

Chair for **INFORMATICS I** Efficient Algorithms and Knowledge-Based Systems



# Drawing Graphs on Few Circles and Few Spheres

Myroslav Kryven

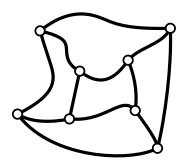
Alexander Ravsky

Alexander Wolff

Julius-Maximilians-Universität Würzburg, Germany

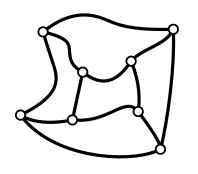
Pidstryhach Institute for Applied Problems of Mechanics and Mathematics, National Academy of Sciences of Ukraine, Lviv, Ukraine

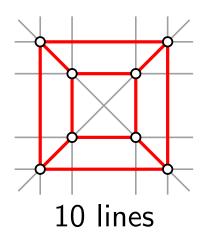
Given a planar graph,...



#### [Chaplick et al., 2016]

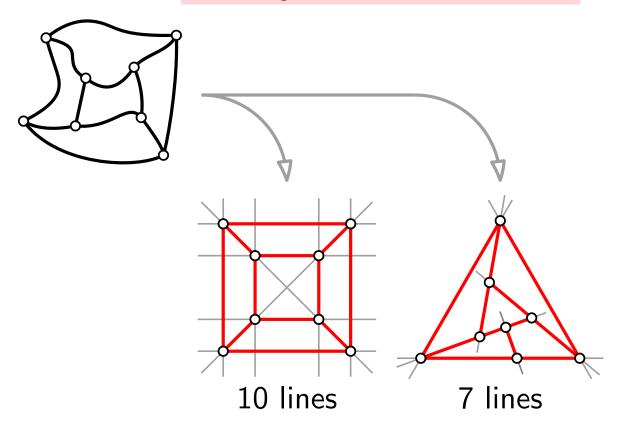
Given a planar graph,... ...find a straight-line drawing with as few lines as possible that together cover the drawing.





#### [Chaplick et al., 2016]

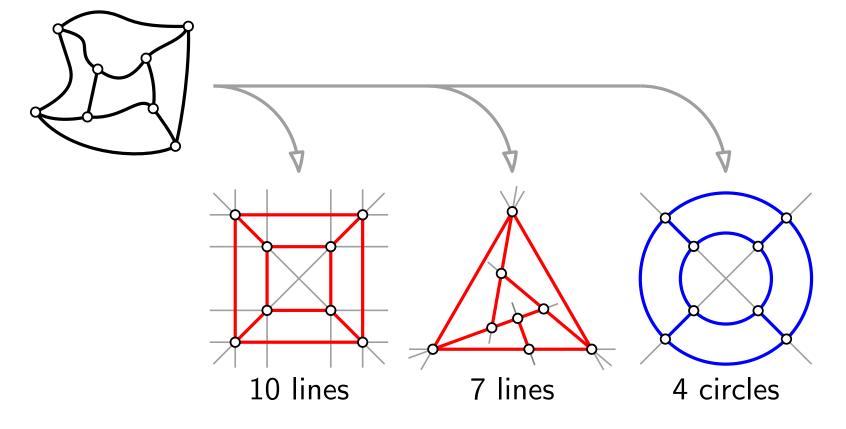
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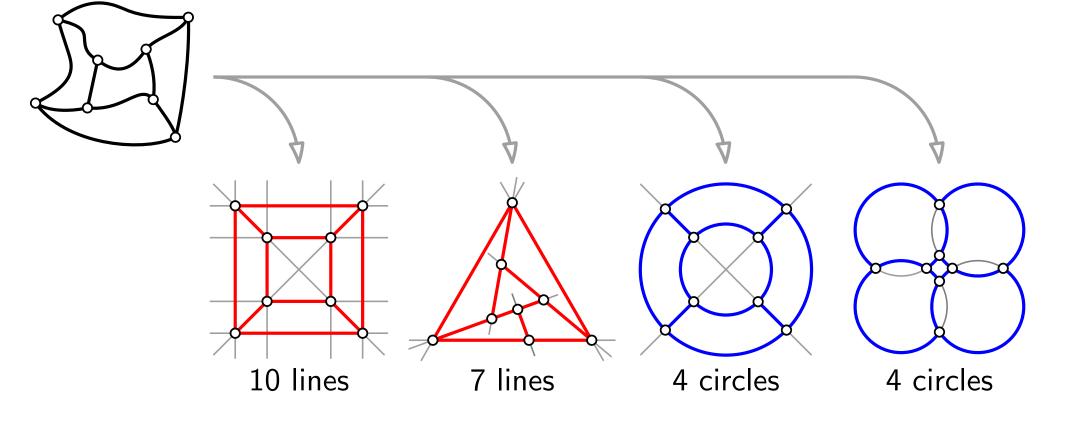
Given a planar graph,... ...find a straight-line drawing with as few lines as possible that together cover the drawing.

...find a circular-arc drawing with as few circles as possible that together cover the drawing.



#### [Chaplick et al., 2016]

Given a planar graph,... ...find a straight-line drawing with as few lines as possible that together cover the drawing. ...find a circular-arc drawing with as few circles as possible that together cover the drawing.

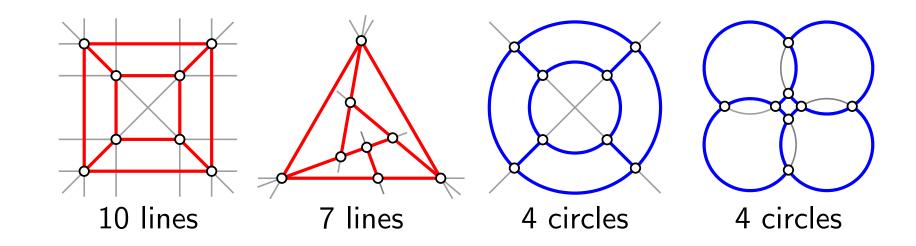


#### [Chaplick et al., 2016]

Given a planar graph,... ...find a straight-line drawing with as few lines as possible that together cover the drawing. ...find a circular-arc drawing with as few circles as possible that together cover the drawing.

#### Advantages:

• Smaller visual complexity

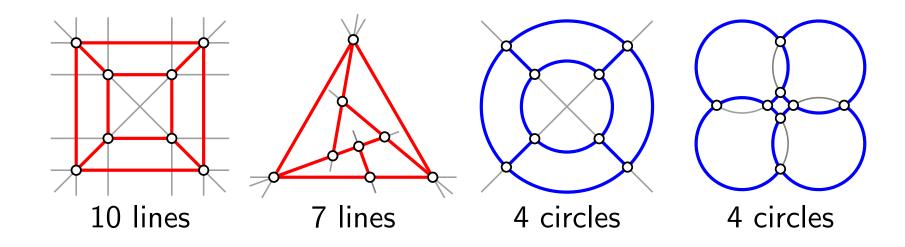


#### [Chaplick et al., 2016]

Given a planar graph,... ...find a straight-line drawing with as few lines as possible that together cover the drawing. ...find a circular-arc drawing with as few circles as possible that together cover the drawing.

#### Advantages:

- Smaller visual complexity
- Better reflects symmetry



## Outline

#### Motivation

#### Formal Definitions

A Combinatorial Lover Bound

#### Platonic solids

- affine cover number
- segment number
- spherical cover number
- arc number

#### Lower Bounds for $\sigma_d^1$ w.r.t. Other Parameters

Open Problem

[\* Chaplick et al., 2016]

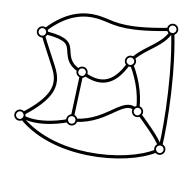
Let G be a graph, and let  $1 \le m < d$ .

**Def.** The affine cover number  $\rho_d^m(G)$  is the minimum number of *m*-dimensional hyperplanes in  $\mathbb{R}^d$ such that *G* has a crossing-free straight-line drawing that is contained in these planes.

[\* Chaplick et al., 2016]

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 $ho_2^1(\mathit{cube}) =$ 

[\* Chaplick et al., 2016]

Let G be a graph, and let  $1 \le m < d$ .

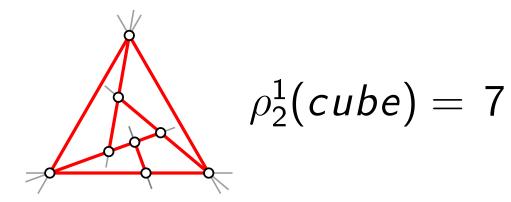
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$$ho_2^1(cube) = 7$$
  $\sigma_2^1(cube) =$ 

[\* Chaplick et al., 2016]

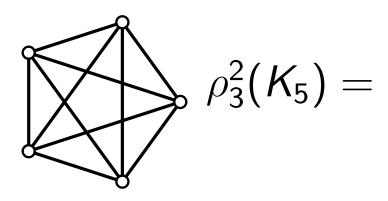
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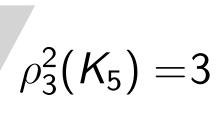
**Def.** The *affine cover number*  $\rho_d^m(G)$  is the minimum number of *m*-dimensional hyperplanes in  $\mathbb{R}^d$ such that *G* has a crossing-free straight-line drawing that is contained in these planes.

 $\rho_3^2(K_5) = 3$ 

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Let G be a graph, and let  $1 \le m < d$ .

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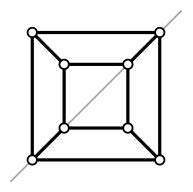
 $\sigma_3^2(K_5) = 2.$ 

### Segment Number and Arc Number

**Def.** The *segment number* of G, seg(G), is the minimum number of line segments formed by the edges of G in a straight-line drawing. [Dujmović, Eppstein,

Suderman, Wood CGTA'07]

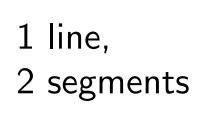
line,
 segments

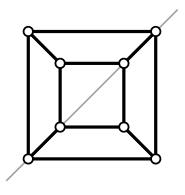


### Segment Number and Arc Number

**Def.** The *segment number* of G, seg(G), is the minimum number of line segments formed by the edges of G in a straight-line drawing. [Dujmović, Eppstein,

Suderman, Wood CGTA'07]





**Def.** The arc number of G,  $\operatorname{arc}(G)$ , is the minimum number of arcs formed by the edges of G in a circular-arc drawing. [Schulz JGAA'15]

## Outline

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#### Formal Definitions

A Combinatorial Lover Bound

#### Platonic solids

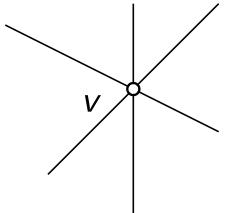
- affine cover number
- segment number
- spherical cover number
- arc number

#### Lower Bounds for $\sigma_d^1$ w.r.t. Other Parameters

Open Problem

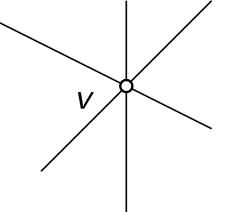
Combinatorial Lower Bounds on  $\rho_2^1$  and  $\sigma_2^1$ [Chaplick et al., 2016] Let G be a graph.

**Obs. 1** Any vertex v of G lies on  $\geq \lceil \deg(v)/2 \rceil$  lines.

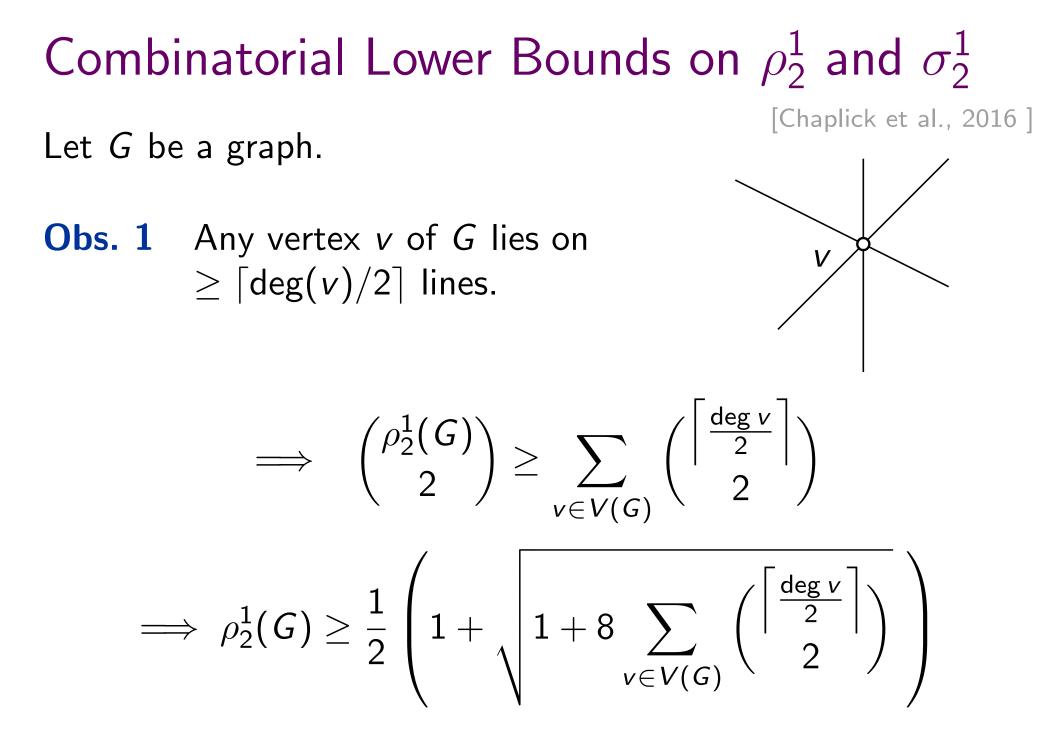


Combinatorial Lower Bounds on  $\rho_2^1$  and  $\sigma_2^1$ [Chaplick et al., 2016] Let *G* be a graph.

**Obs. 1** Any vertex v of G lies on  $\geq \lceil \deg(v)/2 \rceil$  lines.



$$\implies \left( \begin{array}{c} \rho_2^1(G) \\ 2 \end{array} \right) \ge \sum_{v \in V(G)} \left( \begin{array}{c} \left\lceil \frac{\deg v}{2} \right\rceil \\ 2 \end{array} \right)$$



Combinatorial Lower Bounds on  $\rho_2^1$  and  $\sigma_2^1$ Let G be a graph.

**Obs. 2** Any vertex v of G lies on  $> \lceil \deg(v)/2 \rceil$  circles.  $\implies 2\binom{\sigma_2^1(G)}{2} \ge \sum_{v \in \mathcal{O}} \binom{\left\lceil \frac{\deg v}{2} \right\rceil}{2}$  $\implies \sigma_2^1(G) \ge \frac{1}{2} \left( 1 + \sqrt{1 + 4 \sum_{v \in V(G)} \left( \begin{bmatrix} \frac{\deg v}{2} \\ 2 \end{bmatrix} \right)} \right)$ 

## Outline

#### Motivation

#### Formal Definitions

A Combinatorial Lover Bound

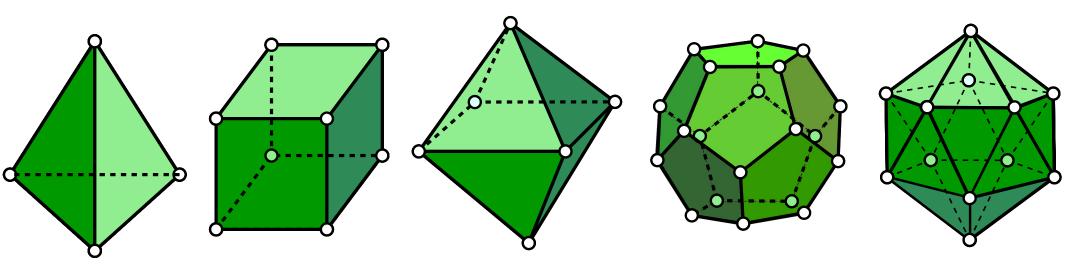
#### Platonic solids

- affine cover number
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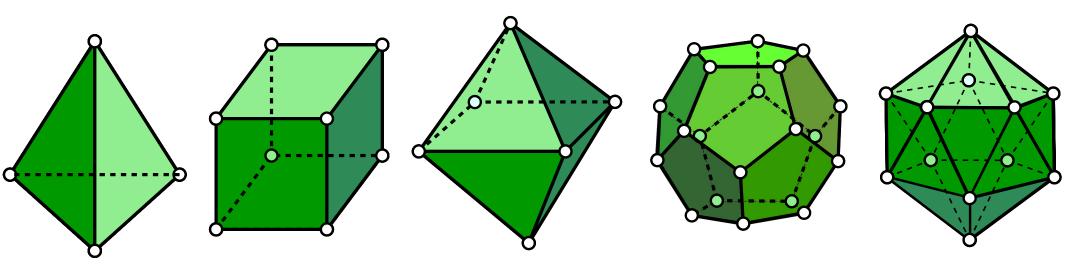
#### Lower Bounds for $\sigma_d^1$ w.r.t. Other Parameters

Open Problem

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4				
octahedron	6	12	8				
cube	8	12	6				
dodecahedron	20	30	12				
icosahedron	12	30	20				



G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4				
octahedron	6	12	8				
cube	8	12	6				
dodecahedron	20	30	12				
icosahedron	12	30	20				



G = (V, E)	V	E	F	$ ho_2^1(G$	seg( $G$ )	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4				
octahedron	6	12	8				
cube	8	12	6				
dodecahedron	20	30	12				
icosahedron	12	30	20				

#### Recall Obs. 1:

$$\rho_2^1(G) \ge \frac{1}{2} \left( 1 + \sqrt{1 + 8 \sum_{v \in V(G)} \left( \begin{bmatrix} \frac{\deg v}{2} \\ 2 \end{bmatrix} \right)} \right)$$

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	$\geq$ 4			
octahedron	6	12	8				
cube	8	12	6				
dodecahedron	20	30	12				
icosahedron	12	30	20				

Recall Obs. 1:

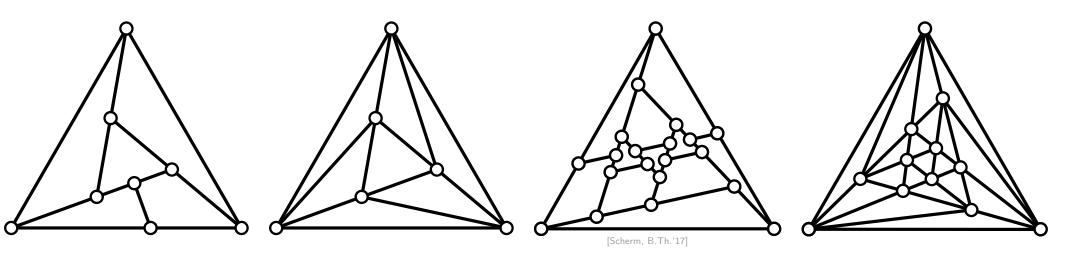
$$\rho_{2}^{1}(\textit{tetrahedron}) \geq \frac{1}{2} \left( 1 + \sqrt{1 + 8 \cdot 4 \begin{pmatrix} \left\lceil \frac{3}{2} \right\rceil \\ 2 \end{pmatrix}} \right) \geq 3.37$$

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	$\geq$ 4			
octahedron	6	12	8	$\geq$ 4			
cube	8	12	6	$\geq$ 5			
dodecahedron	20	30	12	$\geq$ 7			
icosahedron	12	30	20	$\ge$ 9			

Recall Obs. 1:

$$\implies \rho_2^1(G) \ge \frac{1}{2} \left( 1 + \sqrt{1 + 8 \sum_{v \in V(G)} \left( \begin{bmatrix} \frac{\deg v}{2} \\ 2 \end{bmatrix} \right)} \right)$$

G = (V, E)	V	E	F	$\rho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	6			
octahedron	6	12	8	9			
cube	8	12	6	7			
dodecahedron	20	30	12	910			
icosahedron	12	30	20	$13 \dots 15$			



**Arguments**: We use the number of nested cycles and the internal degree of the outer face.

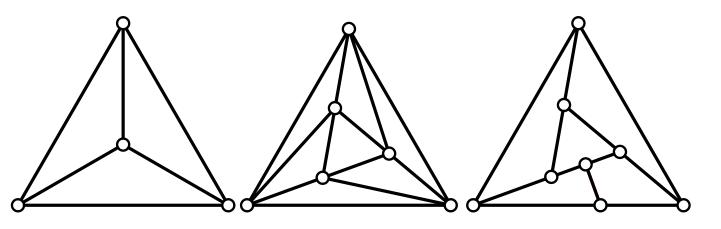
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icosahedron	12	30	20	1315			

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
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dodecahedron	20	30	12	910	$\geq$ 9		
icosahedron	12	30	20	$13 \dots 15$	$\geq$ 13		

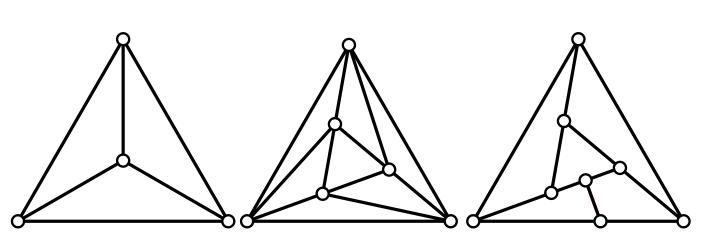
Trivial bound:

 $ho_1^2(G) \leq \mathrm{seg}(G)$ 

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
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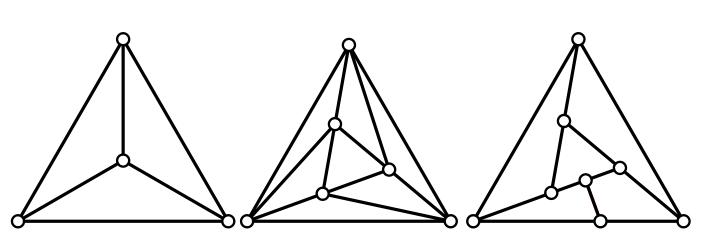
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ILP (For fixed embedding.) Find *locally consistent* angle assignment with maximum number of π-angles.

# Platonic Solids: Segment Numbers

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	6	6		
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cube	8	12	6	7	7		
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icosahedron	12	30	20	$13 \dots 15$	$\geq 15$		

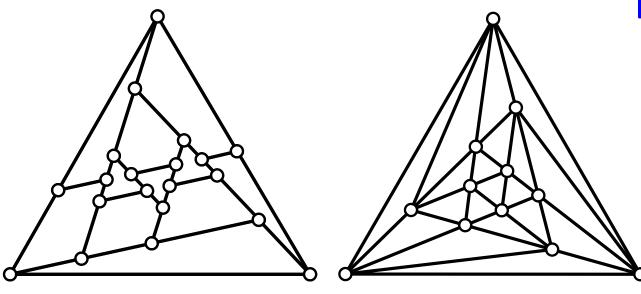


ILP (For fixed embedding.) Find *locally consistent* angle assignment with maximum number of π-angles.

⇒ Lower bounds for
 the minimum number
 of segments in the
 corresponding drawing.

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cube	8	12	6	7	7		
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icosahedron	12	30	20	$13 \dots 15$	15		



13 segments

ILP (For fixed embedding.) Find *locally consistent* angle assignment with maximum number of π-angles.

⇒ Lower bounds for
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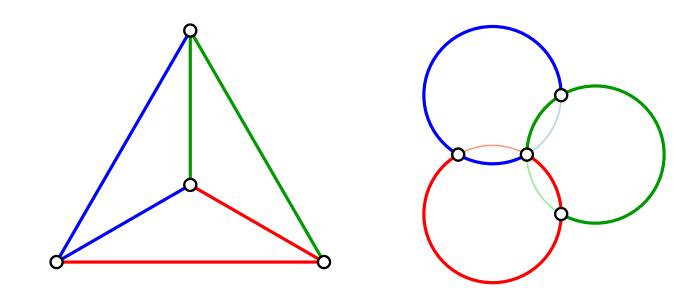
<sup>15</sup> segments

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	6	6	$\geq$ 3	
octahedron	6	12	8	9	9	$\geq$ 3	
cube	8	12	6	7	7	$\geq$ 4	
dodecahedron	20	30	12	910	13	$\geq$ 5	
icosahedron	12	30	20	$13 \dots 15$	15	$\geq$ 7	

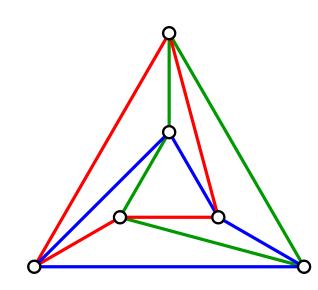
Recall Obs. 2:

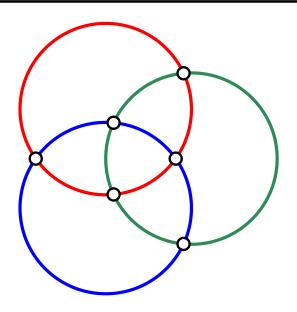
$$\implies \sigma_2^1(G) \ge \frac{1}{2} \left( 1 + \sqrt{1 + 4 \sum_{v \in V(G)} \left( \begin{bmatrix} \frac{\deg v}{2} \\ 2 \end{bmatrix} \right)} \right)$$

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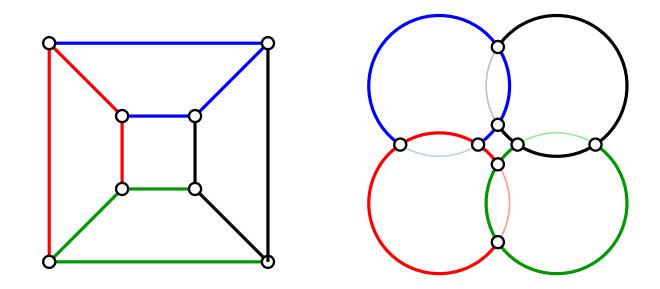


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dodecahedron	20	30	12	910	13	$\geq$ 5	
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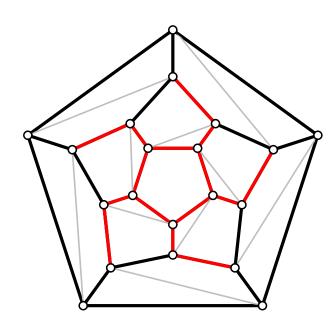


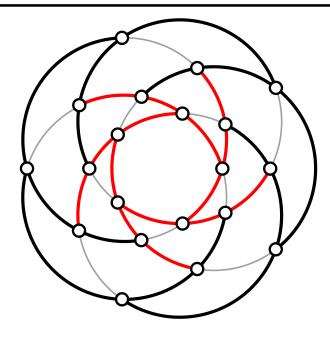


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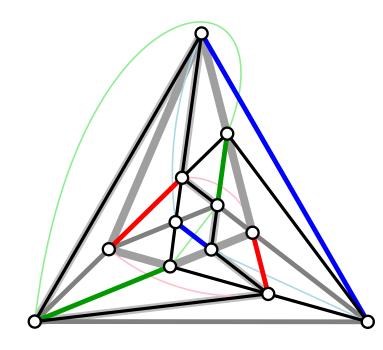
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dodecahedron	20	30	12	910	13	5	
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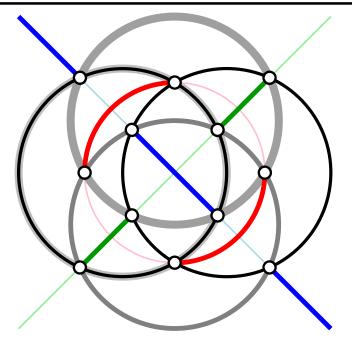




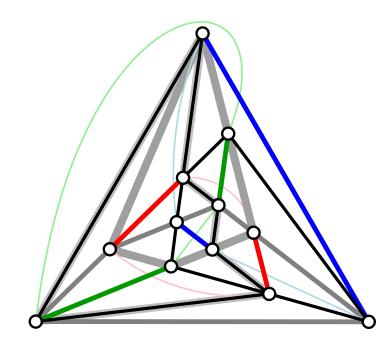
[André Schulz, JGAA'15]

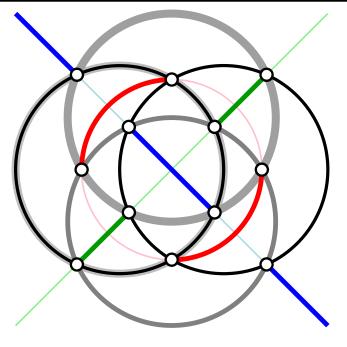
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dodecahedron	20	30	12	910	13	5	
icosahedron	12	30	20	$13 \dots 15$	15	7	





7 circles / 10 arcs

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
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icosahedron	12	30	20	1315	15	7	

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	6	6	3	<u>&gt; 3</u>
octahedron	6	12	8	9	9	3	$\geq$ 3
cube	8	12	6	7	7	4	$\geq$ 4
dodecahedron	20	30	12	910	13	5	$\geq$ 5
icosahedron	12	30	20	$13 \dots 15$	15	7	$\geq$ 7

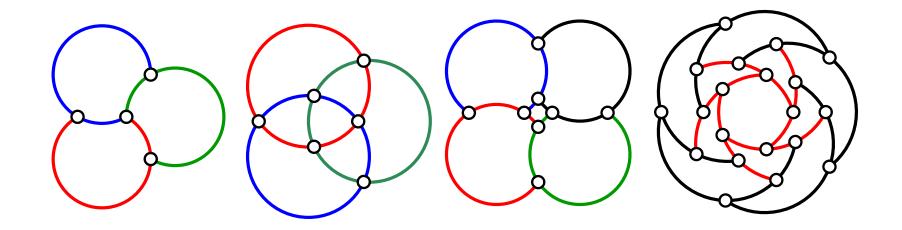
Trivial bound:  $\sigma_1^2(G) \leq \operatorname{arc}(G)$ 

G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
tetrahedron	4	6	4	6	6	3	<u>&gt; 3</u>
octahedron	6	12	8	9	9	3	$\geq$ 3
cube	8	12	6	7	7	4	$\geq$ 4
dodecahedron	20	30	12	910	13	5	$\geq 10$
icosahedron	12	30	20	$13 \dots 15$	15	7	$\geq$ 7

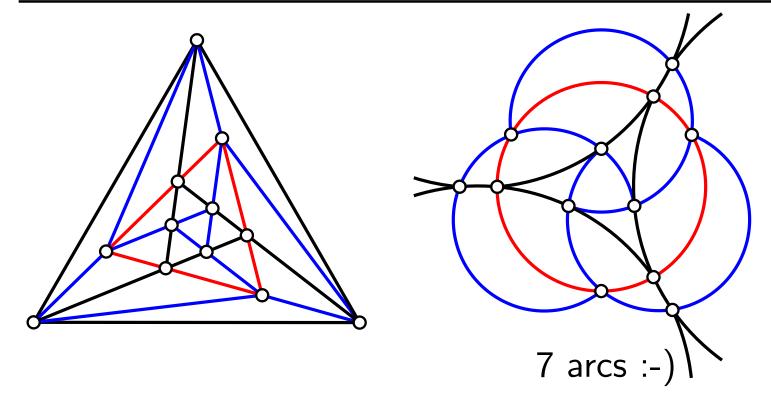
Trivial bound:  $\sigma_1^2(G) \leq \operatorname{arc}(G)$ 

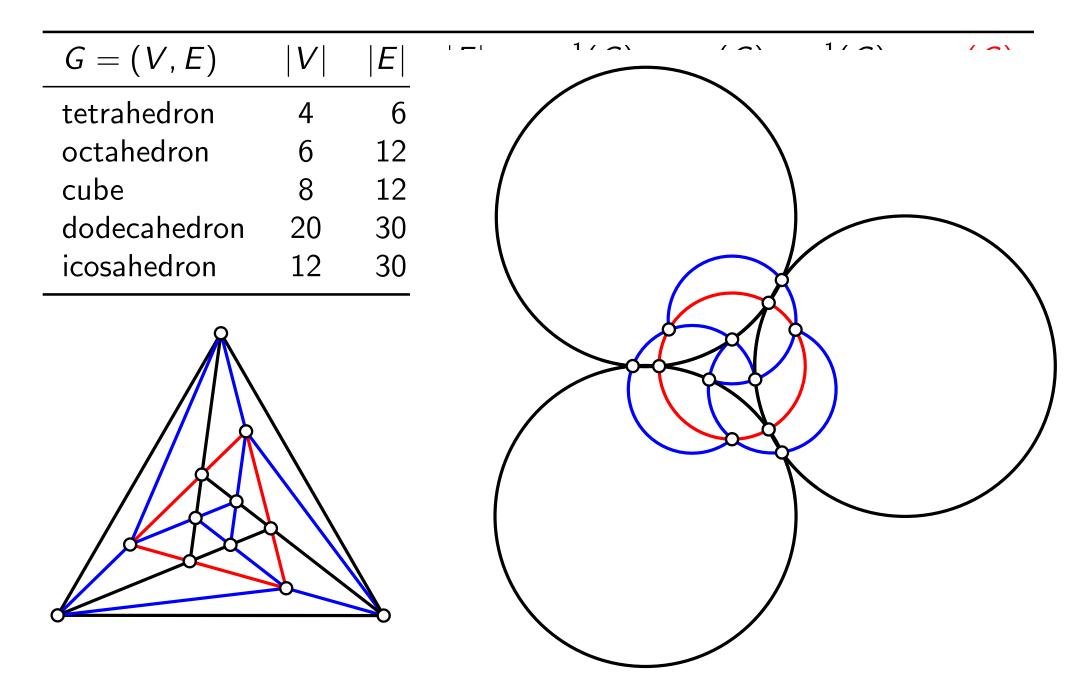
**Obs**: For any graph G,  $\operatorname{arc}(G) \ge \#(\operatorname{odd-deg. vtc. of } G)/2$ [Dujmović, Eppstein, Suderman, Wood CGTA'07]

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tetrahedron	4	6	4	6	6	3	3
octahedron	6	12	8	9	9	3	3
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dodecahedron	20	30	12	910	13	5	10
icosahedron	12	30	20	1315	15	7	$\geq$ 7

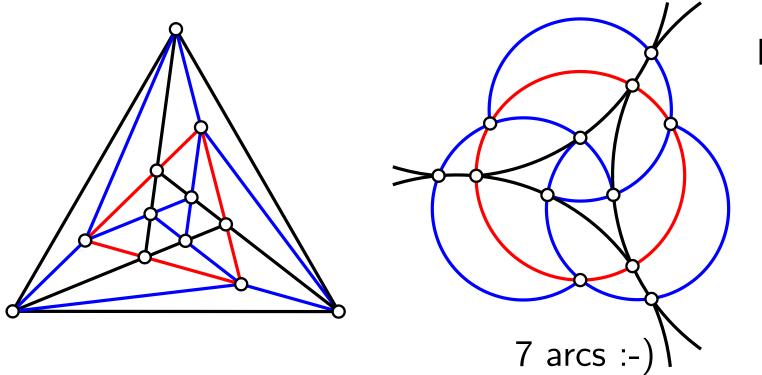


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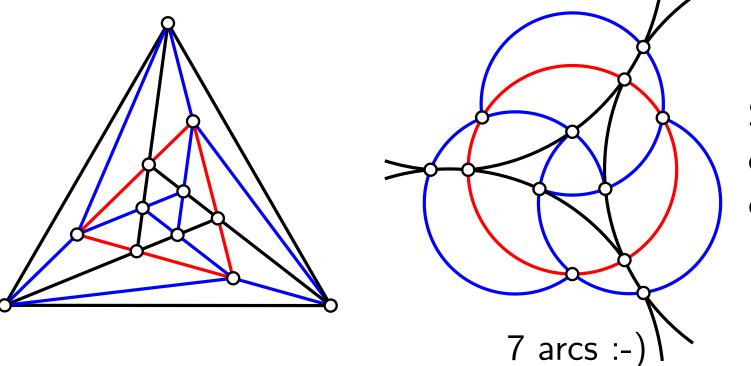


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How to draw?

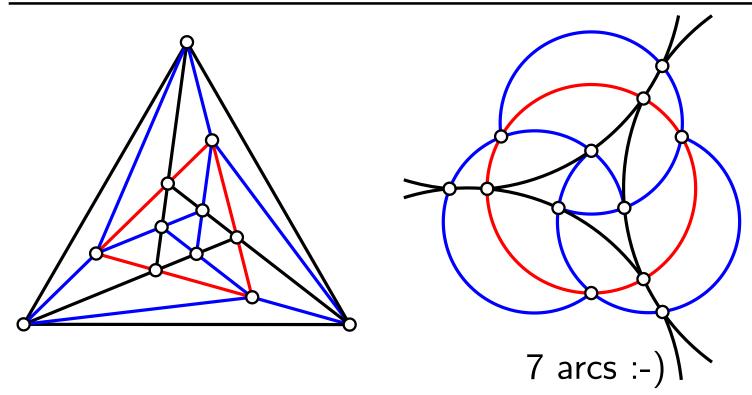
G = (V, E)	V	E	F	$ ho_2^1(G)$	seg(G)	$\sigma_2^1(G)$	$\operatorname{arc}(G)$
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How to draw?

Solve a system of quadratic equations.

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dodecahedron	20	30	12	910	13	5	10
icosahedron	12	30	20	$13 \dots 15$	15	7	7



How to draw?

Solve a system of quadratic equations.

Solution exists!

# Outline

#### Motivation

#### Formal Definitions

A Combinatorial Lover Bound

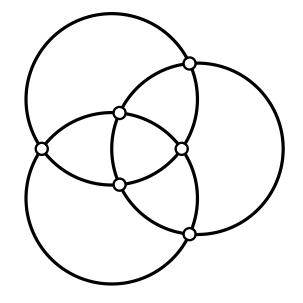
#### Platonic solids

- affine cover number
- segment number
- spherical cover number
- arc number

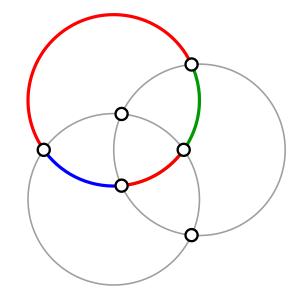
#### Lower Bounds for $\sigma_d^1$ w.r.t. Other Parameters

#### Open Problem

Edge-chromatic #  $\sigma_d^1(G) \ge \chi_e(G)/3$ ,



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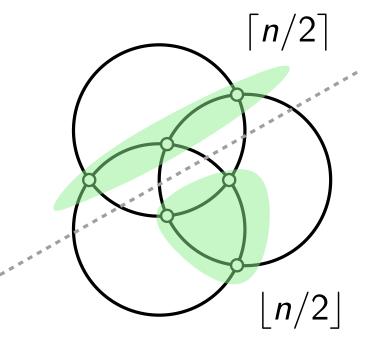


Edge-chromatic #

$$\sigma^1_d(G) \ge \chi_e(G)/3$$
,

bisection width

 $\sigma^1_d(G) \geq \mathsf{bw}(G)/2$ ,



Edge-chromatic #  $\sigma^1_d(G) \ge \chi_e(G)/3$ ,

bisection width

linear arboricity

 $\sigma^1_d(G) \ge \mathsf{bw}(G)/2,$ 

 $\sigma^1_d(G) \geq \frac{2}{3} \operatorname{la}(G),$ 

balanced separator

 $\sigma^1_d(G) \ge \sup_W(G)/2,$ 

for almost all G cubic  $\sigma_d^1(G) > n/10$ , treewidth  $\sigma_d^1(G) \ge tw(G)/6$ .

# Outline

#### Motivation

#### Formal Definitions

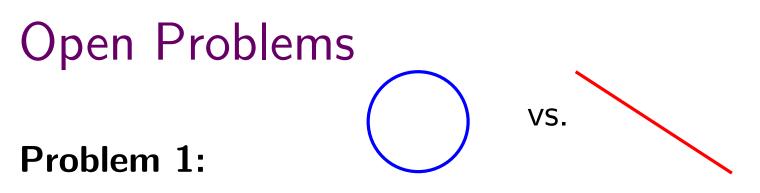
A Combinatorial Lover Bound

#### Platonic solids

- affine cover number
- segment number
- spherical cover number
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#### Lower Bounds for $\sigma_d^1$ w.r.t. Other Parameters

#### **Open Problem**



Is there a family of planar graphs whose circle cover number grows asymptotically more slowly than their line cover number?

# Open Problems Vs.

Is there a family of planar graphs whose circle cover number grows asymptotically more slowly than their line cover number?

#### Problem 2:

Determine the line cover number for the dodecahedron and icosahedron.

