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Probing a Set of Trajectories to Maximize Captured Movement

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Motivation



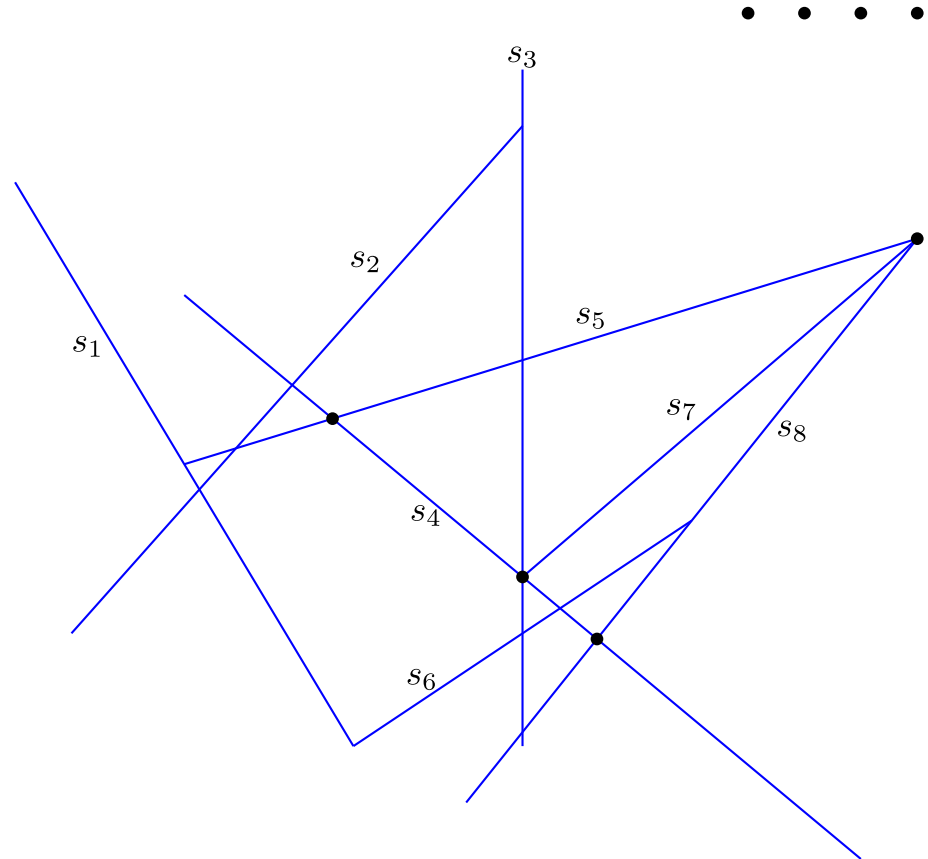
Trajectory Capturing Problem

Given:

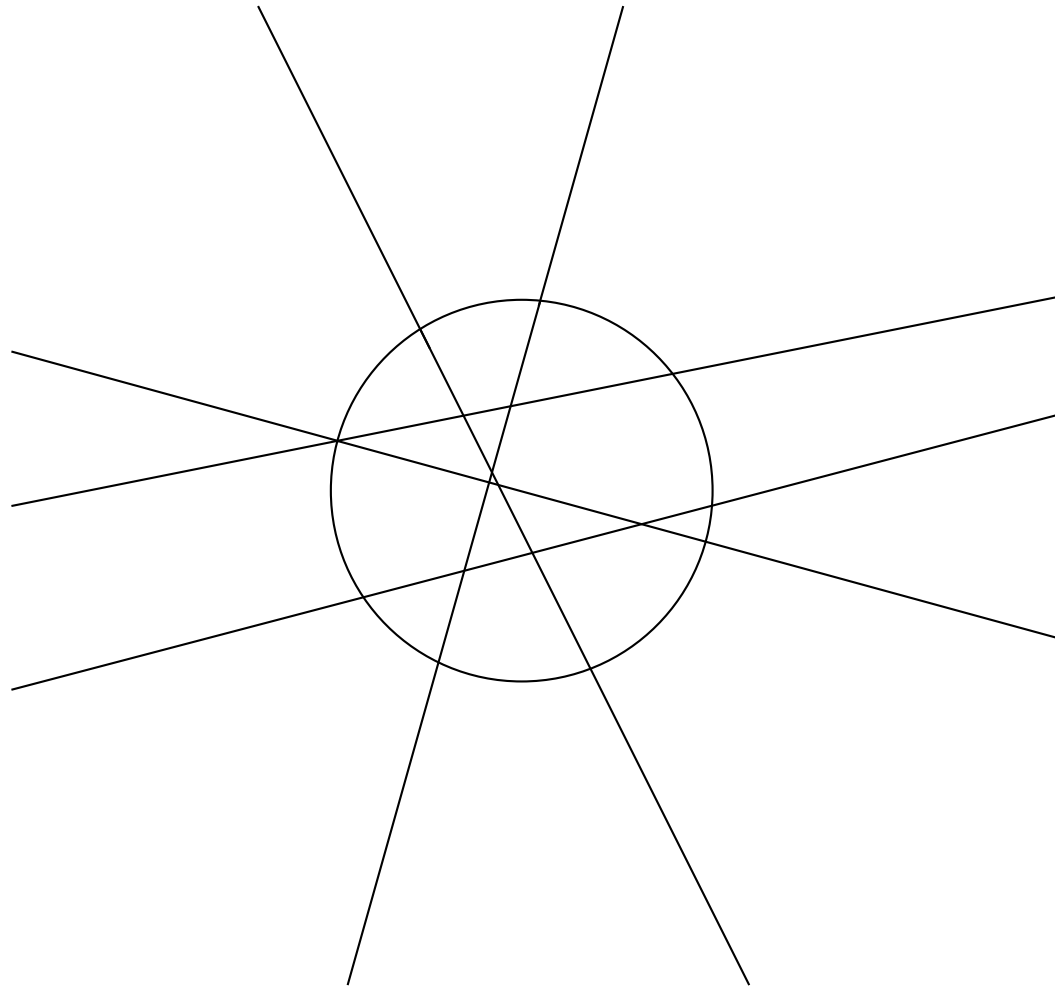
- Set $\mathcal{T}$ of trajectories
- $k \geq 2$ portals/pins

Wanted:

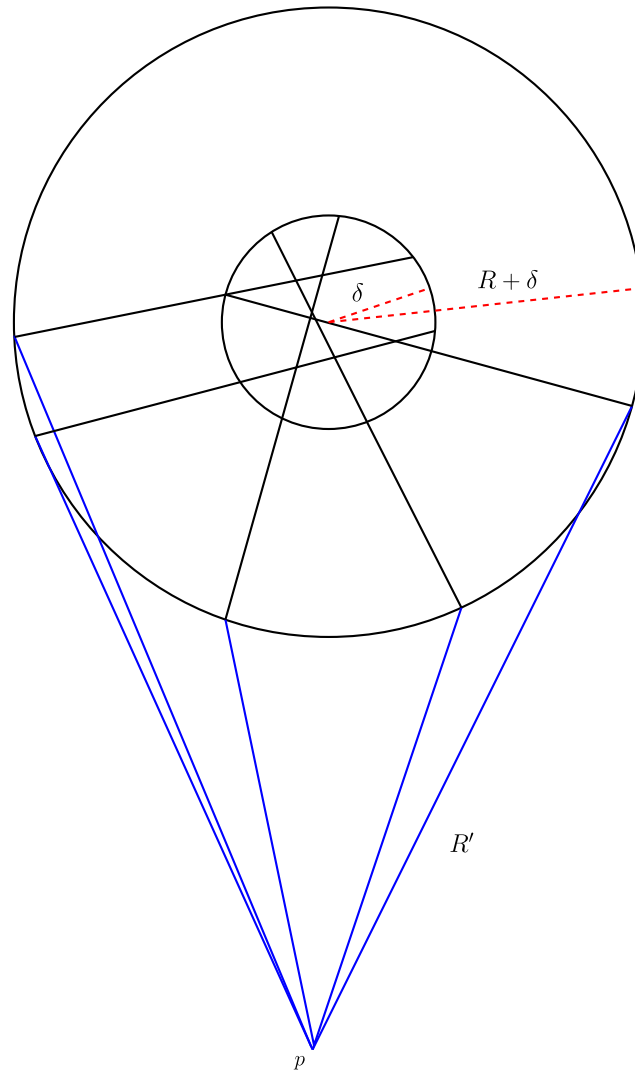
- Place k portals
- maximize the total length
- Measured between two portals



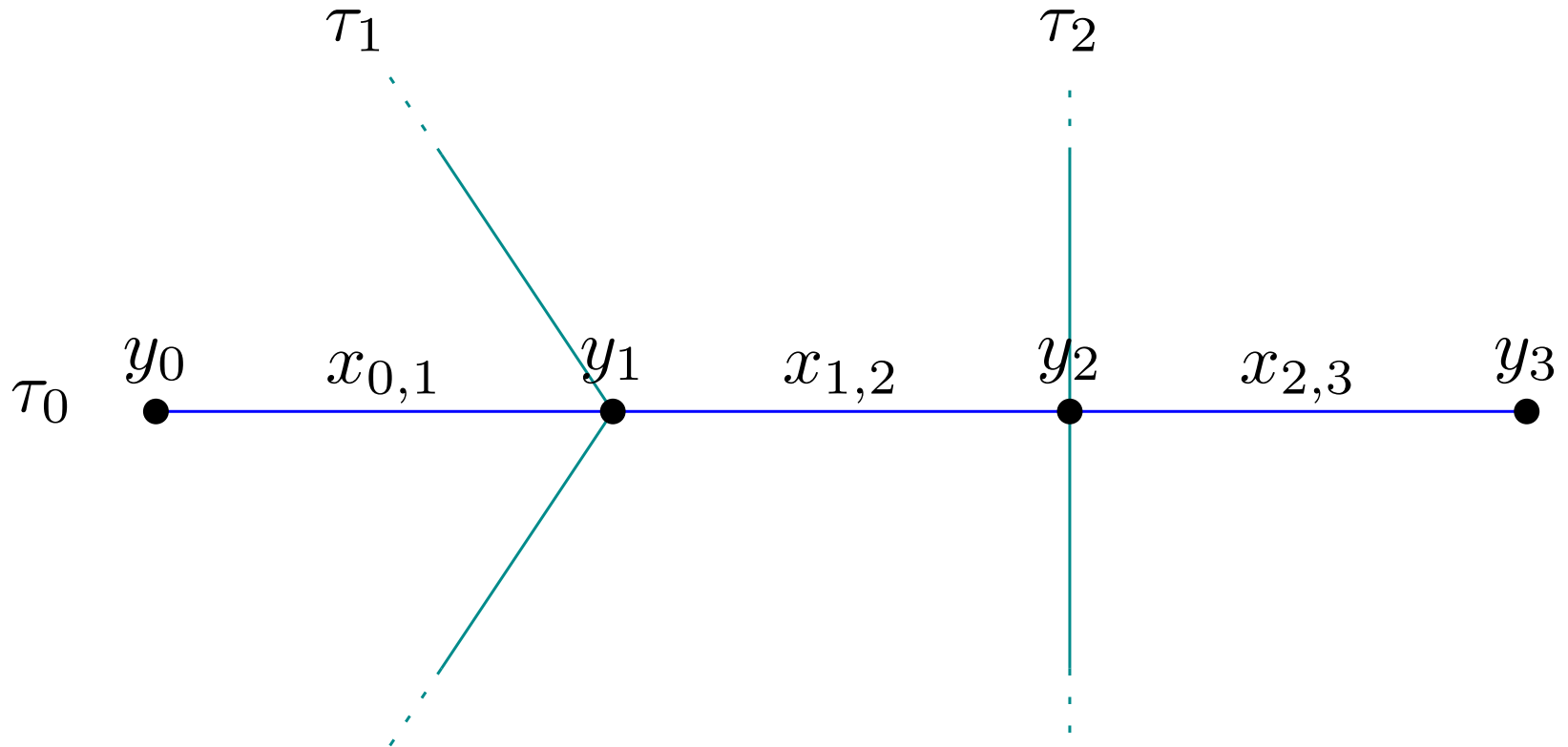
Hardness proof: Geometric graphs



Hardness proof: Geometric graphs



Integer Programming Preliminaries

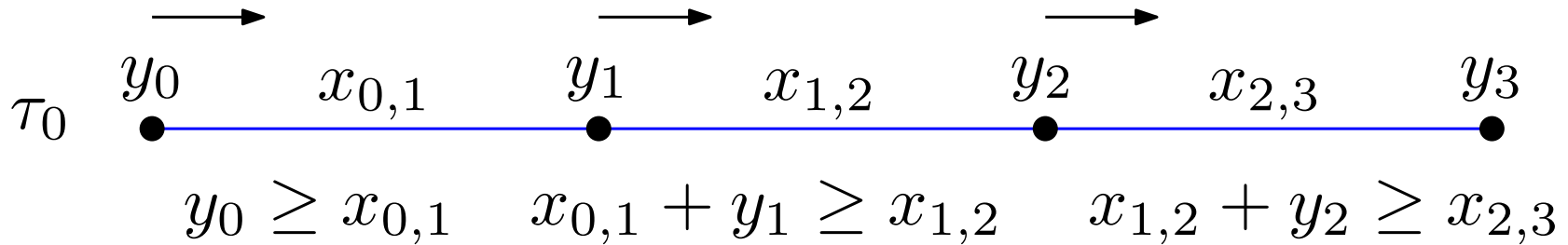


Integer Programming

$$\begin{aligned} \max \quad & \sum_{\tau \in \mathcal{T}, e \in E(\tau)} f(e) x_{\tau, e} \\ & \sum_{v \in V} y_v \leq k \end{aligned}$$

(Constraint 1)

Integer Programming

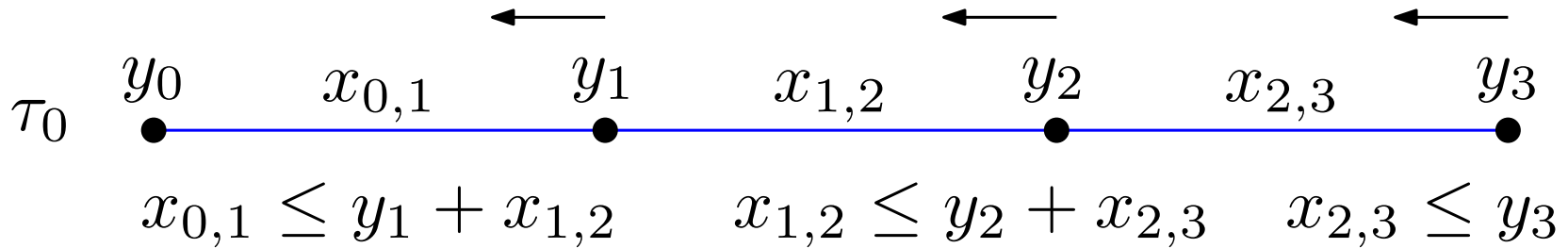


Generalization:

$$\forall \tau = (v_0, \dots, v_l) \in \mathcal{T} :$$

$$\forall i \in 0, \dots, l-1 : \quad \begin{cases} x_{\tau, v_i v_{i+1}} & \leq y_i & \text{if } i = 0, \\ x_{\tau, v_i v_{i+1}} & \leq y_i + x_{\tau, v_{i-1} v_i} & \text{else} \end{cases}$$

Integer Programming



Generalization:

$$\forall \tau = (v_0, \dots, v_l) \in \mathcal{T} :$$

$$\forall i \in 1, \dots, l : \quad \begin{cases} x_{\tau, v_{i-1} v_i} & \leq & y_i & \text{if } i = l, \\ x_{\tau, v_{i-1} v_i} & \leq & y_i + x_{\tau, v_i v_{i+1}} & \text{else} \end{cases}$$

Integer Programming

$$\max \sum_{\tau \in \mathcal{T}, e \in E(\tau)} f(e) x_{\tau, e}$$

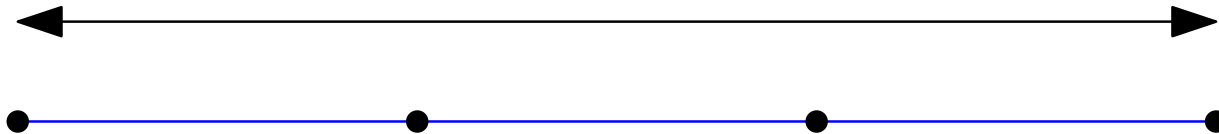
$$\sum_{v \in V} y_v \leq k \quad (\text{Constraint 1})$$

$$\forall \tau = (v_0, \dots, v_l) \in \mathcal{T} :$$

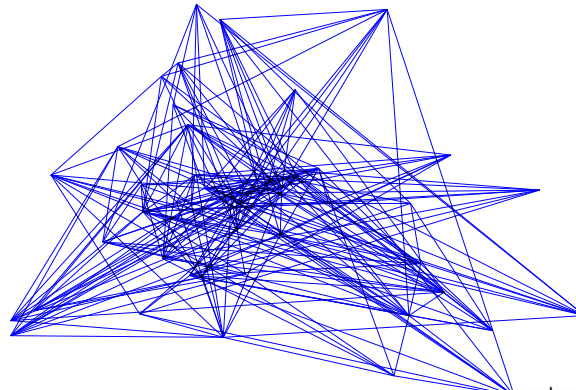
$$\forall i \in 0, \dots, l-1 : \quad \begin{cases} x_{\tau, v_i v_{i+1}} \leq y_i & \text{if } i = 0, \\ x_{\tau, v_i v_{i+1}} \leq y_i + x_{\tau, v_{i-1} v_i} & \text{else} \end{cases} \quad (\text{Constraint 2})$$

$$\forall i \in 1, \dots, l : \quad \begin{cases} x_{\tau, v_{i-1} v_i} \leq y_i & \text{if } i = l, \\ x_{\tau, v_{i-1} v_i} \leq y_i + x_{\tau, v_i v_{i+1}} & \text{else} \end{cases} \quad (\text{Constraint 3})$$

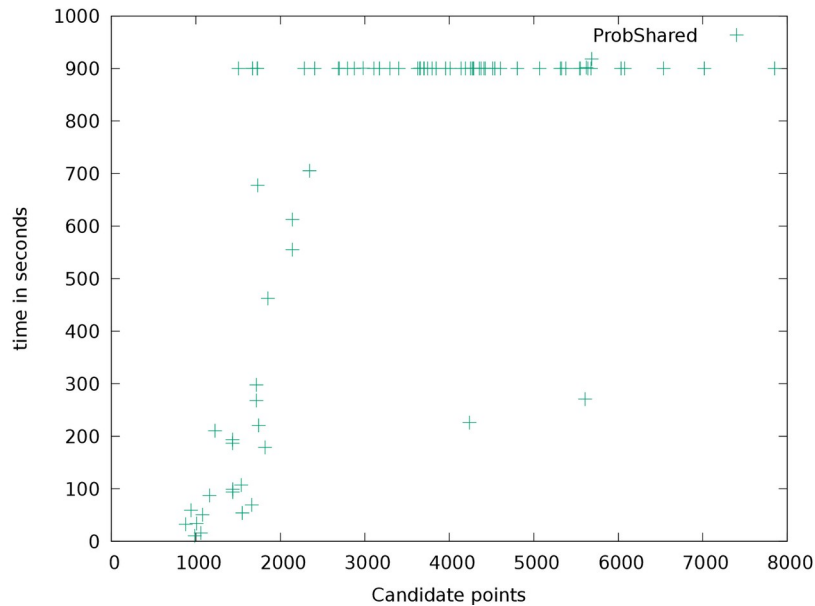
$$\forall v \in V, \tau \in e \in \tau : \quad x_{\tau, e}, y_v \in \{0, 1\}$$



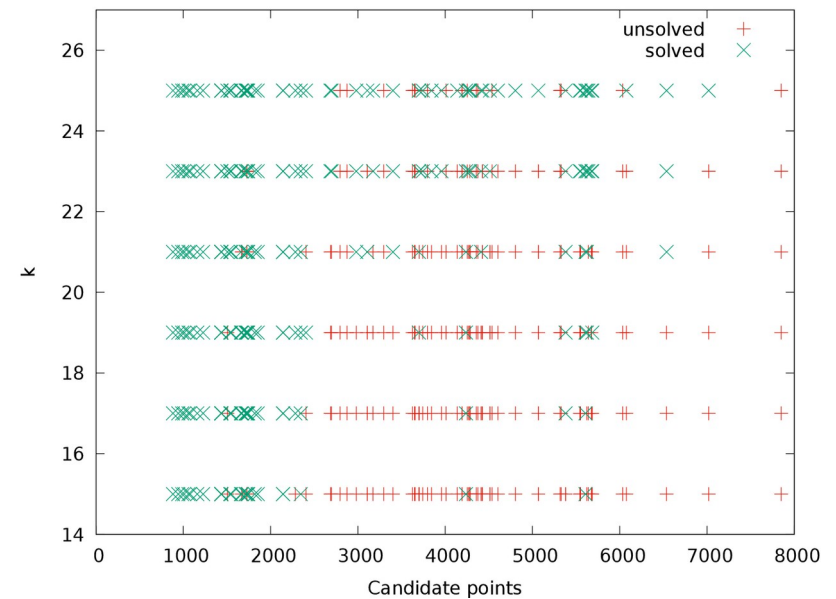
Evaluation: Segments between points



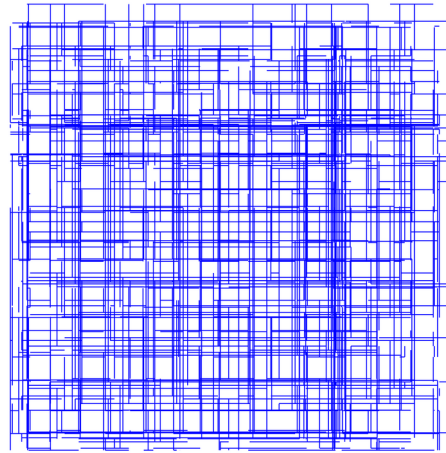
time consumption for $k = 15$



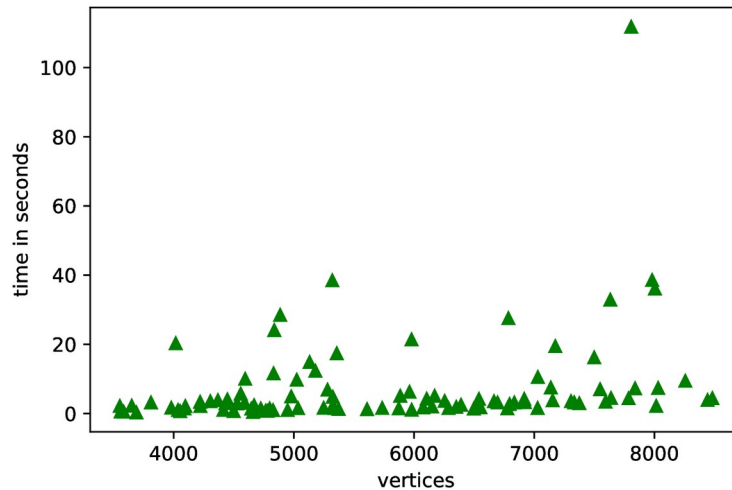
solved and unsolved instances comparison



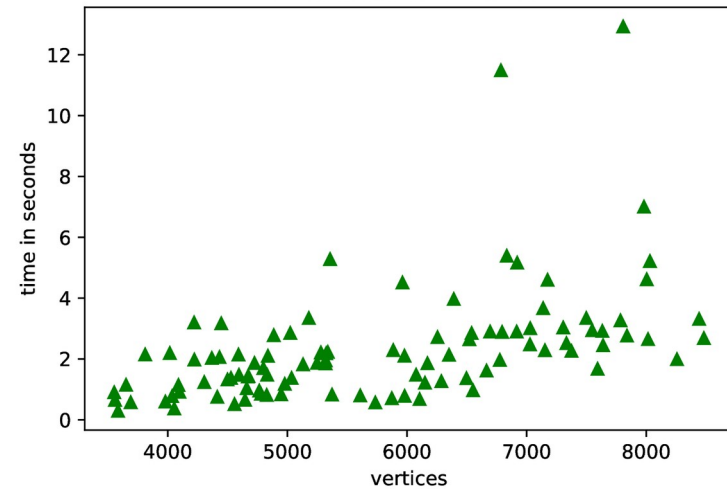
Evaluation: Non-overlapping axis-aligned



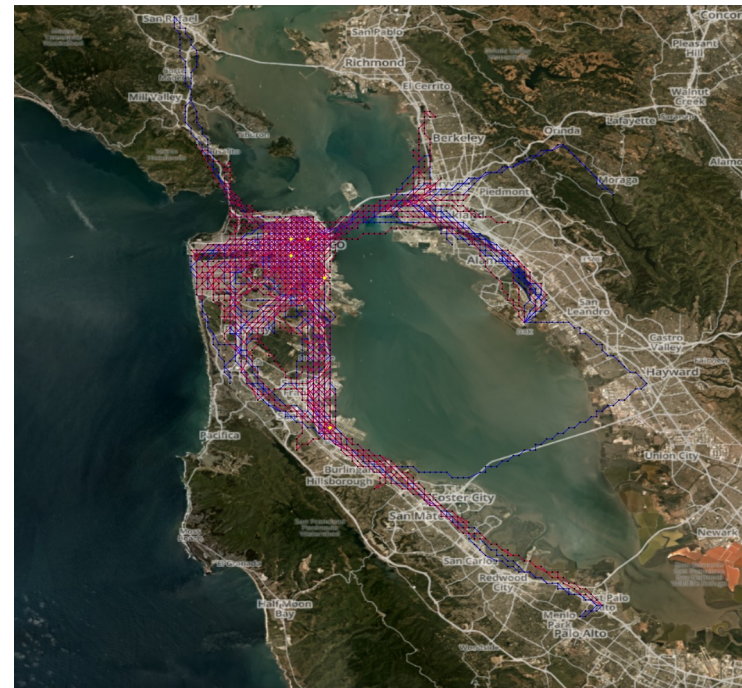
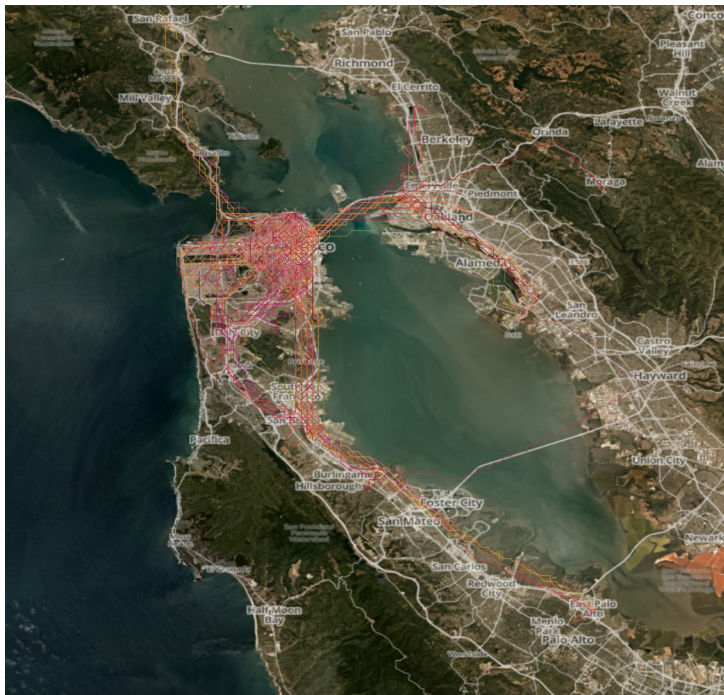
IP time results for $k=5$



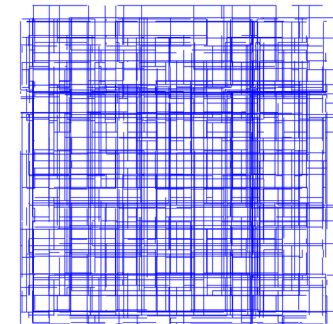
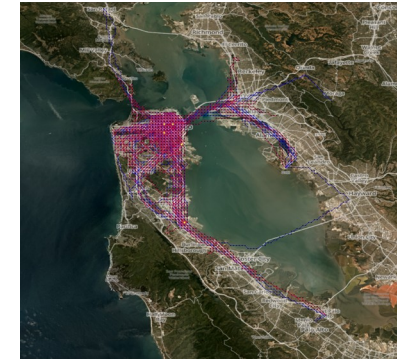
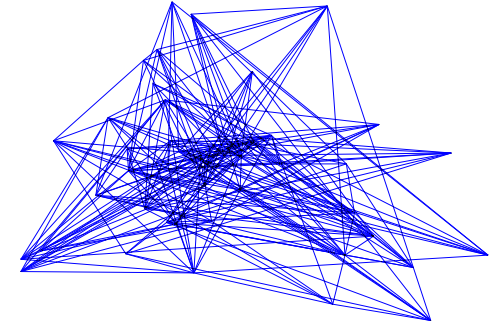
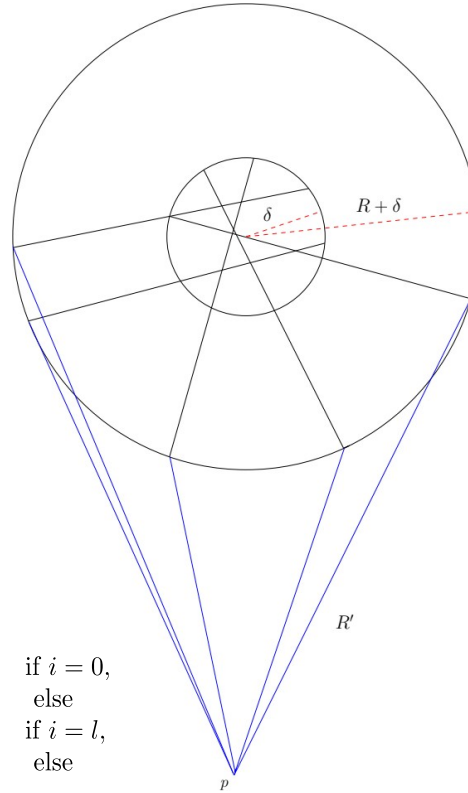
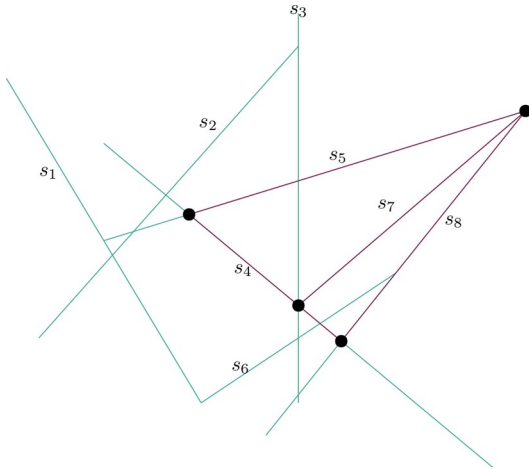
IP time results for $k=10$



Evaluation: Taxi Data



Conclusion



$$\max \sum_{\tau \in \mathcal{T}, e \in E(\tau)} f(e) x_{\tau, e}$$

$$\sum_{v \in V} y_v \leq k$$

$$\forall \tau = (v_0, \dots, v_l) \in \mathcal{T} :$$

$$\forall i \in 0, \dots, l-1 :$$

$$\forall i \in 1, \dots, l :$$

$$\forall v \in V, \tau \in e \in \tau :$$

$$\begin{cases} x_{\tau, v_i v_{i+1}} \leq y_i & \text{if } i = 0, \\ x_{\tau, v_i v_{i+1}} \leq y_i + x_{\tau, v_{i-1} v_i} & \text{else} \\ x_{\tau, v_{i-1} v_i} \leq y_i & \text{if } i = l, \\ x_{\tau, v_{i-1} v_i} \leq y_i + x_{\tau, v_i v_{i+1}} & \text{else} \\ x_{\tau, e}, y_v \in \{0, 1\} \end{cases}$$