

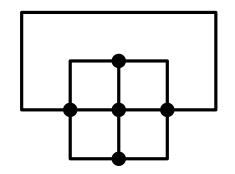


# Smooth Orthogonal Drawings of Planar Graphs

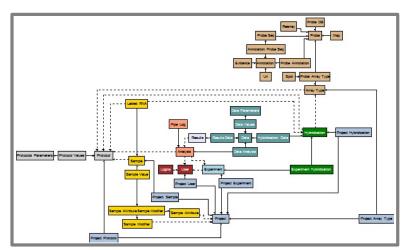
Philipp Kindermann
Chair of Computer Science I
Universität Würzburg

Joint work with Md. Jawaherul Alam, Michael A. Bekos, Michael Kaufmann, Stephen G. Kobourov & Alexander Wolff

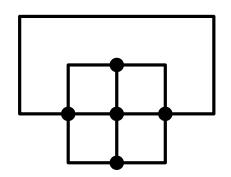
- all edge segments are horizontal or vertical
- a well-studied drawing convention
- many examples in applications



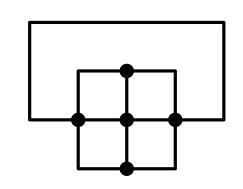
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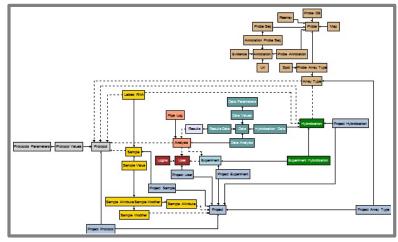


ER diagram in OGDF

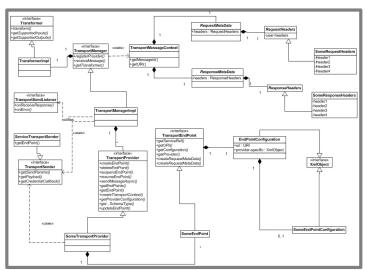


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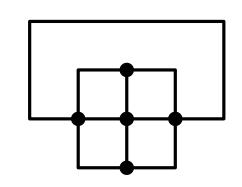


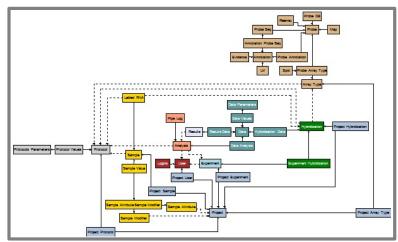
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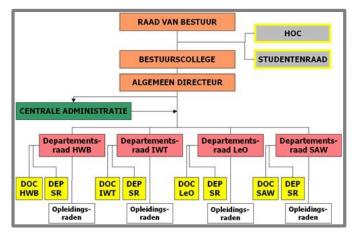
UML diagram by Oracle

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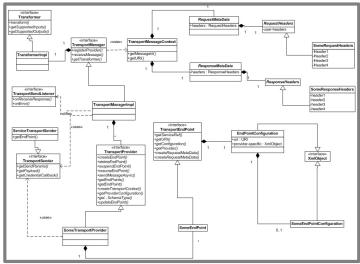




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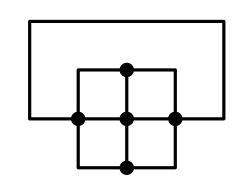


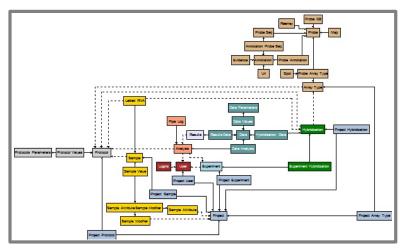
Organigram of HS Limburg



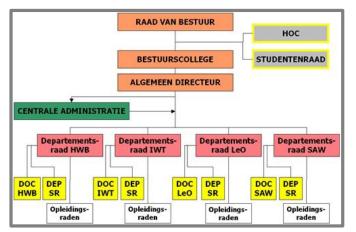
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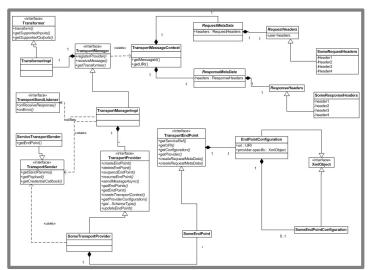




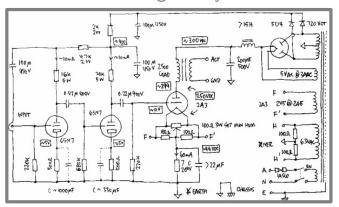
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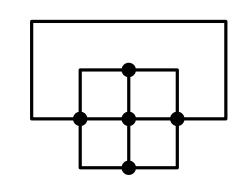


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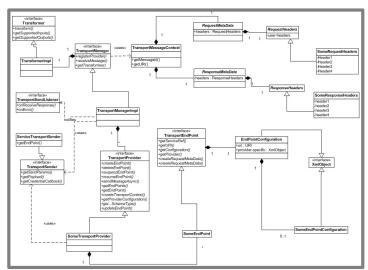
Circuit diagram by Jeff Atwood

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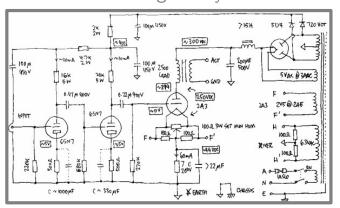




Fused Grid city layouts [By Fgrammen, via Wikimedia Commons]

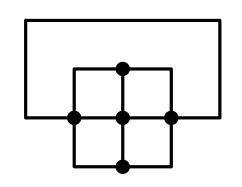


UML diagram by Oracle



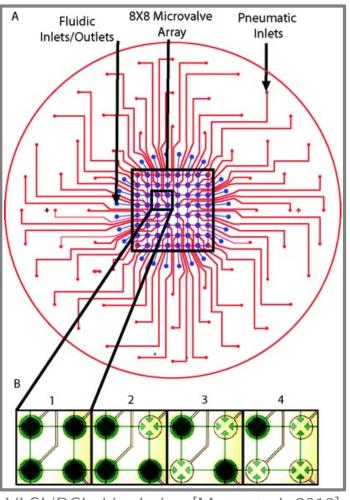
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Fused Grid city layouts [By Fgrammen, via Wikimedia Commons]



VLSI/PCI chip design [Mora et al. 2013]

### Orthogonal Layouts – Well-Known Results

#### [Tamassia, SIAM J Comp'87]

Can minimize number of bends for fixed embedding.

#### [Garg & Tamassia, SIAM J Comp'01]

Without fixed embedding, bend minimization is hard to approx.

### [Biedl & Kant, CGTA'98], [Liu et al., DAM'98]

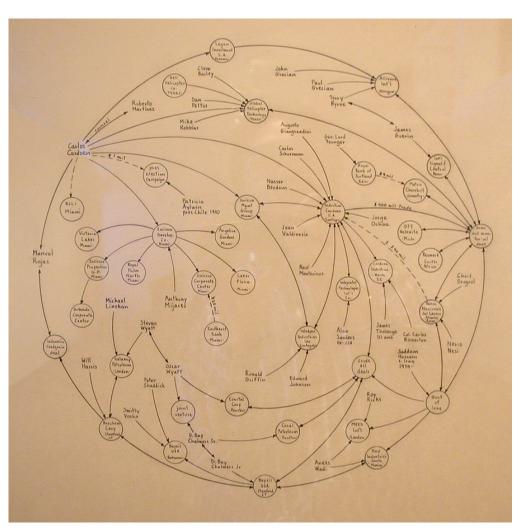
Can compute drawing on the  $(n \times n)$ -grid with  $\leq 2n + 2$  bends for any embedding (and  $\leq 2$  bends/edge – except octahedron)

### [Bläsius et al., '11]

Given an embedding and a function flex:  $E \to \mathbb{N}_{\geq 1}$ , can compute a drawing with  $\leq$  flex(e) bends/edge (if one exists).

#### Lombardi drawings

- circular arc edges
- perfect angular resolution



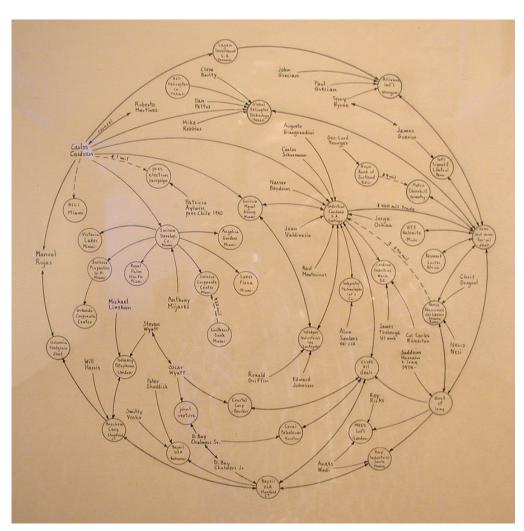
Mark Lombardi (1951–2000)

#### Lombardi drawings

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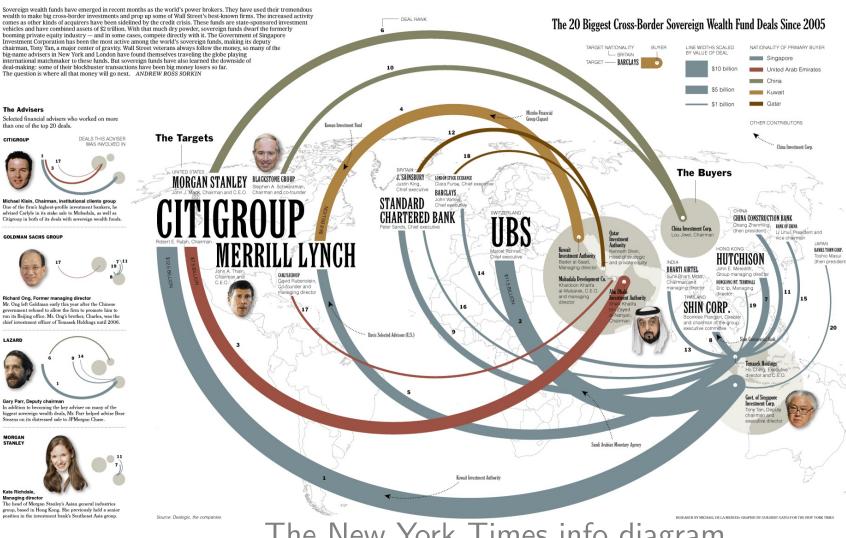
#### *k*-Lombardi drawings

each edge sequence of k
 circular arcs



Mark Lombardi (1951–2000)

#### The New Global Wealth Machine



#### The Lawyers

Selected lawyers who worked on more than one of the top 20 deals.

CLIFFORD CHANCE



James Baird, Partner and global head of private equity Mr. Baird's firm, based in London, was one of the early firms to make a bet on Asia by staffing up there before some of the traditional white-shoe Wall Street firms

DAVIS POLK A

Randall D. Guynn, Partne

As head of the firm's financial institutions group, he has advised on many international deals in Europe and Asia. He also worked on the team that advised Morgan Stanley in its \$5.5 billion stake sale to China's sovereign wealth fund.

LINKLATERS



Based in Singapore, Mr. Good is the firm's man-on-the ground in Asia. He has worked for Linklaters in Asia



A longtime hand in the Middle East, Mr. Besen's deep relationships have helped his firm carve out one of the strongest niches in the region.

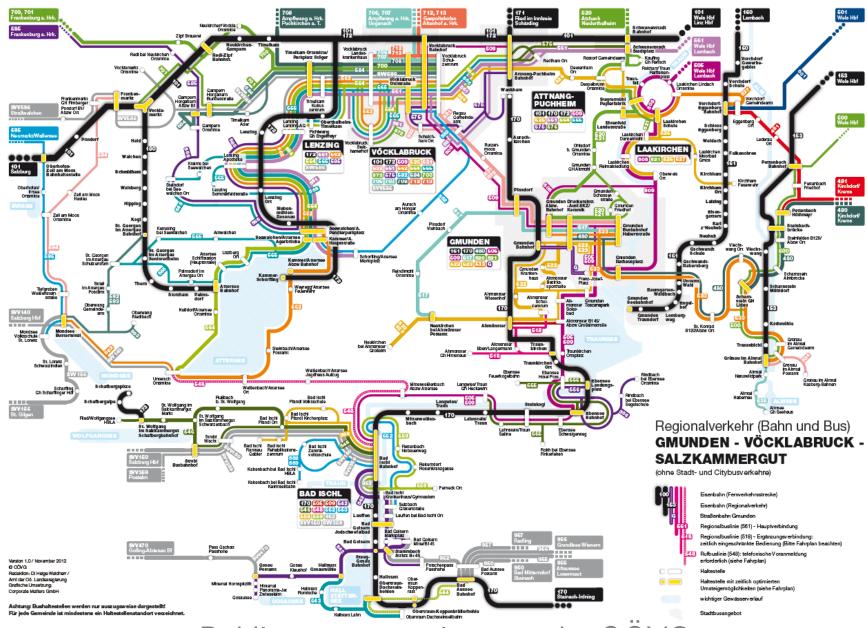




The New York Times info diagram

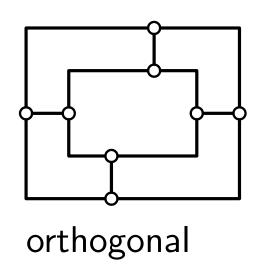


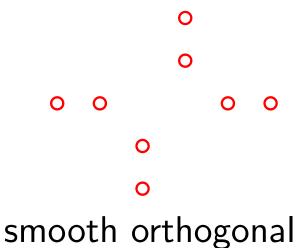
Terra Nova - Design Overdrive by Carlos Barbosa



Public transportation map by OÖVG

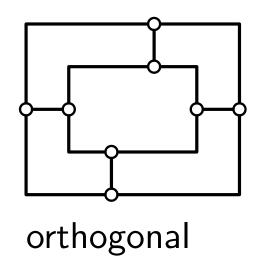
Combine both worlds:

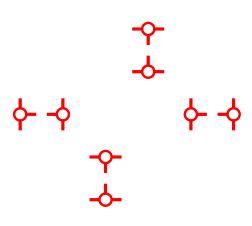




#### Combine both worlds:

edges leave and enter vertices horizontally or vertically

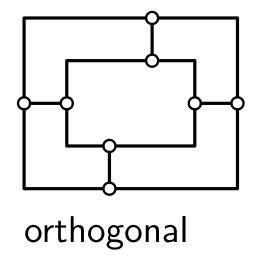


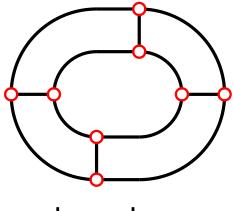


smooth orthogonal

#### Combine both worlds:

- edges leave and enter vertices horizontally or vertically
- each edge is drawn as a sequence of axis-aligned line segments and circular-arc segments without bends

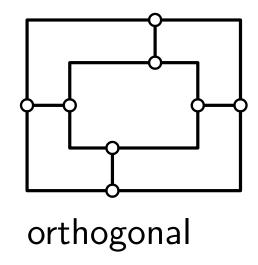


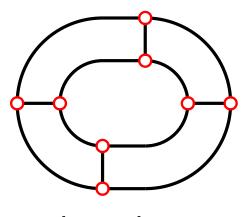


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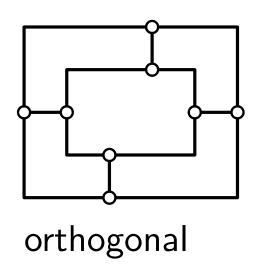
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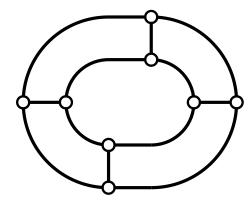
- edges leave and enter vertices horizontally or vertically
- each edge is drawn as a sequence of axis-aligned line segments and circular-arc segments without bends
- there are no edge-crossings (for planar graphs)





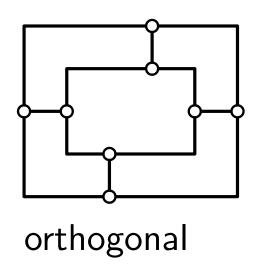
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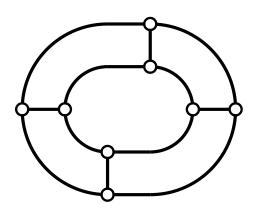




smooth orthogonal

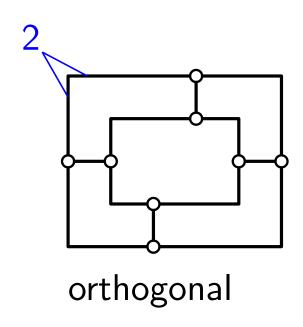
complexity of an edge: number of arcs

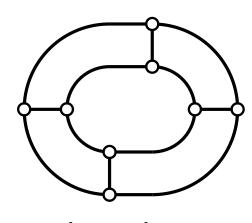




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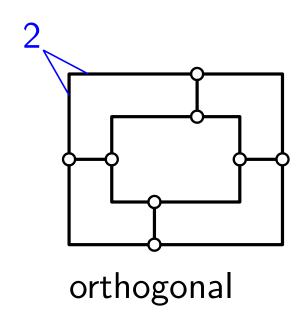
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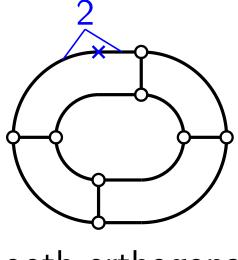




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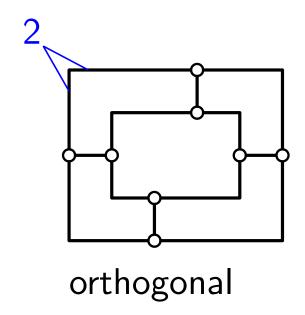


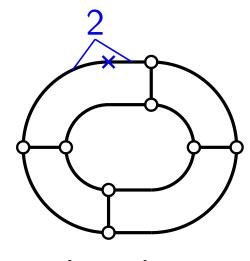


smooth orthogonal

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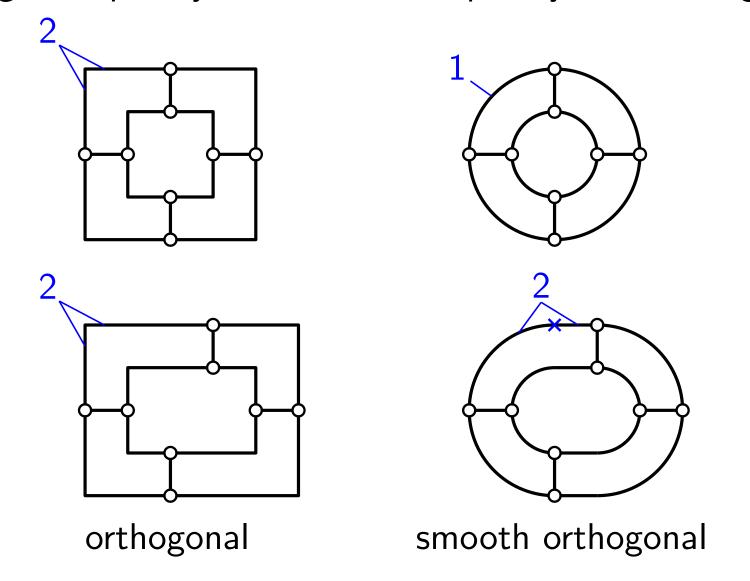
edge complexity: maximum complexity over all edges





smooth orthogonal

complexity of an edge: number of arcs edge complexity: maximum complexity over all edges



biconnected 4-planar graph → orthogonal complexity-3 layout

• choose vertices s and t

- choose vertices s and t
- place vertices by their st-numbering

biconnected 4-planar graph  $\rightarrow$  orthogonal complexity-3 layout

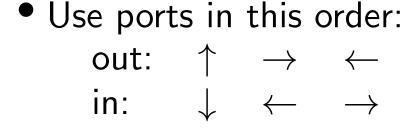
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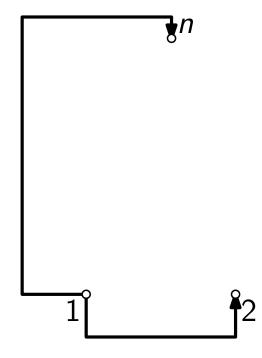
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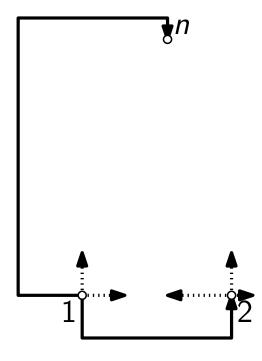
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out: \uparrow \rightarrow \leftarrow in: \downarrow \leftarrow \rightarrow
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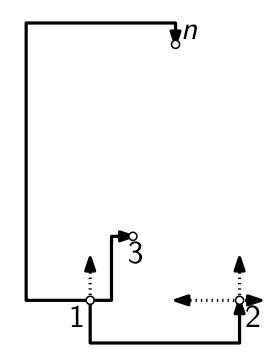




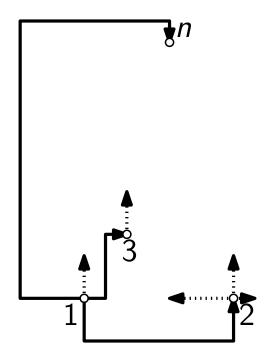
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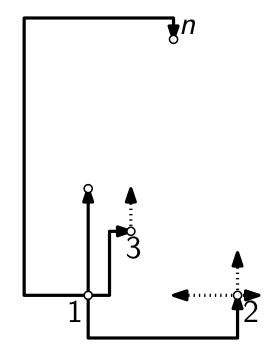
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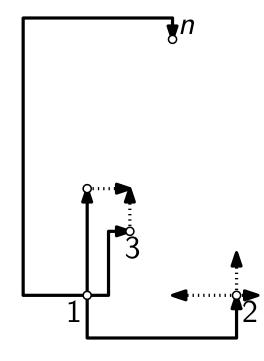
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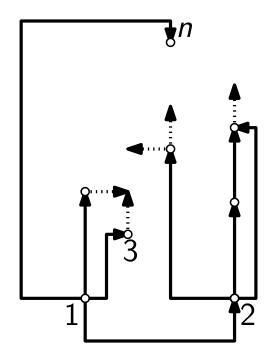
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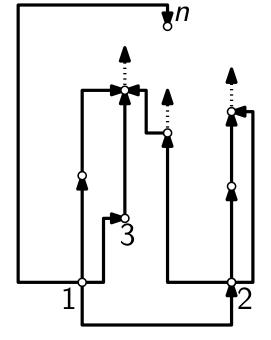
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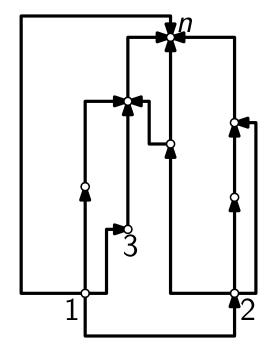
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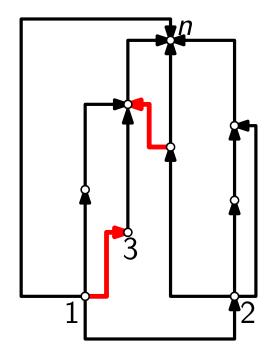


biconnected 4-planar graph  $\rightarrow$  orthogonal complexity-3 layout

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Eliminate S-shapes

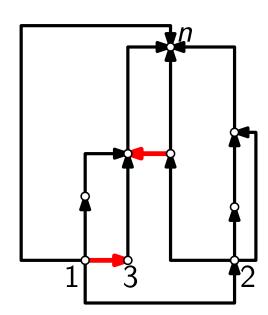


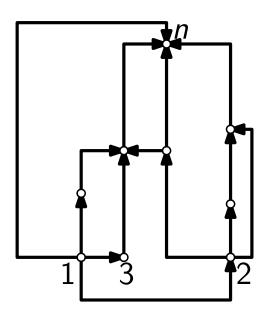
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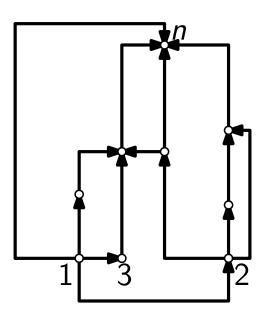
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$$\uparrow$$
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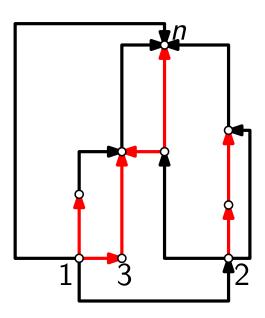


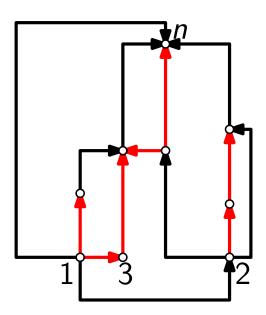


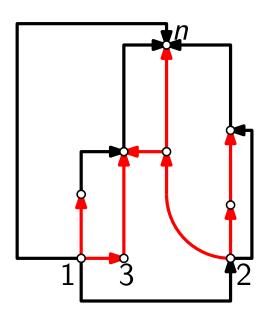
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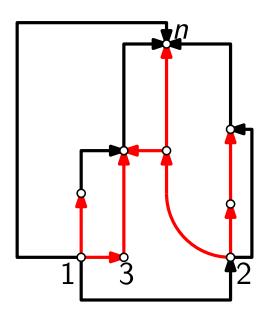
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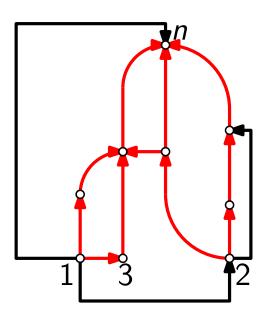


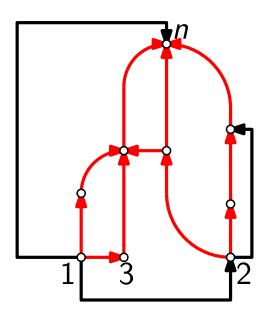


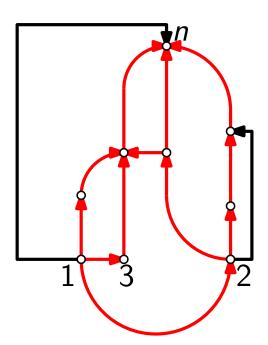
$$| \rightarrow | \rightarrow \downarrow \qquad \Gamma \rightarrow \uparrow$$



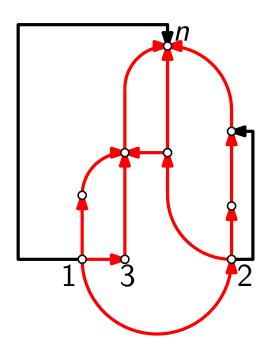
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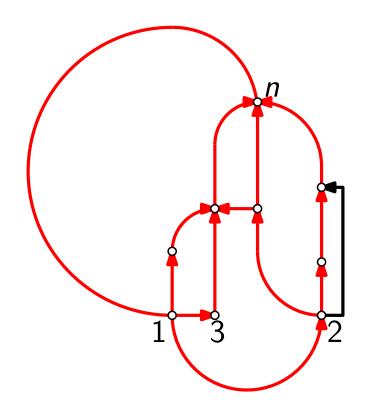


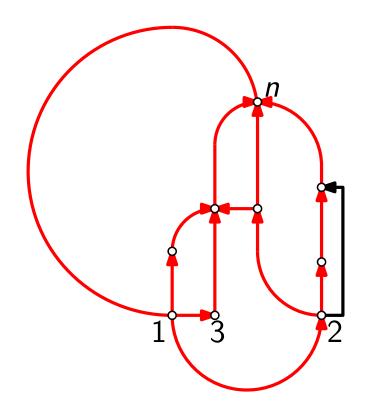


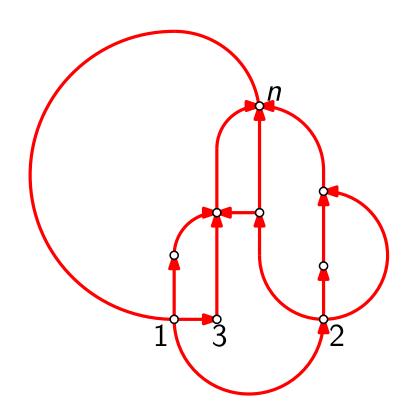
$$| \rightarrow |$$
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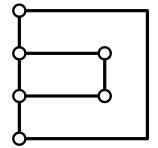


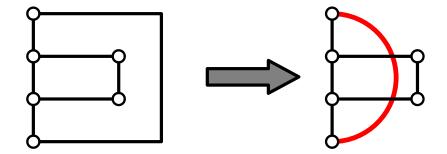
$$| \rightarrow |$$
  $| \rightarrow \downarrow$   $| \rightarrow \downarrow$   $| \rightarrow \downarrow$ 

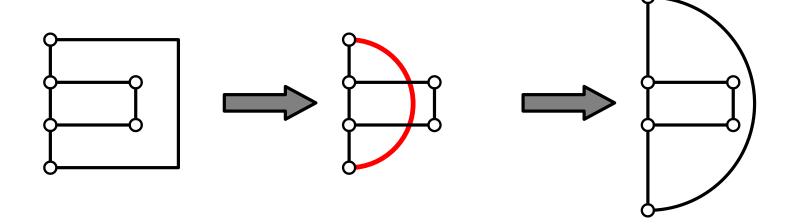


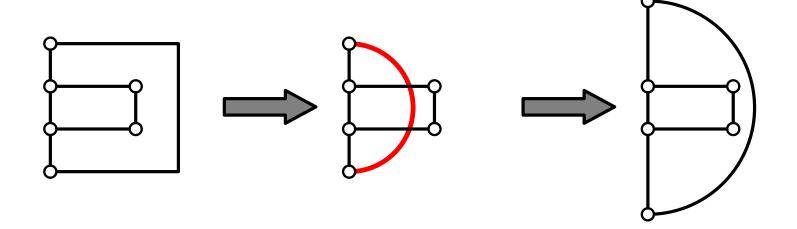


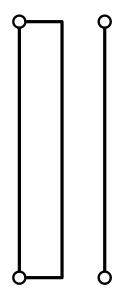


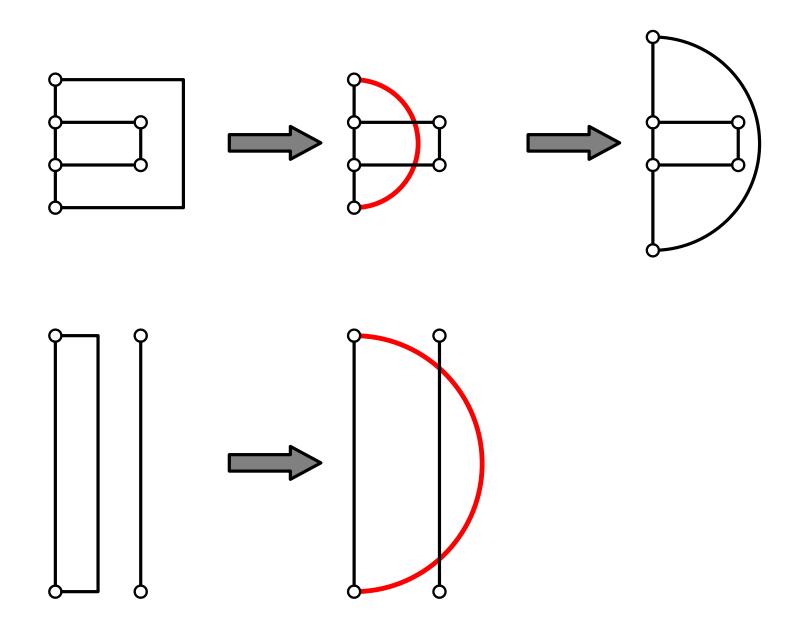


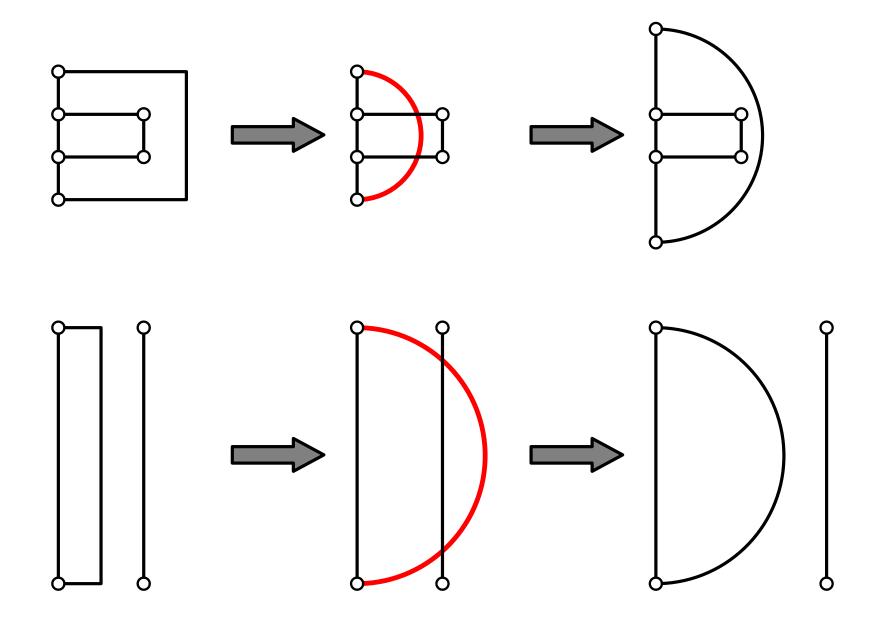






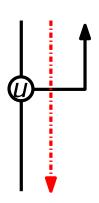






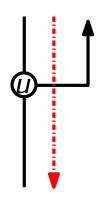
Def.

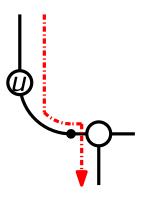
Def. • *y*-monotone curve



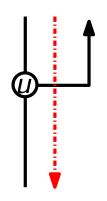
Def. • *y*-monotone curve

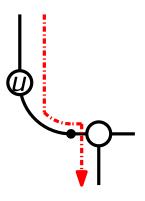
consists of horizontal, vertical and circular segments



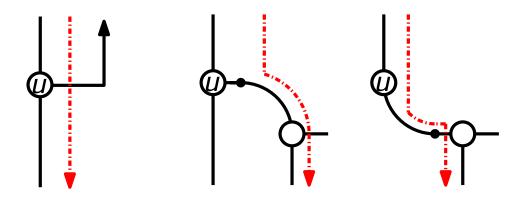


- Def. *y*-monotone curve
  - consists of horizontal, vertical and circular segments
  - divides the current drawing into a left and a right part

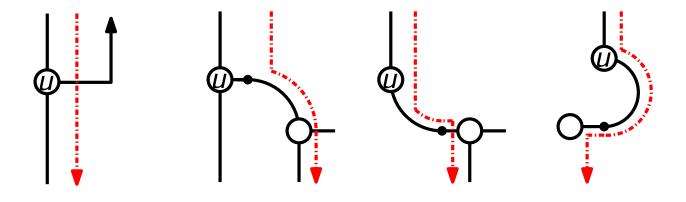




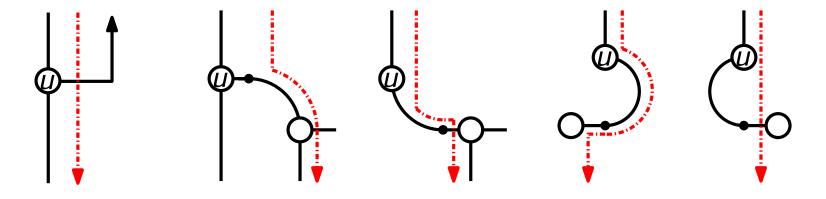
- Def. *y*-monotone curve
  - consists of horizontal, vertical and circular segments
  - divides the current drawing into a left and a right part
  - intersects only horizontal segments



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  - consists of horizontal, vertical and circular segments
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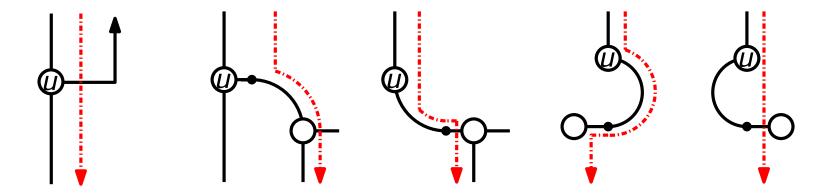


- Def. *y*-monotone curve
  - consists of horizontal, vertical and circular segments
  - divides the current drawing into a left and a right part
  - intersects only horizontal segments



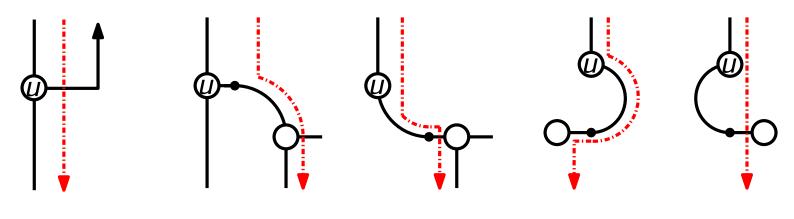
Def. • *y*-monotone curve

- consists of horizontal, vertical and circular segments
- divides the current drawing into a left and a right part
- intersects only horizontal segments

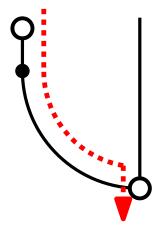


Problems:

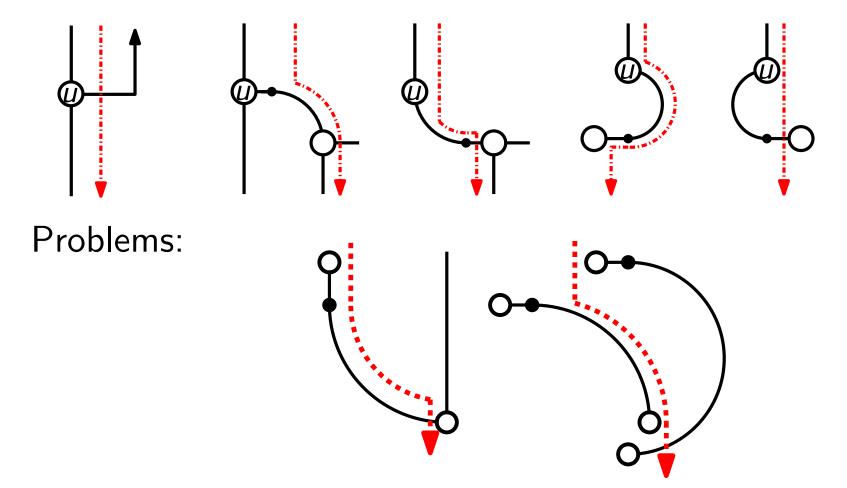
- Def. *y*-monotone curve
  - consists of horizontal, vertical and circular segments
  - divides the current drawing into a left and a right part
  - intersects only horizontal segments



Problems:

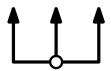


- Def. *y*-monotone curve
  - consists of horizontal, vertical and circular segments
  - divides the current drawing into a left and a right part
  - intersects only horizontal segments



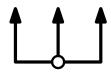
### **Invariants**

 $(I_1)$  Every open edge is associated with a column

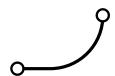


### **Invariants**

 $(I_1)$  Every open edge is associated with a column

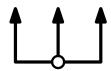


 $(I_2)$  An L-shape always contains a horizontal segment; it never contains a vertical segment.



### **Invariants**

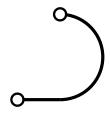
 $(I_1)$  Every open edge is associated with a column



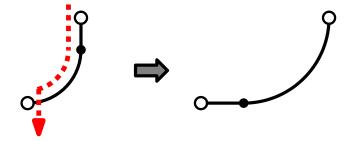
 $(I_2)$  An L-shape always contains a horizontal segment; it never contains a vertical segment.



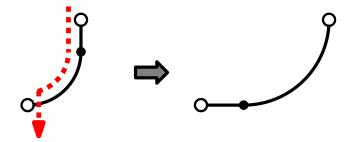
 $(I_3)$  A C-shape always has a horizontal segment incident to its bottom vertex.



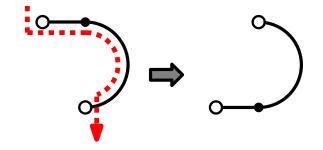
### L-shape



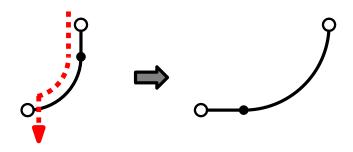
### L-shape



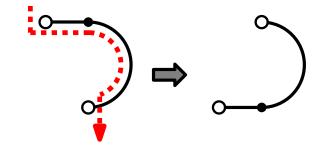
### C-shape



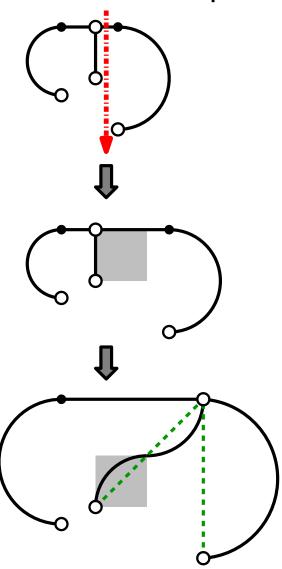
L-shape

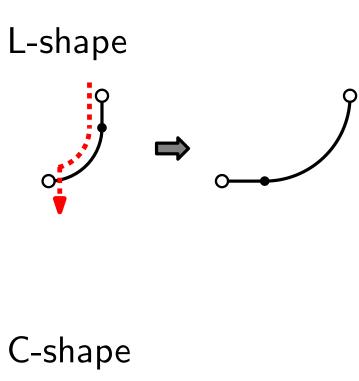


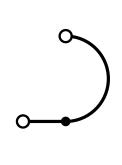
C-shape

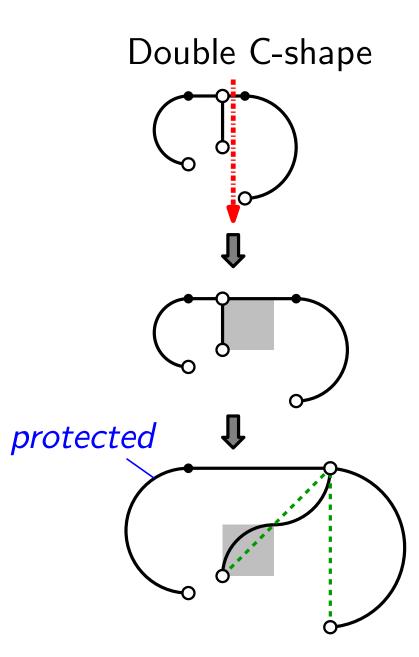


Double C-shape



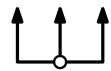






### Invariants, updated

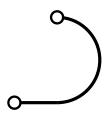
 $(I_1)$  Every open edge is associated with a column

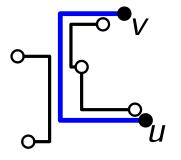


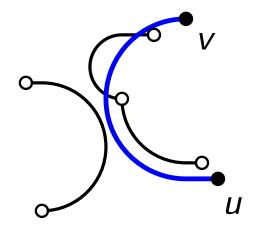
 $(I_2)$  An L-shape always contains a horizontal segment; it never contains a vertical segment.



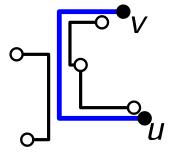
 $(I_3)$  An *unprotected* C-shape always has a horizontal segment incident to its bottom vertex.

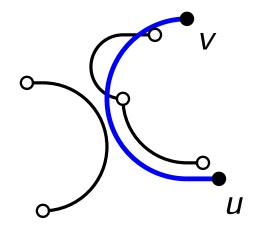




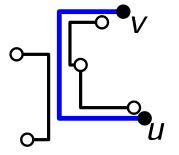


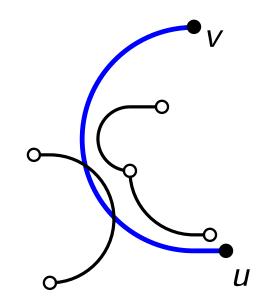
1. move *v* up



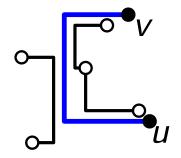


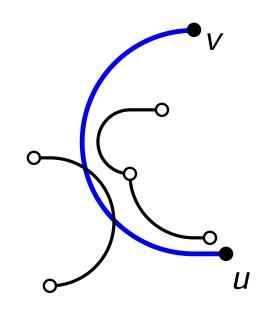
1. move *v* up



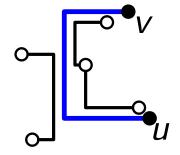


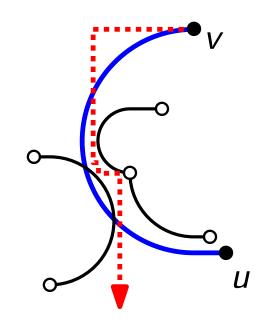
- 1. move *v* up
- 2. find a cut



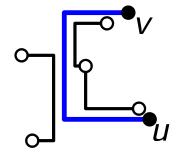


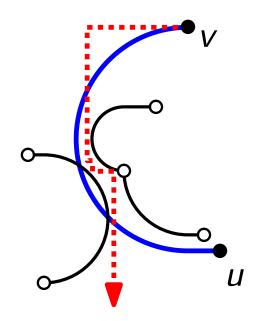
- 1. move *v* up
- 2. find a cut



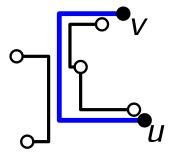


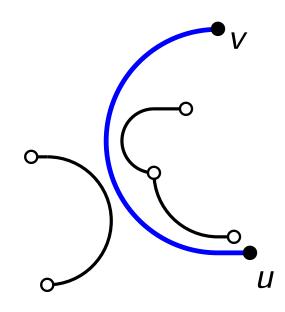
- 1. move *v* up
- 2. find a cut
- 3. move vertices to the left

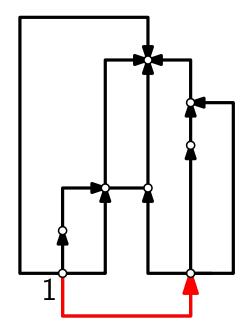


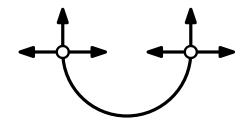


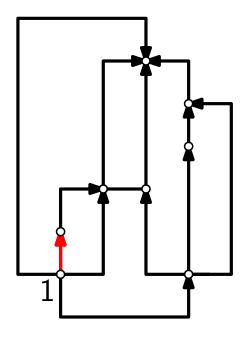
- 1. move *v* up
- 2. find a cut
- 3. move vertices to the left

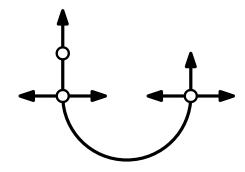


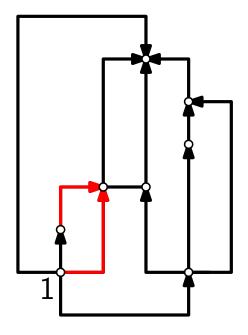


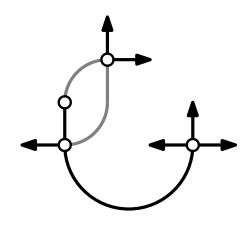


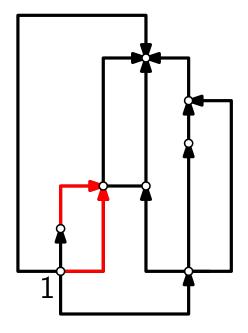


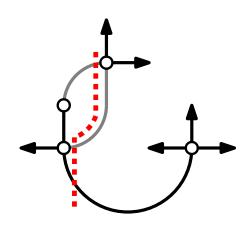


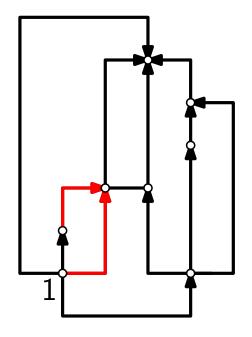


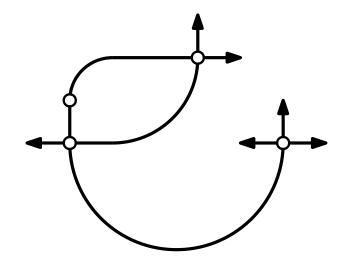


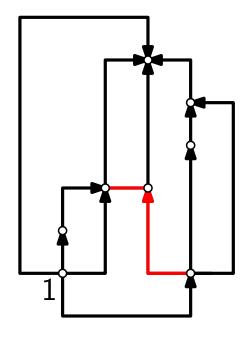


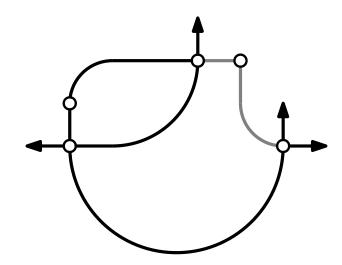


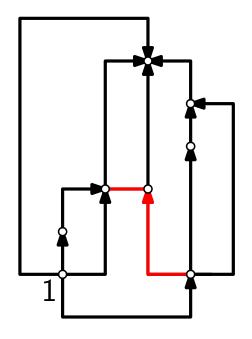


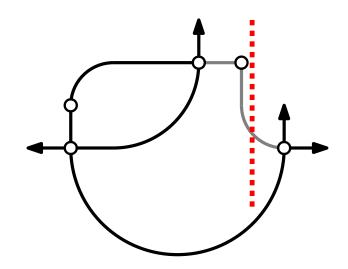


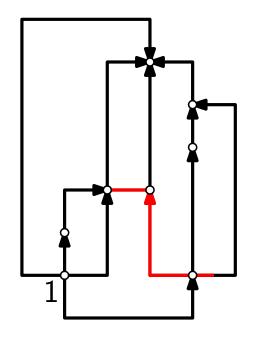


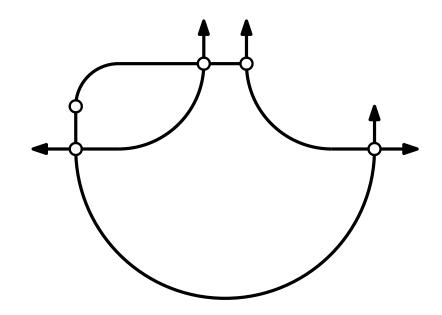


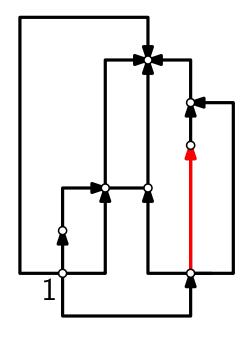


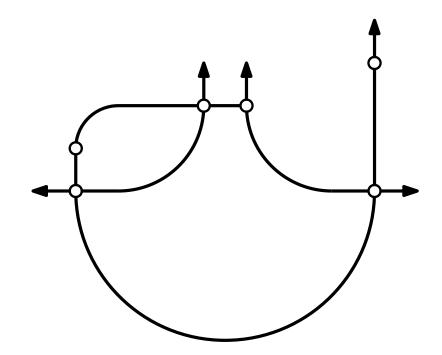


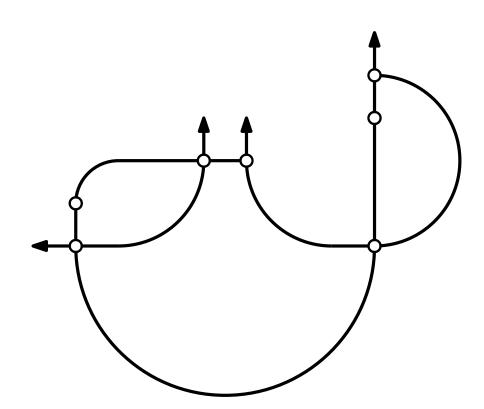


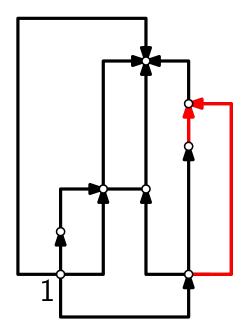


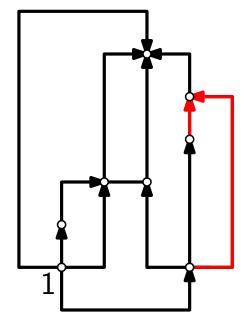


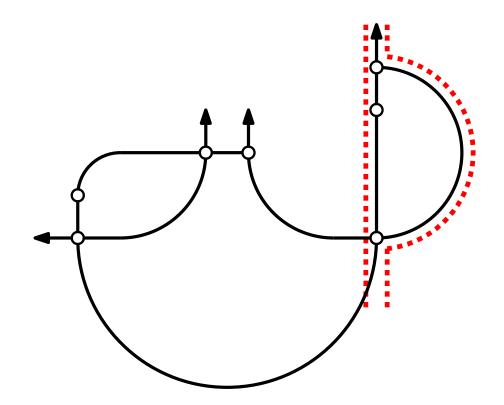


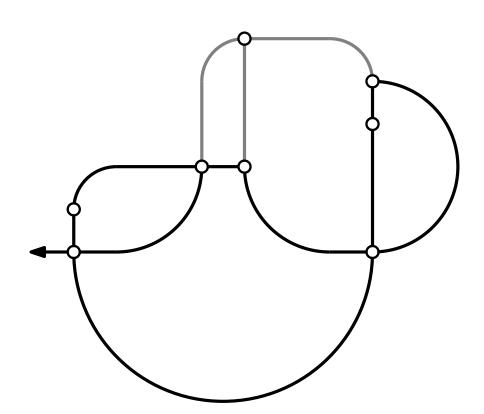


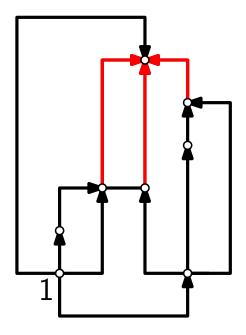


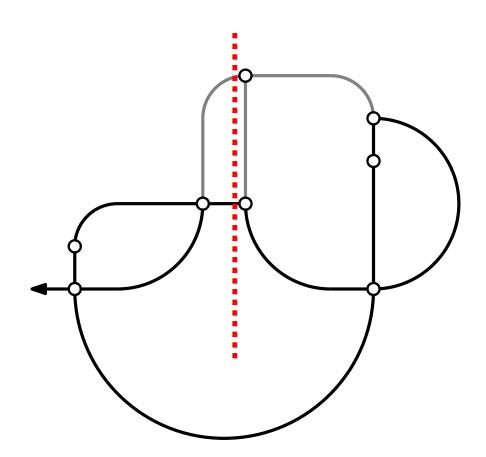


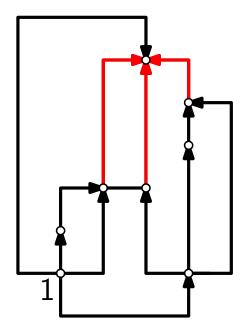


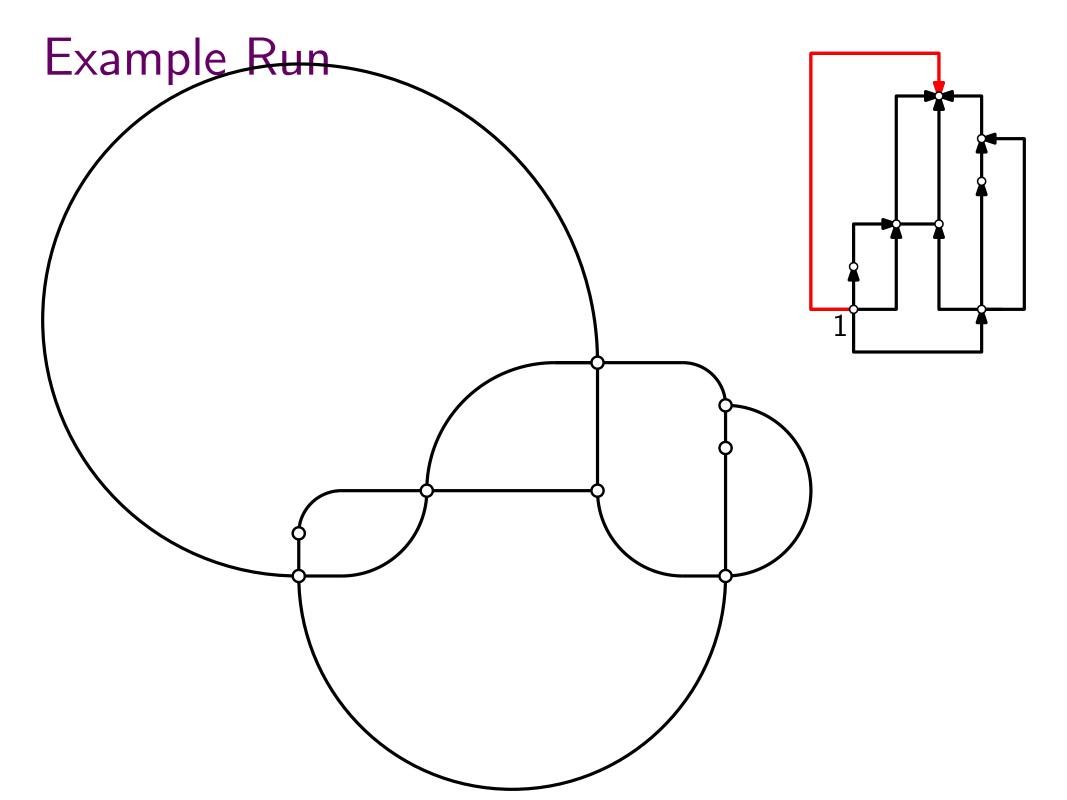


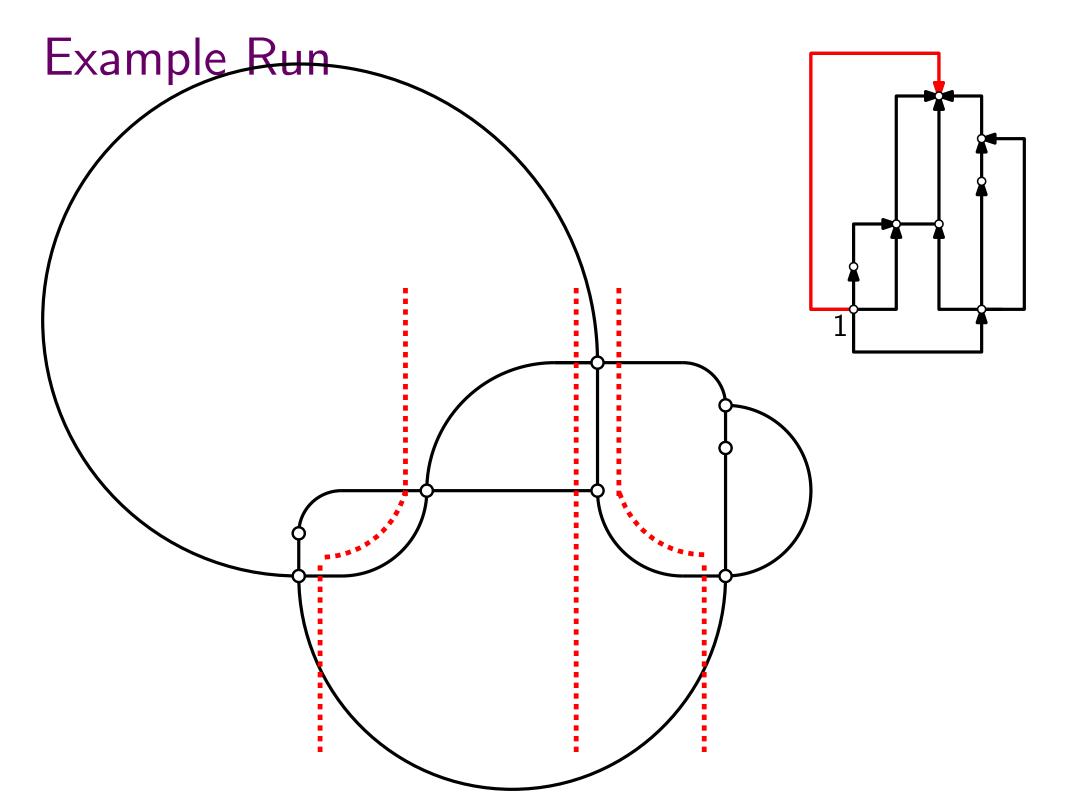


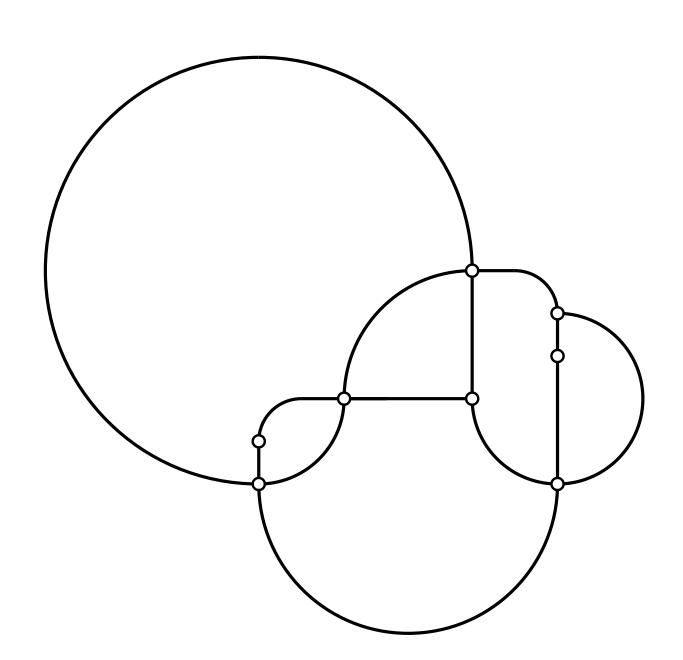


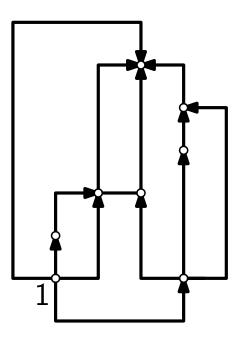


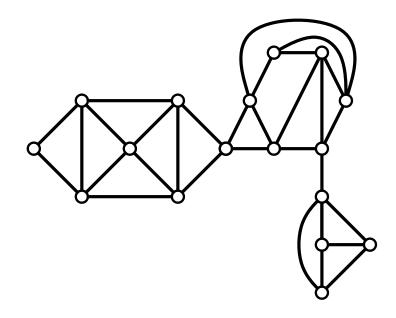




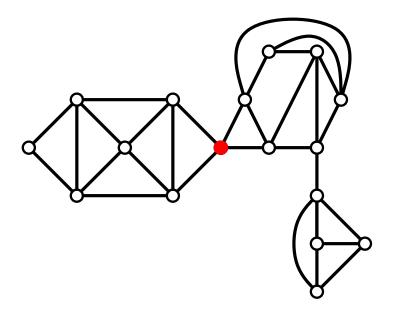




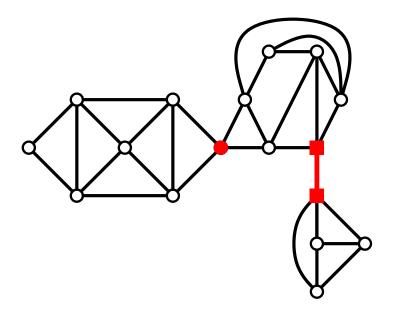




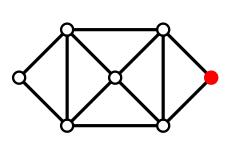
cutvertices

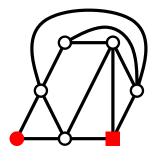


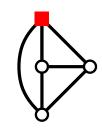
- cutvertices
- bridges



- cutvertices
- bridges

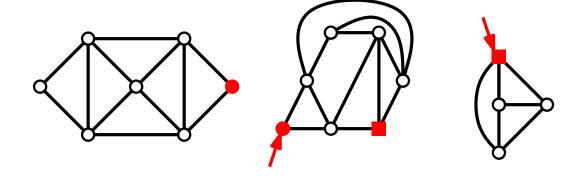






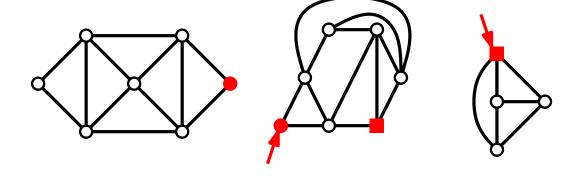
- cutvertices
- bridges

Draw specific vertex on outer face



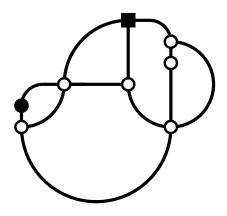
- cutvertices
- bridges

Draw specific vertex on outer face Draw cut vertices with right angles

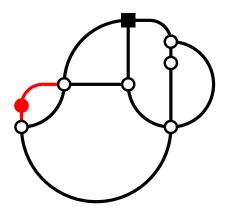


- cutvertices
- bridges

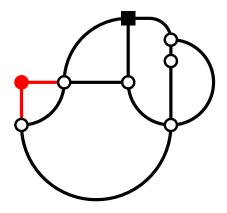
Draw specific vertex on outer face Draw cut vertices with right angles



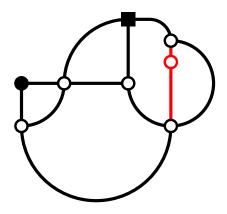
- cutvertices
- bridges



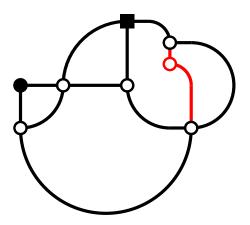
- cutvertices
- bridges



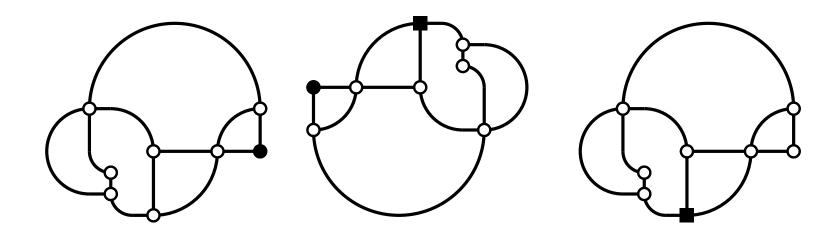
- cutvertices
- bridges



- cutvertices
- bridges



- cutvertices
- bridges



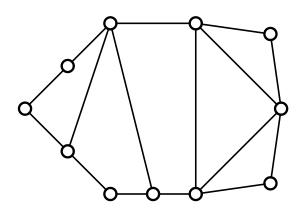
- cutvertices
- bridges

4-planar graph  $\rightarrow$  smooth orthogonal complexity-2 layout

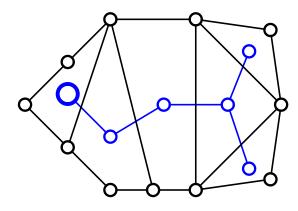
- cutvertices
- bridges

Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.

Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.

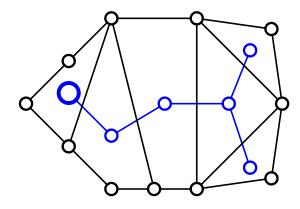


Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.

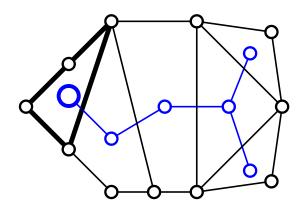


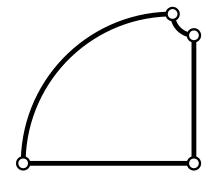
Consider the dual tree

Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.

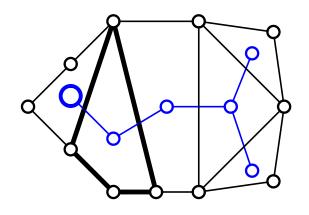


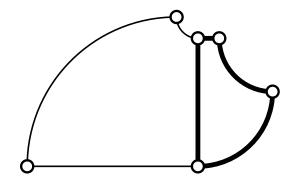
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.



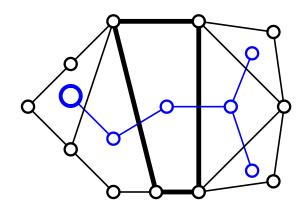


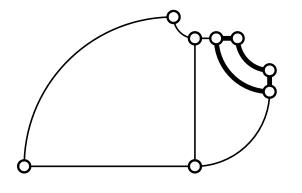
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.



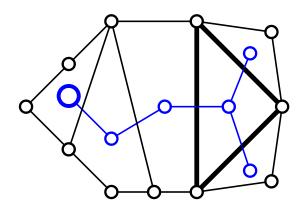


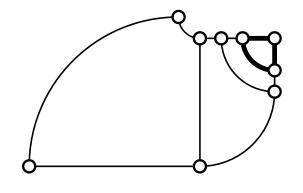
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.



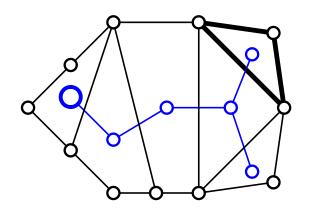


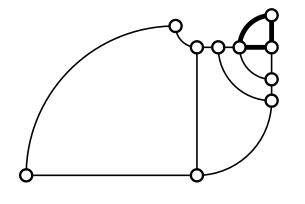
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.



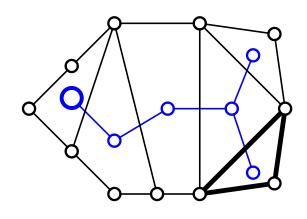


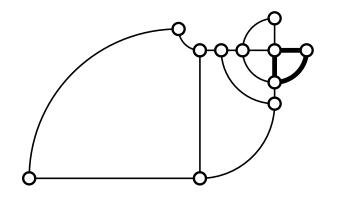
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.



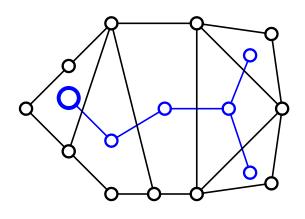


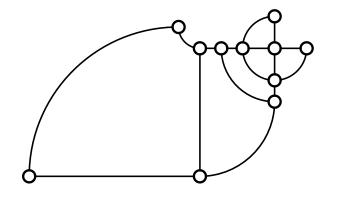
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.





Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.





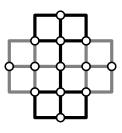
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.

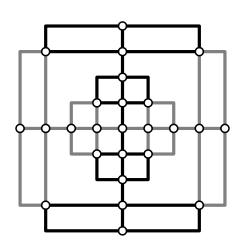
Any triconnected 3-planar graph admits an SC<sub>1</sub>-layout.

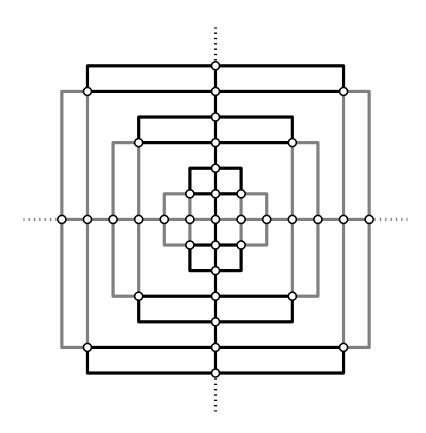
Any biconnected 4-outerplanar graph admits an SC<sub>1</sub>-layout.

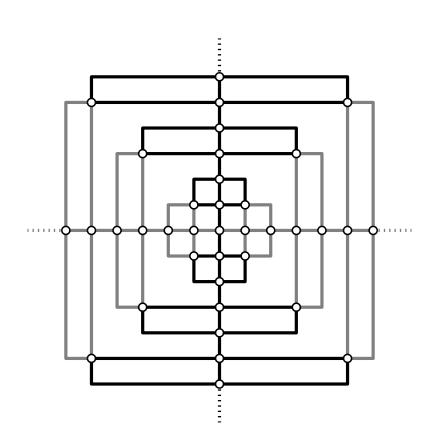
Any triconnected 3-planar graph admits an SC<sub>1</sub>-layout.

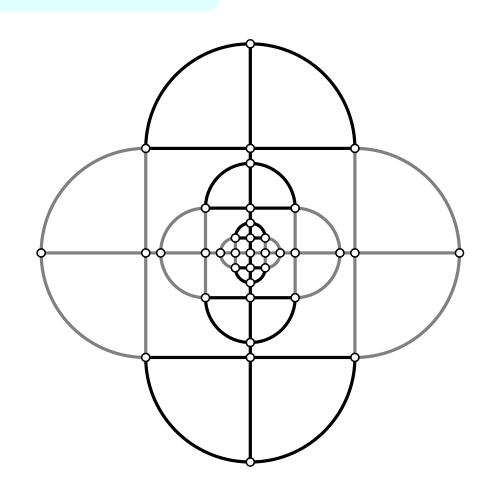
Any Hamiltonian 3-planar graph admits an SC<sub>1</sub>-layout.









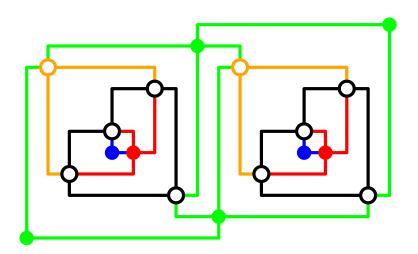


# Biconnected Graphs without SC<sub>1</sub>-Layout

There exists a biconnected 4-planar graph that admits an  $OC_2$ -layout, but does not admit an  $SC_1$ -layout.

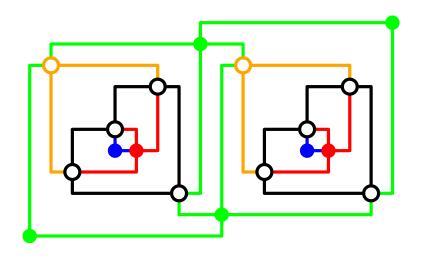
## Biconnected Graphs without SC<sub>1</sub>-Layout

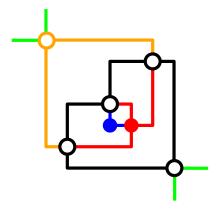
There exists a biconnected 4-planar graph that admits an  $OC_2$ -layout, but does not admit an  $SC_1$ -layout.



## Biconnected Graphs without SC<sub>1</sub>-Layout

There exists a biconnected 4-planar graph that admits an  $OC_2$ -layout, but does not admit an  $SC_1$ -layout.





 Do all 4-planar graphs admit an SC<sub>2</sub>-layout in polynomial area?

- Do all 4-planar graphs admit an SC<sub>2</sub>-layout in polynomial area?
- Do all 4-outerplanar graphs admit an SC<sub>1</sub>-layout?

- Do all 4-planar graphs admit an SC<sub>2</sub>-layout in polynomial area?
- Do all 4-outerplanar graphs admit an SC<sub>1</sub>-layout?
- Do all 3-planar graphs admit an SC<sub>1</sub>-layout?

- Do all 4-planar graphs admit an SC<sub>2</sub>-layout in polynomial area?
- Do all 4-outerplanar graphs admit an SC<sub>1</sub>-layout?
- Do all 3-planar graphs admit an SC<sub>1</sub>-layout?
- Is it NP-hard to decide whether a 4-planar graph admits an SC<sub>1</sub>-layout?