## Selecting the Aspect Ratio of a Scatter Plot

# Based on Its Delaunay Triangulation 

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Joint work with

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- ... reveal trends ...



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- .. . heavily rely on the chosen aspect ratio.


WHENEVER SOMEONE UPLOADS A LETTERBOXED 16:9 VDEO RESCALED TO 4:3, I DO THIS TO THER CAR.

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task:
automatically select a good aspect ratio



## Previous Work

- aspect-ratio selection for line charts



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results depend on initial aspect ratio


## Our Approach

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- optimization criteria:
- maximize smallest angle
- minimize total edge length
- optimize compactness of triangles
- etc.



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- runtime: $\Theta(k n \log n)$
- approximation? which intermediate ratios?


## Overview

1. Maintaining the Delaunay Triangulation
2. Maximizing the Smallest Angle
3. Minimizing the Total Edge Length
4. Other Optimization Criteria

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criterion: empty circumcircle of 4 points easy to check


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- $O\left(n^{2}\right)$ flips
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- total runtime: $O\left(n^{2} \log n\right)$ for traversing all topologically different Delaunay triangulations


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- traverse lower envelope
- Davenport-Schinzel sequences \& [Agarwall + Sharir, 1995]: yields globally optimal aspect ratio in $O\left(n^{2} \log n\right)$ time


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- also works for other optimization criteria


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perimeter $\sqrt{\text { area }}$
- more:
- maximize mean inradius
- minimize sum of squared angles


## User Study

- What do users want?
- let participants choose

Please participate: www1.informatik.uni-wuerzburg.de/scatterplots

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Adjust aspect ratio and accept

Accept current drawing
I can't decide-skip this drawing
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- 18 tested instances, e.g. ...


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- maximize minimum angle



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- maximize mean inradius

|  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |

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preliminary results



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- more than one good quality measure

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