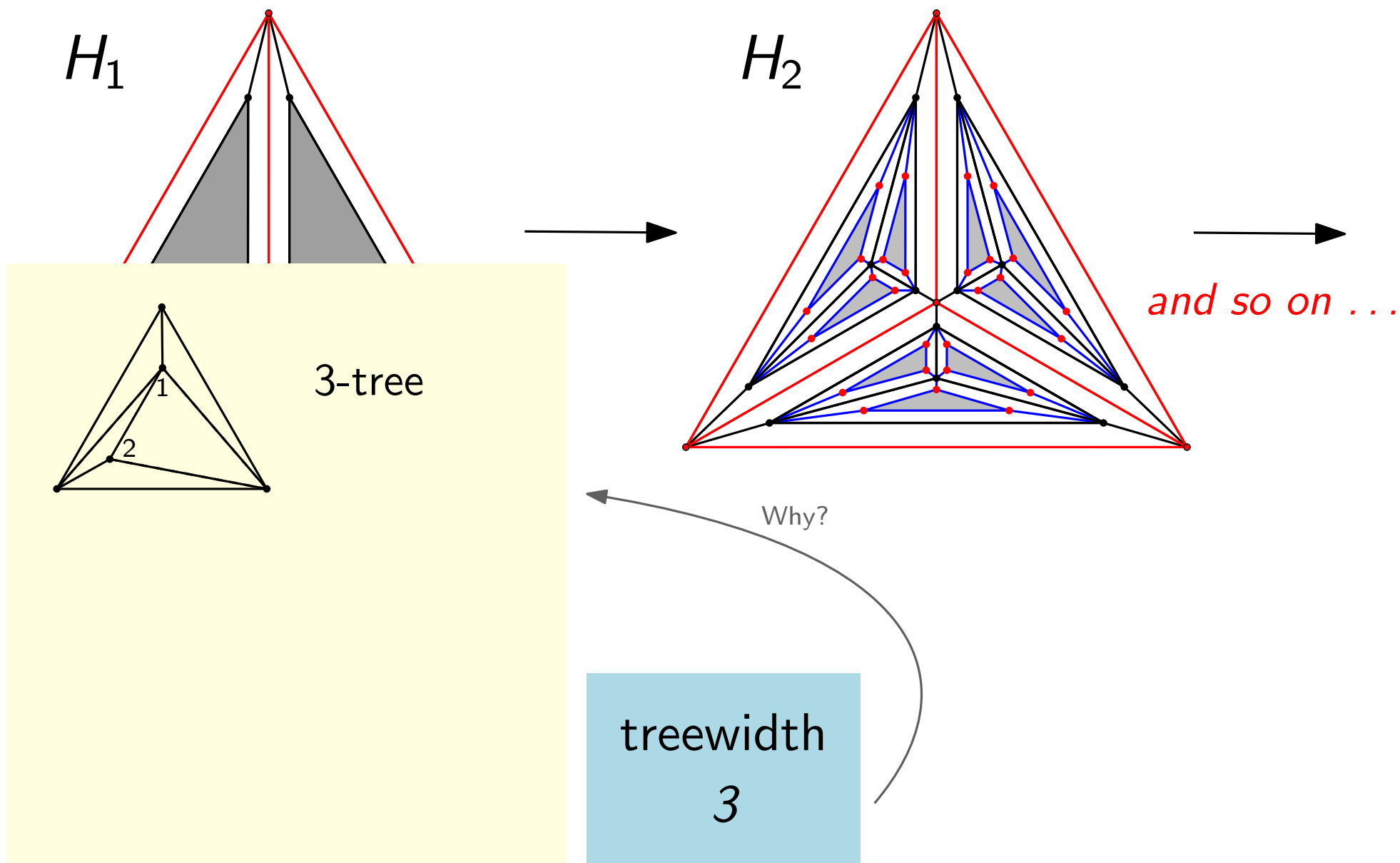
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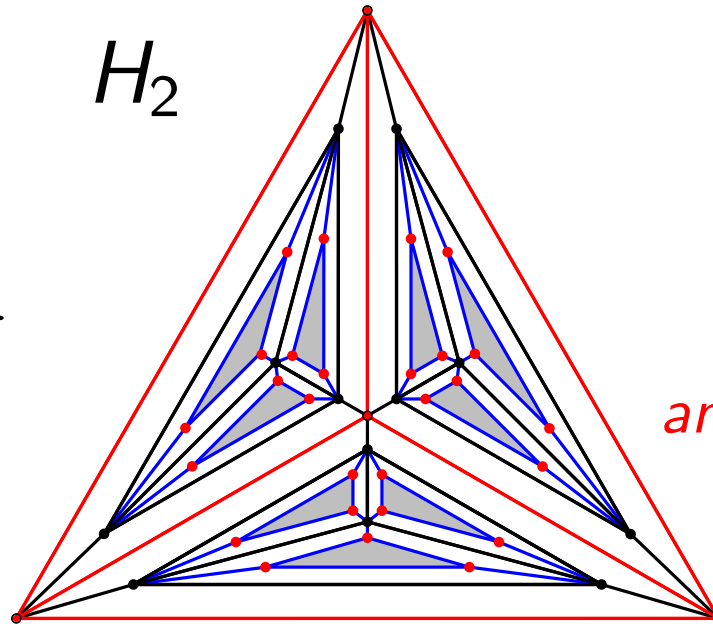
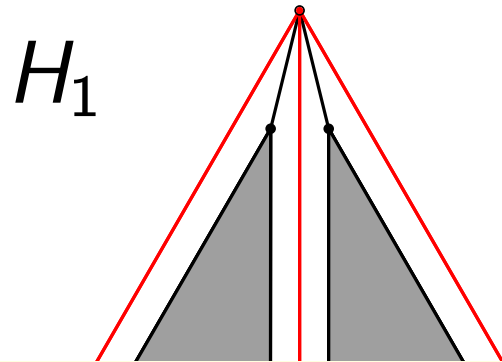
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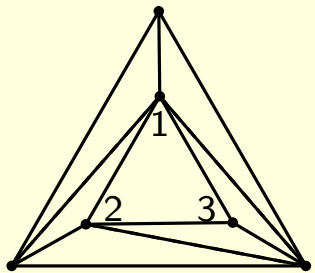
Properties of the family of graphs



Properties of the family of graphs



and so on ...



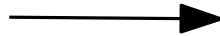
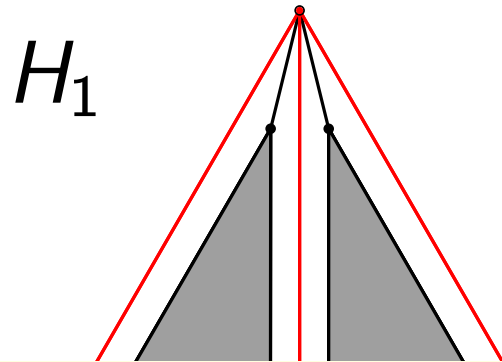
3-tree

treewidth

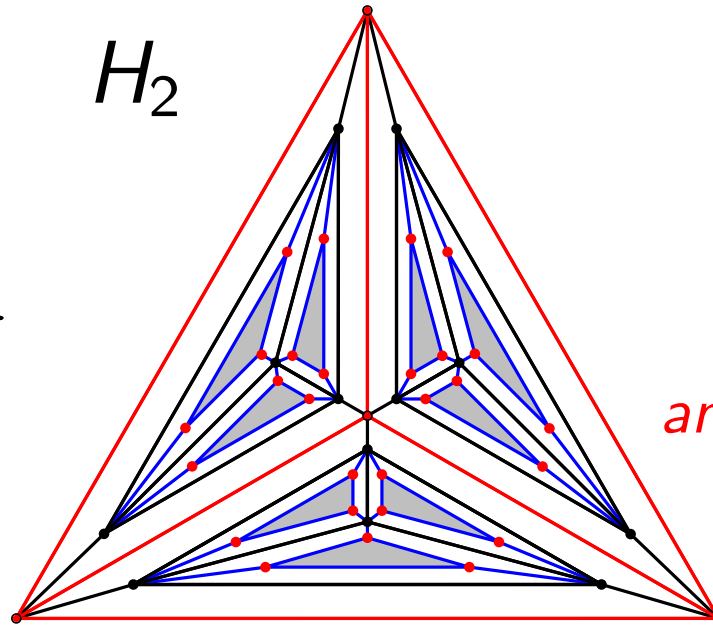
3

Why?

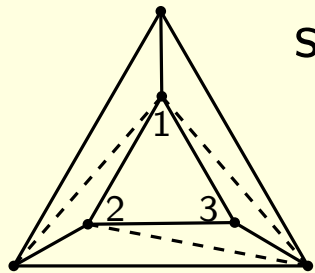
Properties of the family of graphs



H_2



and so on ...



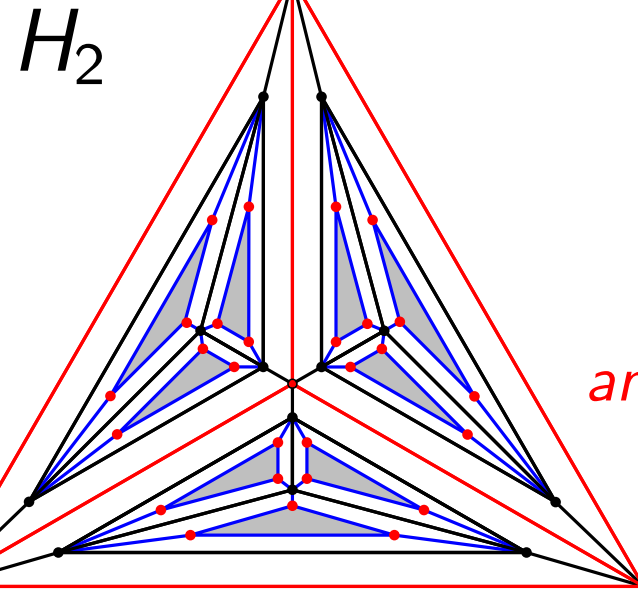
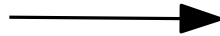
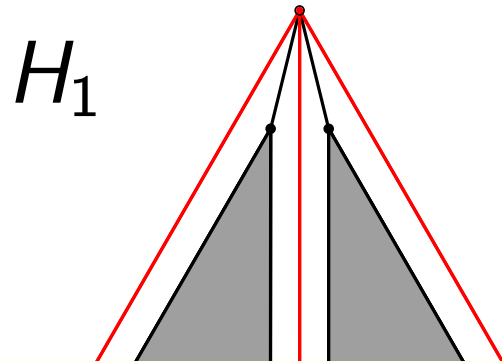
subgraph of
3-tree

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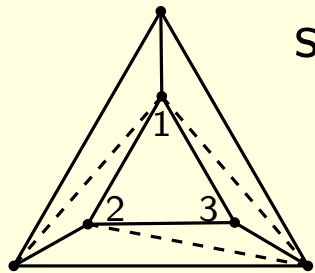
treewidth

3

Properties of the family of graphs



and so on ...



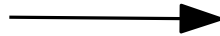
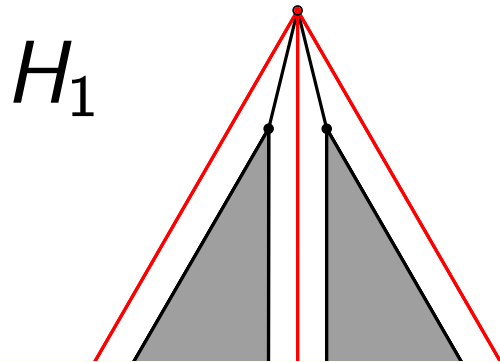
subgraph of
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$$\text{tw}(G_i) \leq 3$$

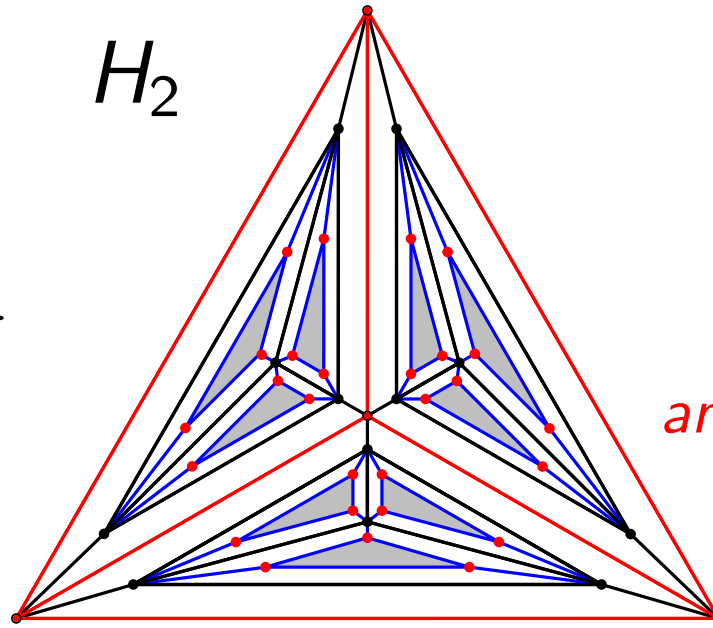
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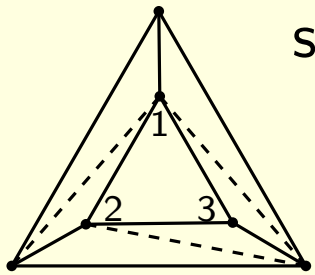
Properties of the family of graphs



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and so on ...



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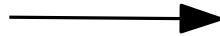
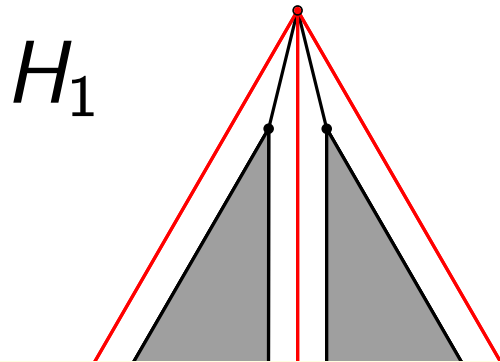
$$\text{tw}(G_i) \geq \text{mindeg}(G_i) =$$

treewidth

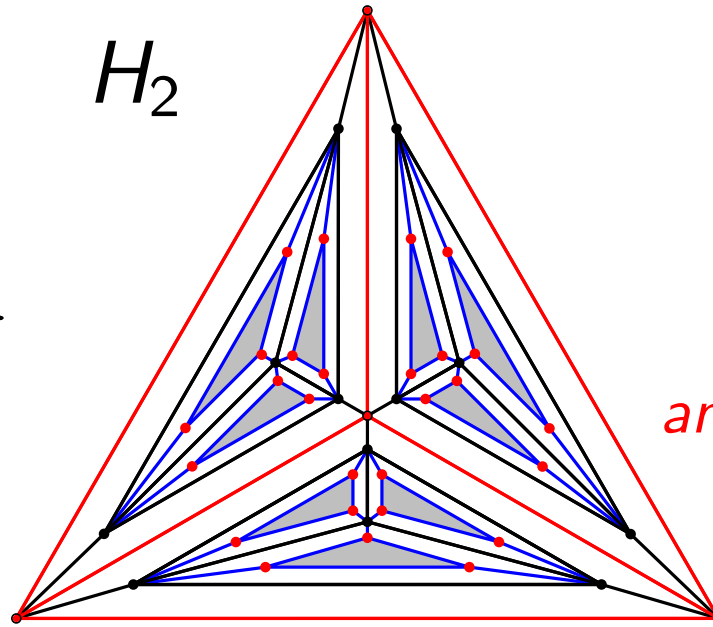
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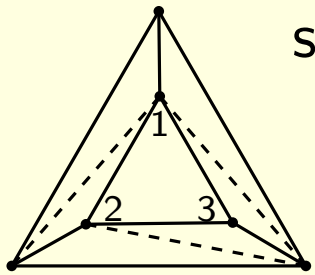
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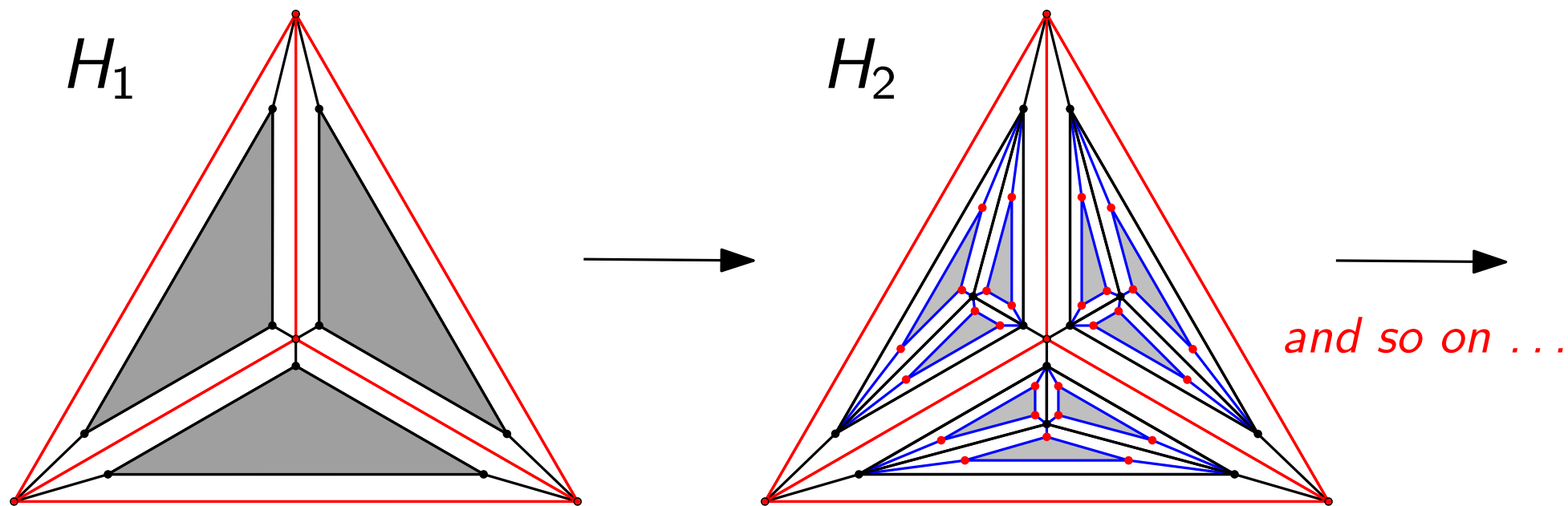
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Why?

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Properties of the family of graphs



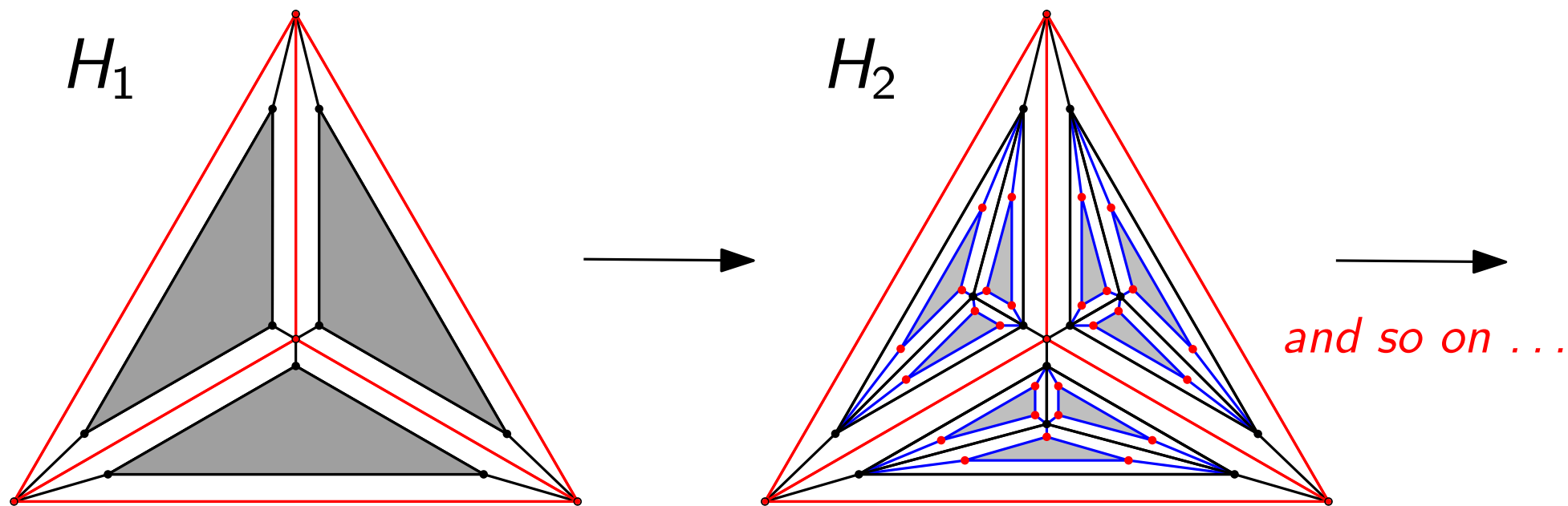
maximum degree

6

treewidth

3

Properties of the family of graphs



maximum degree

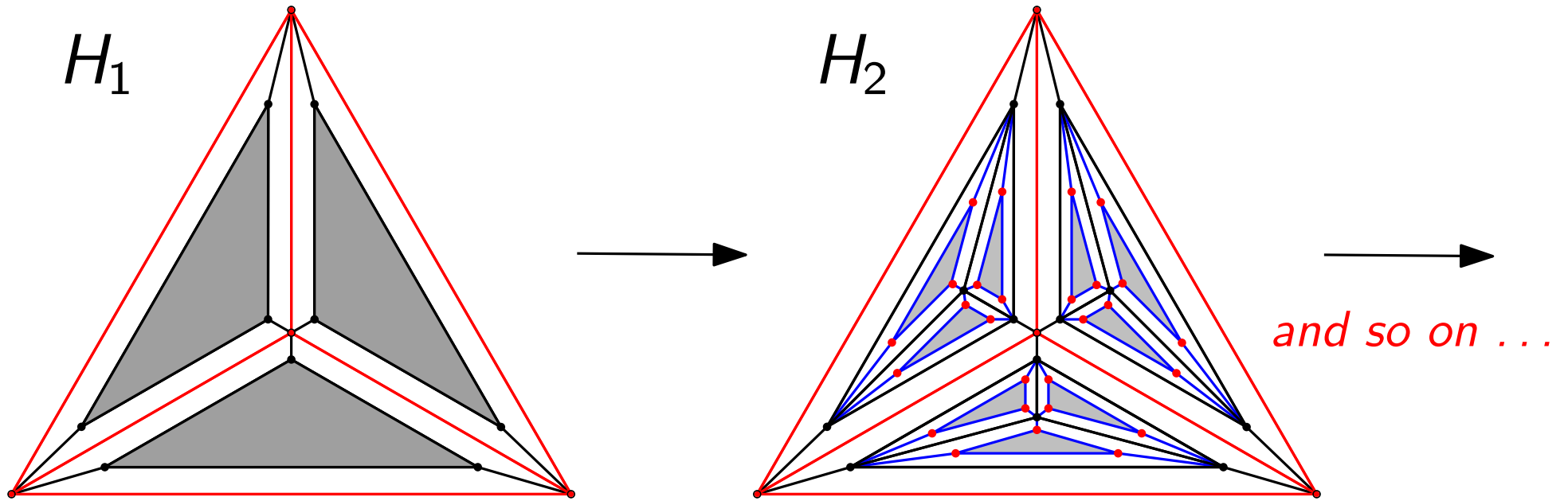
6

treewidth

3

2D weak line
cover number
unbounded

Properties of the family of graphs



Why?

maximum degree

6

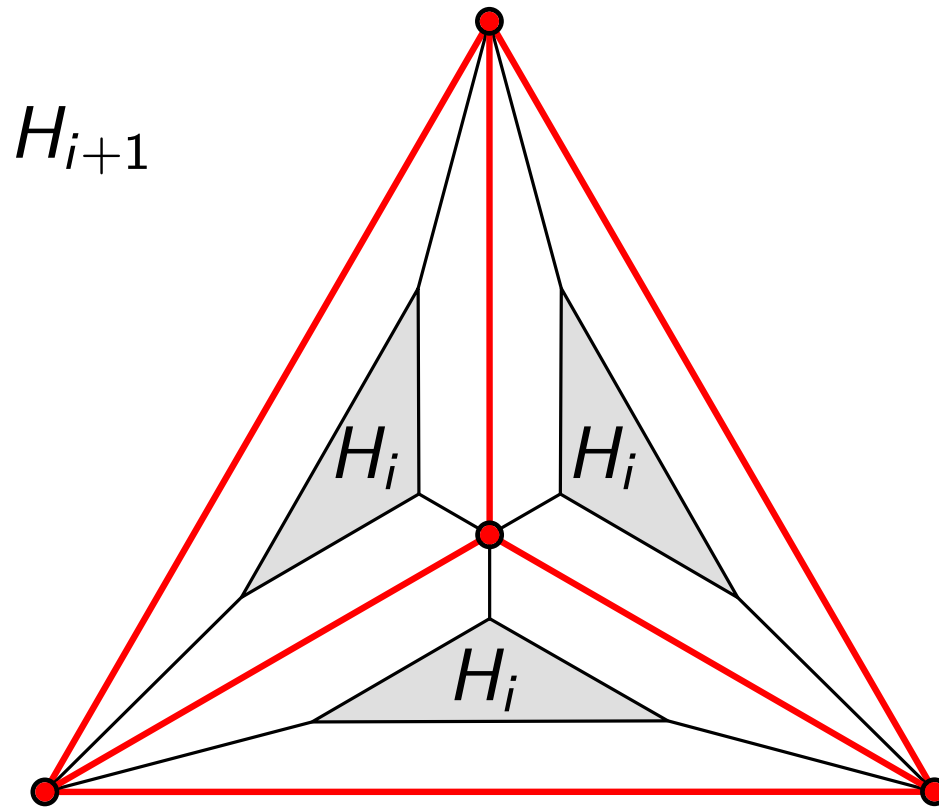
treewidth

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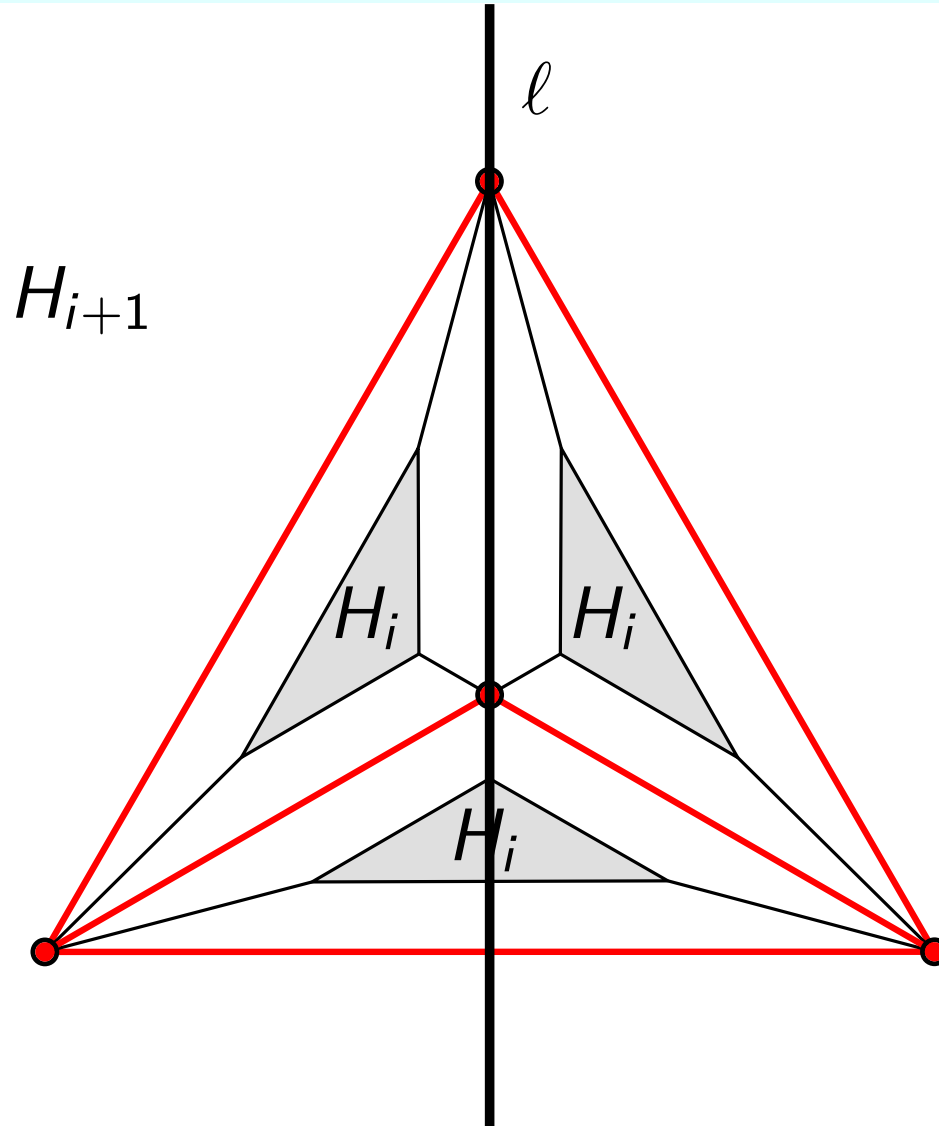
Short proof

Consider the graph H_{i+1} , $i = 1, 2, 3, \dots$



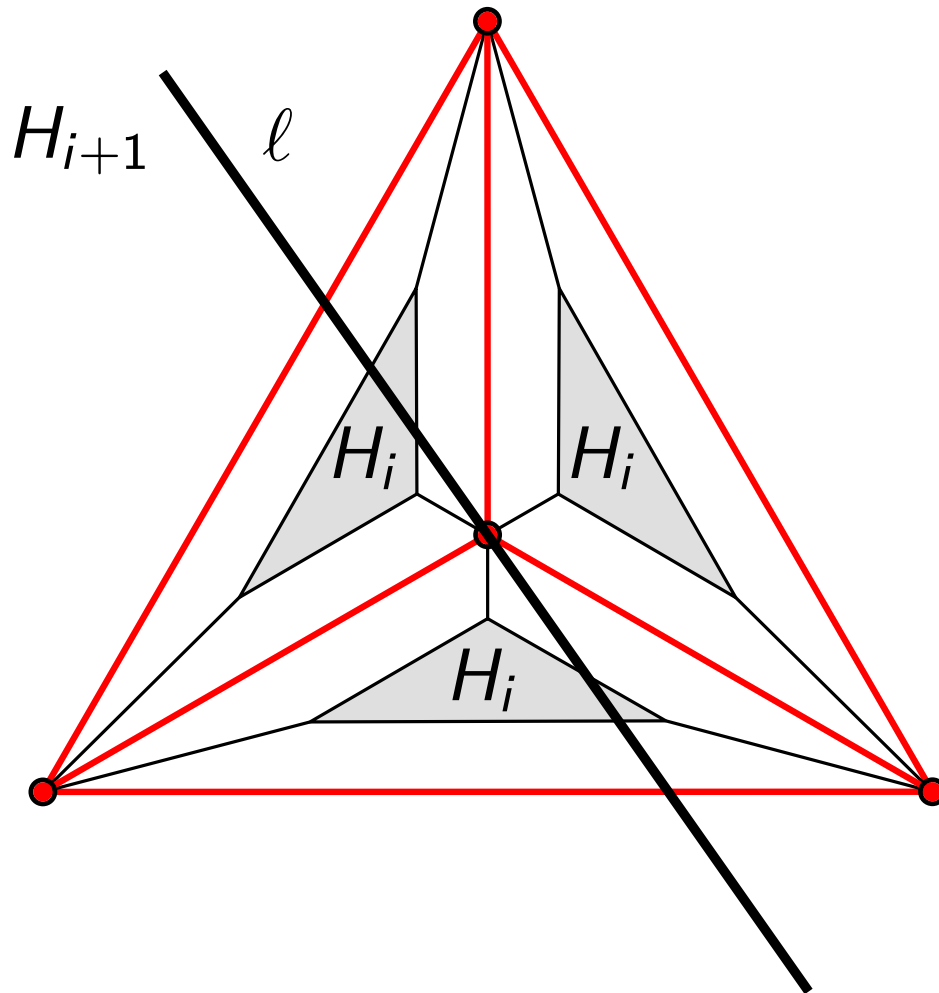
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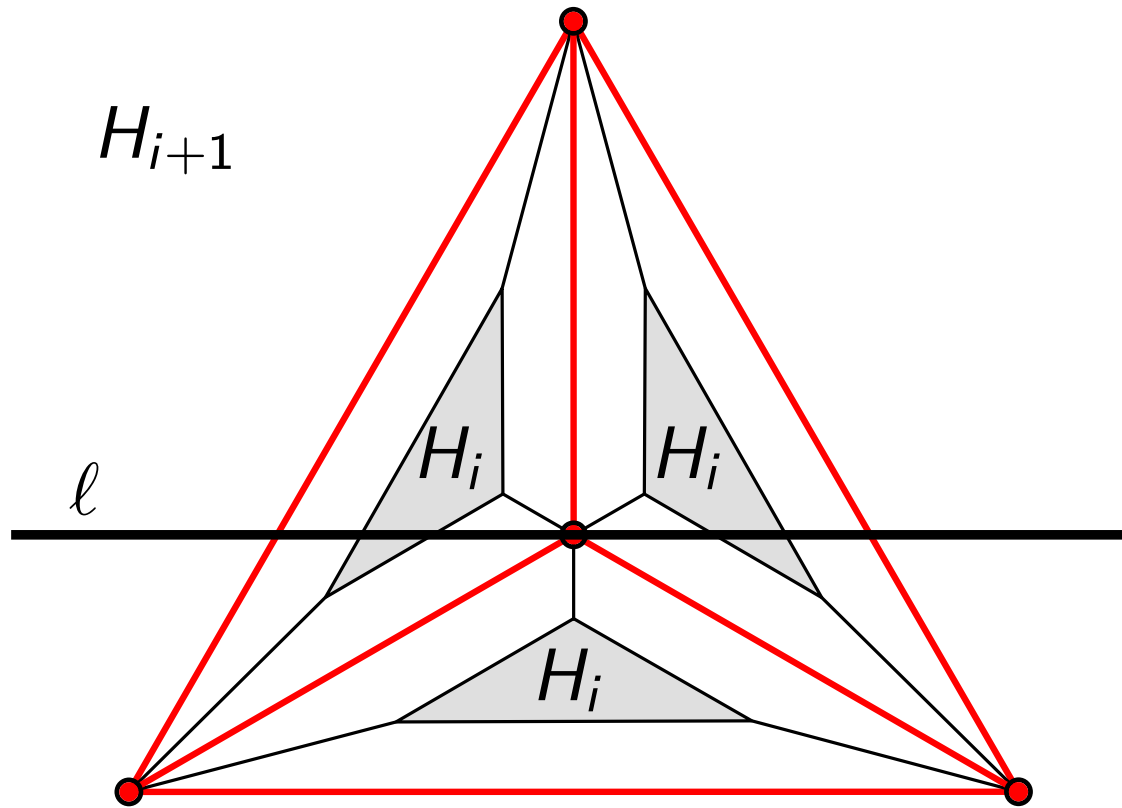
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Short proof

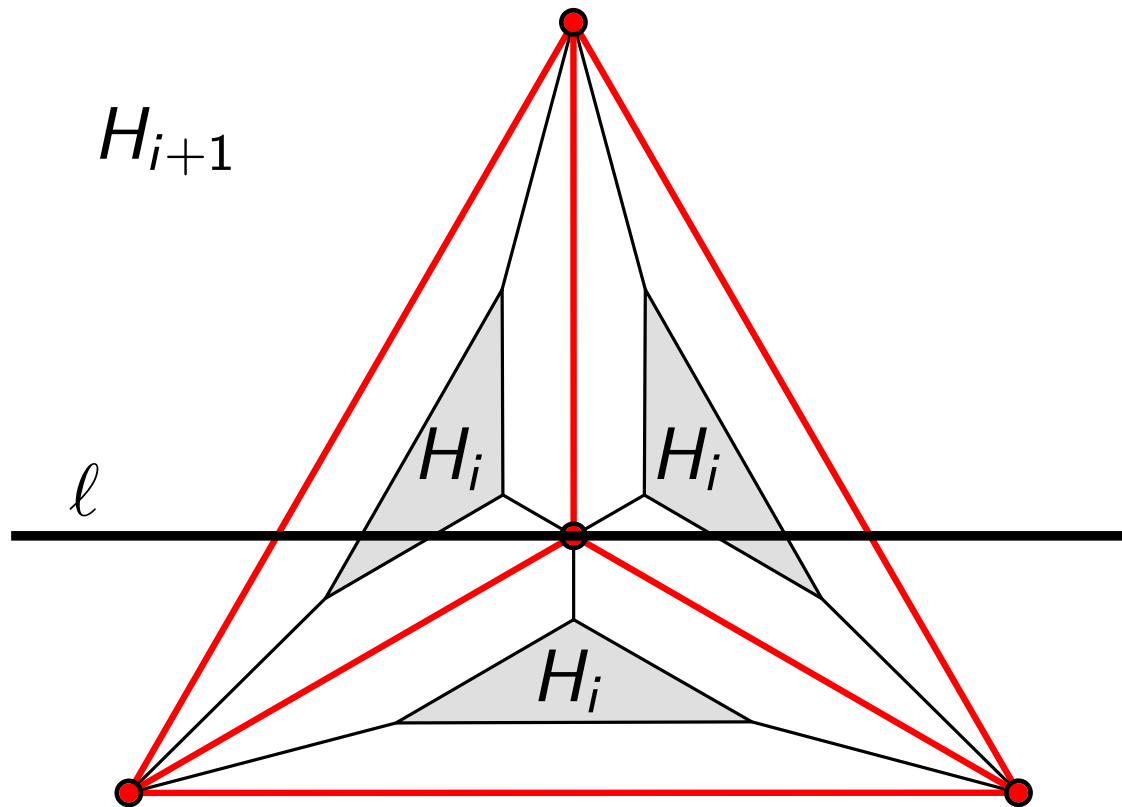
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Short proof

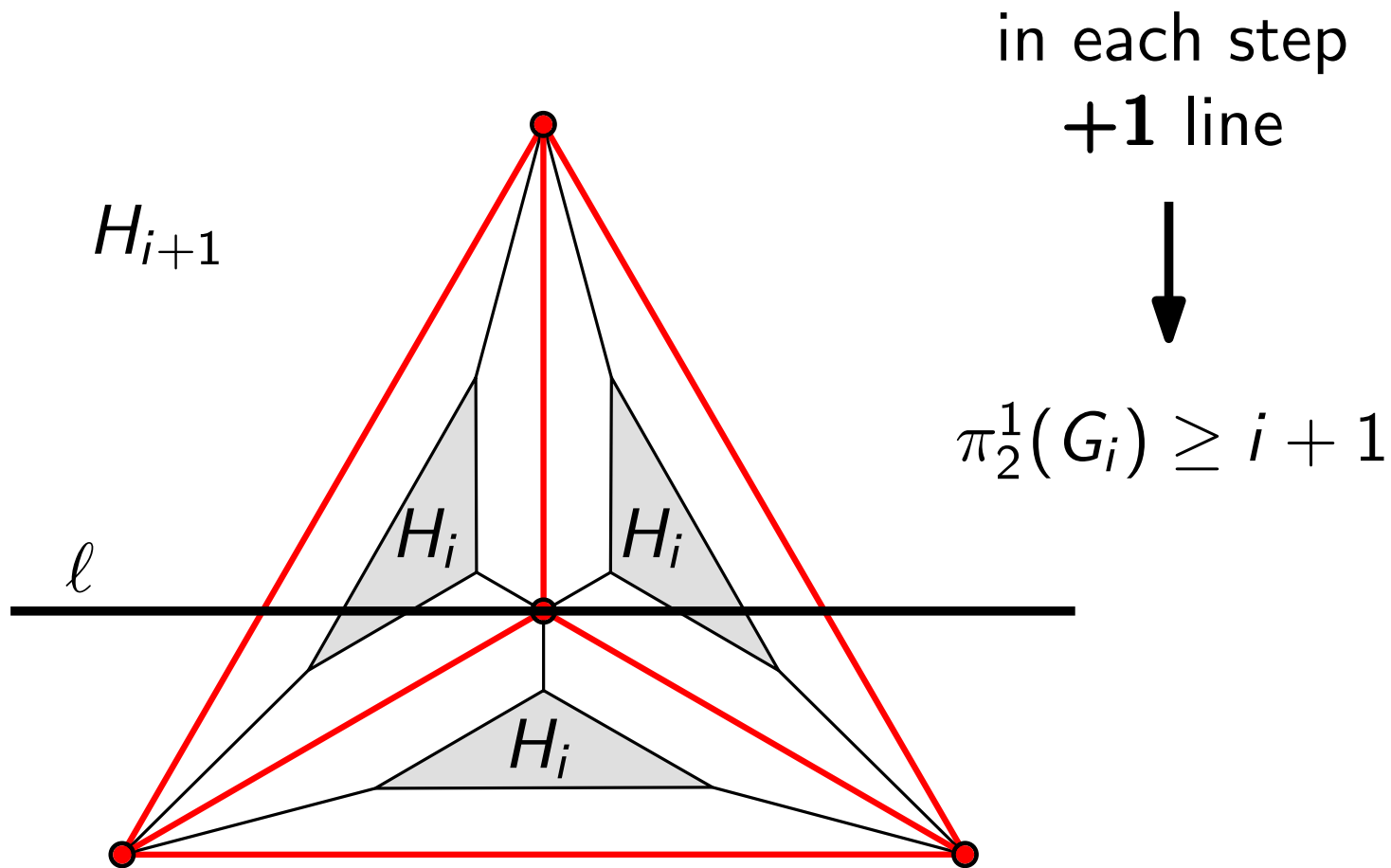
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in each step
+1 line



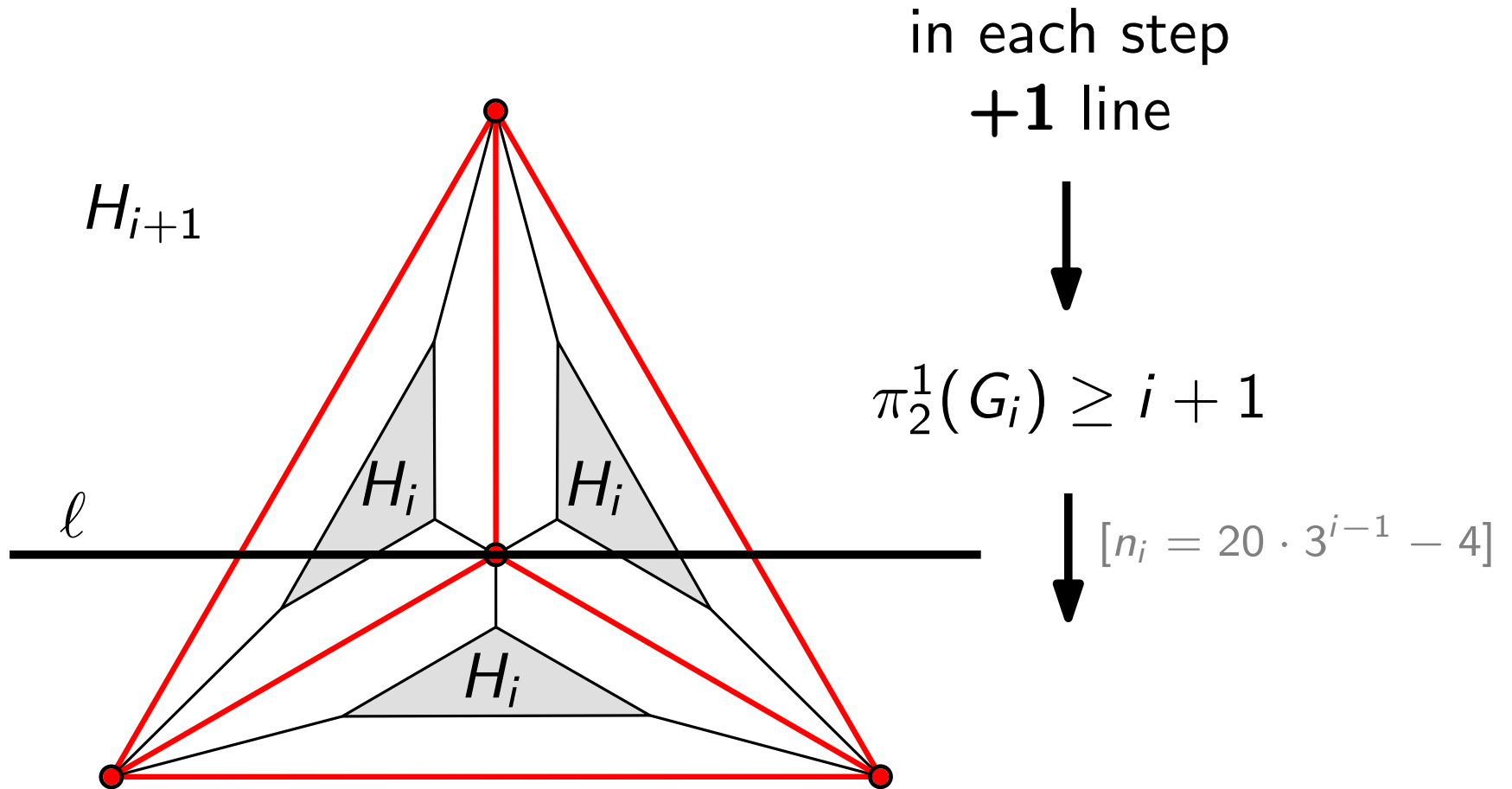
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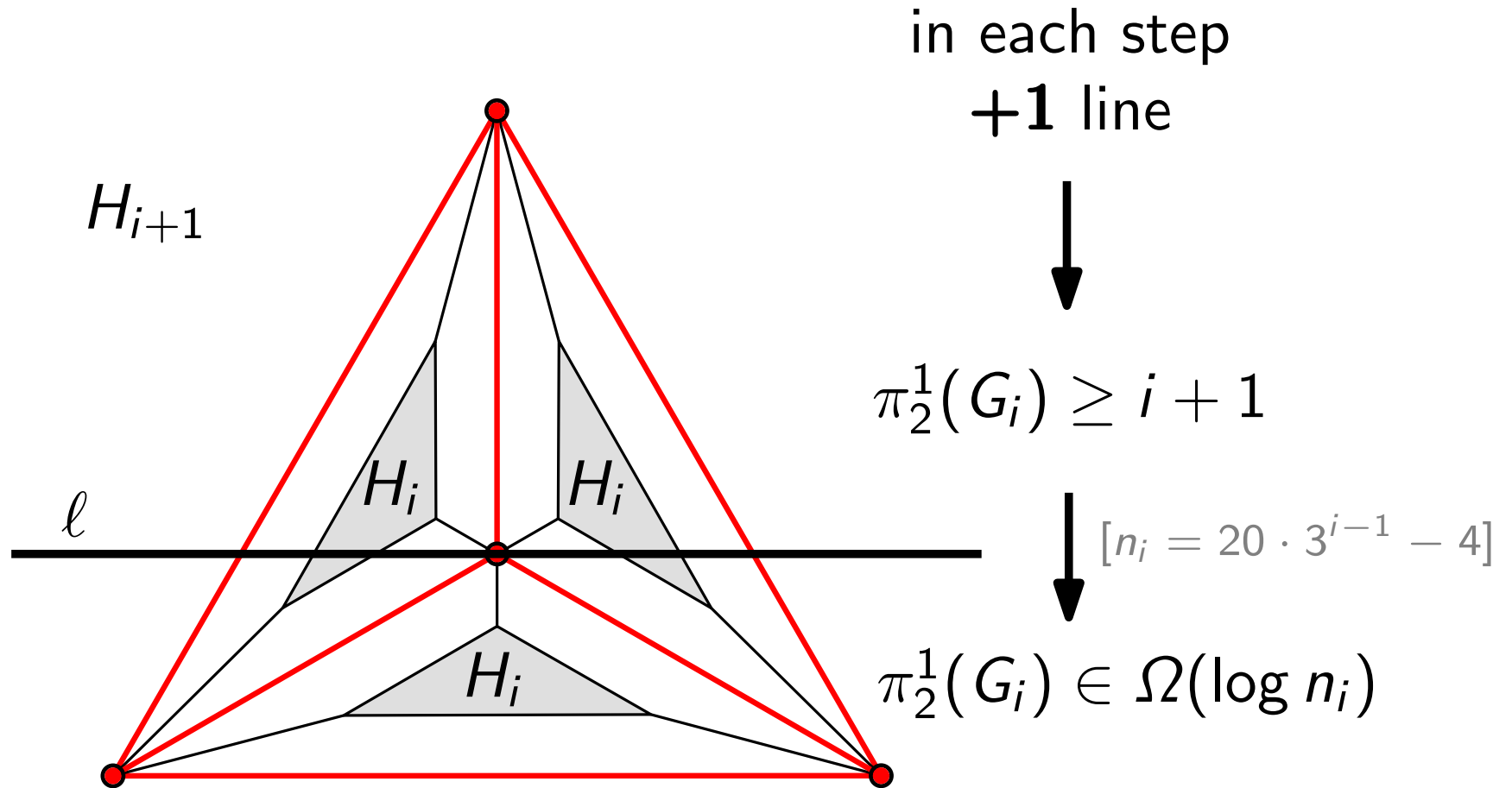
$$\pi_2^1(G_i) \geq i + 1$$



$$[n_i = 20 \cdot 3^{i-1} - 4]$$

Short proof

Consider the graph H_{i+1} , $i = 1, 2, 3, \dots$



in each step
+1 line



$$\pi_2^1(G_i) \geq i + 1$$



$$[n_i = 20 \cdot 3^{i-1} - 4]$$

$$\pi_2^1(G_i) \in \Omega(\log n_i)$$

Open problems

Problem 1

How small can we make **the maximum degree** in a family of planar graphs such that their π_2^1 -value is still unbounded?

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Is it **NP-hard** to compute $\pi_2^1(G)$ for a given graph G ?

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Problem 3

Is it **NP-hard** to compute $\pi_2^1(G)$ for a given graph G ?

Yes, by reduction from (a restricted version of) Level Planarity.

[Biedl, Evans, Felsner, Lazard, Meijer, Valtr, Whitesides, Wismath, Wolff 2018]